# PRECONDITIONED AOR ITERATIVE METHODS FOR SOLVING MULTI-LINEAR SYSTEMS WITH $\mathcal{M}$-TENSOR ${ }^{\dagger}$ 

MENG QI, XINHUI SHAO*


#### Abstract

Some problems in engineering and science can be equivalently transformed into solving multi-linear systems. In this paper, we propose two preconditioned AOR iteration methods to solve multi-linear systems with -tensor. Based on these methods, the general conditions of preconditioners are given. We give the convergence theorem and comparison theorem of the two methods. The results of numerical examples show that methods we propose are more effective.


AMS Mathematics Subject Classification : 65H05, 65F10.
Key words and phrases : M-tensor, multi-linear systems, iteration method, AOR type method, preconditioner.

## 1. Introduction

We consider the multi-linear system

$$
\begin{equation*}
\mathcal{A} \mathbf{x}^{m-1}=\mathbf{b} \tag{1.1}
\end{equation*}
$$

where $\mathcal{A}=\left(a_{i i_{2} \cdots i_{m}}\right)$ is an order m dimension n tensor, $\mathbf{x}$ and $\mathbf{b}$ are n dimensional vectors. The tensor-vector product $\mathcal{A} \mathbf{x}^{m-1}$ is defined by

$$
\begin{equation*}
\left(\mathcal{A} \mathbf{x}^{m-1}\right)_{i}=\sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}, i=1,2, \cdots, n \tag{1.2}
\end{equation*}
$$

where $\mathbf{x}=\left(x_{1}, x_{2}, \cdots, x_{n}\right)^{T}$. There are some practical applications of the multilinear systems in engineering and science fields [1,2], for instance, numerical partial differential equations [3], tensor complementarity problems [4], data mining [5], tensor absolute value equations [6] and so on.

[^0]Ding and Wei [3] proved the multi-linear system (1.1) always has a unique positive solution if $\mathcal{A}$ is a strong $\mathcal{M}$-tensor and $\mathbf{b}$ is a positive vector. In [12], the authors gave the definition of tensor splitting $\mathcal{A}=\mathcal{E}-\mathcal{F}$. The tensor splitting method for solving the system (1.1) is defined by

$$
\mathbf{x}_{k}=\left[M(\mathcal{E})^{-1} \mathcal{F} \mathbf{x}_{k-1}^{m-1}+M(\mathcal{E}) \mathbf{b}\right]^{\left[\frac{1}{m-1}\right]}, \quad k=1,2, \cdots, n
$$

According to this method, Li et al. [13] proposed preconditioned multi-linear systems based on the preconditioned technique of linear systems. Cui et al. [14] proposed a new preconditioner to solve the system (1.1).

Li, Liu and Vong [13] considered the preconditioner

$$
P_{\alpha}=I+S_{\alpha}=\left[\begin{array}{ccccc}
1 & -\alpha_{1} a_{12 \cdots 2} & 0 & \cdots & 0 \\
0 & 1 & -\alpha_{2} a_{23 \cdots 3} & \cdots & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & 0 & \cdots & -\alpha_{n-1} a_{n-1, n \cdots n} \\
0 & 0 & 0 & \cdots & 1
\end{array}\right]
$$

Cui, Li and Song [14] proposed a new preconditioner x

$$
P_{\max }=\left(I+S_{\max }\right)
$$

where $S_{\text {max }}$ is defined by

$$
S_{\max }=\left(s_{i, k_{i}}^{m}\right)=\left\{\begin{array}{cc}
-a_{i, k_{i}, \cdots k_{i}}, & i=1, \cdots n-1, k_{i}>i, \\
0, & \text { otherwise },
\end{array}\right.
$$

where $k_{i}=\min \left\{j\left|\max _{j}\right| a_{i j, \cdots j} \mid, i<n, j>i\right\}$.
In this paper, $\mathcal{A}$ is a strong $\mathcal{M}$-tensor and $\mathbf{b}>\mathbf{0}$. Without loss of generality, we assume that the all diagonal entries of $\mathcal{A}$ are 1. The preconditioned multi-linear system is $P \mathcal{A} \mathbf{x}^{m-1}=P \mathbf{b}$ where $P$ is a nonsingular and nonnegative matrix with unit diagonal entries. Let $\widehat{\mathcal{A}}=P \mathcal{A}$ and $\widehat{\mathbf{b}}=P \mathbf{b}$.

## 2. AOR iterative method

### 2.1. Proposed method.

The order $m$ dimension $n$ unit tensor is denoted by $\mathcal{I}_{m}$. The majorization matrix of tensor $\mathcal{A}$ is denoted by $M(\mathcal{A})$ and $M(\mathcal{A})_{i j}=a_{i j \cdots j}, i, j=1,2, \cdots, n$. Let $\widehat{\mathcal{A}}=\widehat{\mathcal{D}}-\widehat{\mathcal{L}}-\widehat{\mathcal{F}}$, where $\widehat{\mathcal{D}}=\widehat{D} \mathcal{I}_{m}, \widehat{\mathcal{L}}=\widehat{L} \mathcal{I}_{m}, \widehat{D}$ and $-\widehat{L}$ are the diagonal and strictly lower triangular parts of $M(\widehat{\mathcal{A}})$, respectively. The matrix $P$ is nonsingular and nonnegative with unit diagonal entries. Then

$$
\widehat{\mathcal{A}}=\left(\hat{a}_{i i_{2} \cdots i_{m}}\right)=\sum_{k=1}^{n} p_{i k} a_{k i_{2} \cdots i_{m}}
$$

We define

$$
M(\mathcal{A})=I-L-U
$$

$$
\begin{gathered}
P=I+P_{1}+P_{2} \\
P_{1} U=E_{1}+F_{1}+G_{1} \\
P_{2} L=E_{2}+F_{2}+G_{2}
\end{gathered}
$$

where $E_{1}$ and $E_{2}$ are diagonal matrixes, $L, P_{1}, F_{1}$ and $F_{2}$ are strictly lower triangular matrixes, $U, P_{2}, G_{1}$ and $G_{2}$ are strictly upper triangular matrixes. It is obvious that the matrices mentioned above are nonnegative.

$$
\begin{gathered}
\widehat{\mathcal{D}}=\left(I-E_{1}-E_{2}\right) \mathcal{I}_{m} \\
\widehat{\mathcal{L}}=\left(L-P_{1}+P_{1} L+F_{1}+F_{2}\right) \mathcal{I}_{m} \\
\widehat{\mathcal{F}}=\left(U-P_{2}+P_{2} U+G_{1}+G_{2}\right) \mathcal{I}_{m}+P \mathcal{F}
\end{gathered}
$$

Hadjidimos [18] proposed an accelerated overrelaxation (AOR) iterative method to solve linear systems. Based on this method, we propose the following AOR iterative algorithm:

$$
\mathbf{x}_{k}^{[m-1]}=(\widehat{D}-r \widehat{L})^{-1}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] \mathbf{x}_{k-1}^{m-1}+w(\widehat{D}-r \widehat{L})^{-1} \widehat{\mathbf{b}}
$$

$k=1,2, \cdots, n$, where $\mathbf{x}_{k}{ }^{[m-1]}=\left(x_{1}^{m-1}, x_{2}^{m-1}, \cdots, x_{n}^{m-1}\right)^{T}, w$ and $r$ are real parameters with $0 \leq r \leq w \leq 1(w \neq 0)$. The iteration tensor of AOR iterative methods is

$$
\widehat{\mathcal{T}}_{r, w}=(\widehat{D}-r \widehat{L})^{-1}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] .
$$

### 2.2. Convergence analysis of the proposed method.

First, we give the following lemma to show that $\widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor.
Lemma 2.1. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. If $P=\left(p_{i j}\right)$ is a nonsingular and nonnegative matrix with $p_{i i}=1$ and

$$
\sum_{k=1}^{n} p_{i k} a_{k j \cdots j} \leq 01 \leq i \neq j \leq n
$$

then $\widehat{\mathcal{A}}=P \mathcal{A}$ is a strong $\mathcal{M}$-tensor.
Proof. When $\left(i_{2}, \cdots, i_{m}\right) \neq(j, \cdots, j)$, since $p_{i j} \geq 0$, we have $\sum_{k=1}^{n} p_{i k} a_{k i_{2} \cdots i_{m}} \leq$ 0. Noticing that $\sum_{k=1}^{n} p_{i k} a_{k j \cdots j} \leq 0$ for $1 \leq i \neq j \leq n$, we obtain $\sum_{k=1}^{n} p_{i k} a_{k i_{2} \cdots i_{m}} \leq$ 0 for $\left(i, i_{2}, \cdots, i_{m}\right) \neq(i, i, \cdots, i)$. Consequently, $\widehat{\mathcal{A}}$ is a $\mathcal{Z}$-tensor.

According to the theorem 3 in [15], we assume that $\mathbf{x} \geq \mathbf{0}$ and $\mathcal{A} \mathbf{x}^{m-1}>\mathbf{0}$. It is obvious that $P \geq 0$, then we have $P \mathcal{A} \mathbf{x}^{m-1} \geq 0$. Therefore, there exists $\mathbf{x} \geq \mathbf{0}$ such that $P \mathcal{A} \mathbf{x}^{m-1} \geq 0$. By the theorem 3 in $[15], \widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor.

We give the splitting $\widehat{\mathcal{A}}=\widehat{\mathcal{M}}-\widehat{\mathcal{N}}$, where $\widehat{\mathcal{N}}=\frac{1}{w}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+$ $w \widehat{\mathcal{F}}]$ and $\widehat{\mathcal{M}}=\frac{1}{w}(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})$. By the definition and related theorems of tensor splitting in [12], we have the follow theorem 2.2.

Theorem 2.2. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. If $P=\left(p_{i j}\right)$ is a nonsingular and nonnegative matrix with $p_{i i}=1$ and

$$
\sum_{k=1}^{n} p_{i k} a_{k j \cdots j} \leq 0,1 \leq i \neq j \leq n
$$

then $\widehat{\mathcal{A}}=\widehat{\mathcal{M}}-\widehat{\mathcal{N}}$ is a convergent regular splitting for $0 \leq r \leq w \leq 1(w \neq 0)$.
Proof. By lemma 2.1, $\widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor. It is obvious that $\widehat{\mathcal{F}} \geq \mathcal{O}$ and $\widehat{L} \geq \mathbf{0}$.Accordingly, $\widehat{\mathcal{N}}=\frac{1}{w}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] \geq 0$ for $0 \leq r \leq w \leq$ $1(w \neq 0)$. By the Neumann series, we have

$$
\begin{aligned}
& M(\widehat{\mathcal{M}})^{-1}=w(\widehat{D}-r \widehat{L})^{-1}=w\left(I-r \widehat{D}^{-1} \widehat{L}\right)^{-1} \widehat{D}^{-1} \\
= & w\left[I+r \widehat{D}^{-1} \widehat{L}+\left(r \widehat{D}^{-1} \widehat{L}\right)^{2}+\left(r \widehat{D}^{-1} \widehat{L}\right)^{3}+\cdots+\left(r \widehat{D}^{-1} \widehat{L}\right)^{n-1}\right] \widehat{D}^{-1}
\end{aligned}
$$

By the theorem 3 [15] and proposition 4 [15], we get $0<\sum_{k=1}^{n} p_{i k} a_{k i \cdots i} \leq 1$, i.e., $\widehat{D}^{-1} \geq I$. It is easy to know that $M(\widehat{\mathcal{M}})^{-1} \geq \mathbf{0}$ and $\widehat{\mathcal{A}}=\widehat{\mathcal{M}}-\widehat{\mathcal{N}}$ is a regular splitting. By the lemma 3.16 in [12], $\widehat{\mathcal{A}}=\widehat{\mathcal{M}}-\widehat{\mathcal{N}}$ is a convergent splitting.

According to the theorem 5.4 [12] and lemma 2.1, $\widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor and the AOR iterative method is convergent.

### 2.3. The comparison theorem.

Before the comparison theorem, we give the following lemma first.
Lemma 2.3. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. Then there exists $\varepsilon_{0}>0$ such that, for any $0<\varepsilon<\varepsilon_{0}, \mathcal{A}(\varepsilon)=\left(a_{i i_{2} \cdots i_{m}}(\varepsilon)\right)$ is also a strong $\mathcal{M}$-tensor, where

$$
a_{i i_{2} \cdots i_{m}}(\varepsilon)=\left\{\begin{array}{cl}
a_{i i_{2} \cdots i_{m}}, & a_{i i_{2} \cdots i_{m}} \neq 0 \\
-\varepsilon, & a_{i i_{2} \cdots i_{m}}=0
\end{array}\right.
$$

Proof. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor, it is not hard to check that $\mathcal{A}(\varepsilon)$ is a $\mathcal{Z}$ tensor. By the theorem 3 in [15], there exists $\mathbf{x} \geq \mathbf{0}$ such that $\mathcal{A} \mathbf{x}^{m-1}>\mathbf{0}$. Thus,

$$
\sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}>0 i=1,2, \cdots, n
$$

Assuming that $\delta=\sum_{i_{2}, \cdots, i_{m}=1}^{n} x_{i_{2}} \cdots x_{i_{m}}>0$, let

$$
\varepsilon_{0}=\frac{1}{\delta} \min \left\{\sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}, i=1,2, \cdots, n\right\} .
$$

Then, we have $\varepsilon_{0}>0$ and

$$
\sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}-\delta \varepsilon_{0} \geq 0, i=1,2, \cdots, n
$$

For any $0<\varepsilon<\varepsilon_{0}$, we obtain

$$
\begin{aligned}
& \sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}}(\varepsilon) x_{i_{2}} \cdots x_{i_{m}} \\
> & \sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}-\varepsilon \sum_{i_{2}, \cdots, i_{m}=1}^{n} x_{i_{2}} \cdots x_{i_{m}} \\
\geq & \sum_{i_{2}, \cdots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}-\delta \varepsilon_{0} \\
\geq & i=1,2, \cdots n
\end{aligned}
$$

Thus $\mathcal{A}(\varepsilon) \mathbf{x}^{m-1}>0$. By the theorem 3 in [15], we have $\mathcal{A}(\varepsilon)=\left(a_{i i_{2} \cdots i_{m}}(\varepsilon)\right)$ is also a strong $\mathcal{M}$-tensor.

Let $\mathcal{A}=\mathcal{I}_{m}-\mathcal{L}-\mathcal{F}$, where $\mathcal{L}=L \mathcal{I}_{m},-L$ is the strictly lower triangular part of $M(\widehat{\mathcal{A}})$. We give the following splittings:

$$
\begin{aligned}
& \mathcal{A}=\mathcal{M}-\mathcal{N}=\frac{1}{w}\left(\mathcal{I}_{m}-r \mathcal{L}\right)-\frac{1}{w}\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right], \\
& \widehat{\mathcal{A}}=\widehat{\mathcal{M}}-\widehat{\mathcal{N}}=\frac{1}{w}(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})-\frac{1}{w}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] .
\end{aligned}
$$

The iteration tensor:

$$
\begin{gathered}
\mathcal{T}_{r, w}=M(\mathcal{M}) \mathcal{N}=(I-r L)^{-1}\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right], \\
\widehat{\mathcal{T}}_{r, w}=M(\widehat{\mathcal{M}}) \widehat{\mathcal{N}}=(\widehat{D}-r \widehat{L})^{-1}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] .
\end{gathered}
$$

The comparison theorem of spectral radius between preconditioned AOR iterative method with and without preconditioner is given.

Theorem 2.4. Let $\mathcal{A}$ be an $\mathcal{M}$-tensor. If $P=\left(p_{i j}\right)$ is a nonsingular and nonnegative matrix with $p_{i i}=1$ and

$$
\sum_{k=1}^{n} p_{i k} a_{k j \cdots j} \leq 0,1 \leq i \neq j \leq n
$$

we have $\widehat{\mathcal{T}}_{r, w} \leq \mathcal{T}_{r, w}<1$ for $0 \leq r \leq w \leq 1(w \neq 0)$.
Proof. Since $\mathcal{A}$ is a $\mathcal{M}$-tensor, we can get $\mathcal{N}=\frac{1}{w}\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right] \geq$ $\mathcal{O}$. By Neumann series, we have $M(\mathcal{M})^{-1} \geq \mathbf{0}$. Thus, $M(\mathcal{M})^{-1} \mathcal{N}$ is a nonnegative tensor. By theorem 2.2, $M(\widehat{\mathcal{M}})^{-1} \widehat{\mathcal{N}}$ is a nonnegative tensor. According to theorem 1.3 in [22], there exists a nonnegative vector $\mathbf{x} \neq \mathbf{0}$ such that

$$
M(\mathcal{M})^{-1} \mathcal{N} \mathbf{x}^{m-1}=\lambda \mathbf{x}^{[m-1]} \text { and } 0 \leq \rho\left(\mathcal{T}_{r, w}\right)=\lambda<1
$$

or equivalently,

$$
\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right] \mathbf{x}^{m-1}=\lambda\left(\mathcal{I}_{m}-r \mathcal{L}\right) \mathbf{x}^{m-1}
$$

Then we obtain
$\mathcal{A} \mathbf{x}=M(\mathcal{M})\left(\mathcal{I}_{m}-M(\mathcal{M})^{-1} \mathcal{N}\right) \mathbf{x}^{m-1}=(1-\lambda) M(\mathcal{M}) \mathcal{I}_{m} \mathbf{x}^{m-1}=(1-\lambda) \mathcal{M} \mathbf{x}^{m-1}$.
If $\lambda>0$, we have $w-r+r \lambda \neq 0$ and $\mathcal{L} \mathbf{x}^{m-1}=\frac{(-1+w+\lambda) \mathcal{I}_{m}-w \mathcal{F}}{w-r+r \lambda} \mathbf{x}^{m-1}$.
Hence,

$$
\begin{align*}
& M(\widehat{\mathcal{M}})^{-1} \widehat{\mathcal{N}} \mathbf{x}^{m-1}-\lambda \mathbf{x}^{[m-1]} \\
= & M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}[(1-w) \widehat{\mathcal{D}}+(w-r) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}-\lambda(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})] \mathbf{x}^{m-1} \\
= & M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}[(1-w-\lambda) \widehat{\mathcal{D}}+(w-r+r \lambda) \widehat{\mathcal{L}}+w \widehat{\mathcal{F}}] \mathbf{x}^{m-1} \\
= & M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}[(1-\lambda) \widehat{\mathcal{D}}+(r \lambda-r) \widehat{\mathcal{L}}-w \widehat{\mathcal{A}}] \mathbf{x}^{m-1} \\
= & M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}[(1-\lambda) \widehat{\mathcal{D}}+\mathrm{r}(\lambda-1) \widehat{\mathcal{L}}-w P \mathcal{A}] \mathbf{x}^{m-1} \\
= & M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}\left[(1-\lambda) \widehat{\mathcal{D}}-\mathrm{r}(1-\lambda) \widehat{\mathcal{L}}-(1-\lambda) P\left(\mathcal{I}_{m}-r \mathcal{L}\right)\right] \mathbf{x}^{m-1} \\
= & (\lambda-1) M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}\left[-\left(I-E_{1}-E_{2}\right) \mathcal{I}_{m}+\mathrm{r}\left(L-P_{1}+P_{1} L+\right.\right. \\
& \left.\left.F_{1}+F_{2}\right) \mathcal{I}_{m}+\left(I+P_{1}+P_{2}\right) \mathcal{I}_{m}-r\left(L+P_{1} L+P_{2} L\right) \mathcal{I}_{m}\right] \mathbf{x}^{m-1} \\
= & (\lambda-1) M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}\left[\left(E_{1}+E_{2}+r\left(F_{1}+F_{2}\right)+(1-r) P_{1}+P_{2}\right.\right. \\
& \left.\left.-r P_{2} L\right) \mathcal{I}_{m}\right] \mathbf{x}^{m-1} \\
= & (\lambda-1) M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}\left[\left(E_{1}+E_{2}+r\left(F_{1}+F_{2}\right)+(1-r) P_{1}\right) \mathcal{I}_{m}\right. \\
& \left.+P_{2} \mathcal{I}_{m}-r P_{2} \frac{(-1+w+\lambda) \mathcal{I}_{m}-w \mathcal{F}}{w-r+r \lambda}\right] \mathbf{x}^{m-1} \\
= & (\lambda-1) M(\widehat{\mathcal{D}}-r \widehat{\mathcal{L}})^{-1}\left[\left(E_{1}+E_{2}+r\left(F_{1}+F_{2}\right)+(1-r) P_{1}\right) \mathcal{I}_{m}\right. \\
& \left.+w P_{2} \frac{(1-r) \mathcal{I}_{m}+r \mathcal{F}}{w-r+r \lambda}\right] \mathbf{x}^{m-1} . \tag{4.1}
\end{align*}
$$

Case 1: $\mathcal{A}$ is irreducible. It is easy to get that

$$
\begin{aligned}
M(\mathcal{M})^{-1} \mathcal{N} & =(I-r L)^{-1}\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right] \\
& =\left[I+r L+(r L)^{2}+\cdots+(r L)^{n-1}\right]\left[(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}\right] \\
& \geq(1-w) \mathcal{I}_{m}+(w-r) \mathcal{L}+w \mathcal{F}+r(1-w) \mathcal{L} \\
& =(1-w) \mathcal{I}_{m}+w(1-r) \mathcal{L}+w \mathcal{F}
\end{aligned}
$$

When $0 \leq r<1$, since $\mathcal{A}$ is irreducible, $M(\mathcal{M})^{-1} \mathcal{N}$ is also irreducible. By theorem 1.4 in [21], it is easy to know that $\lambda>0$ and the Perron vector $\mathbf{x}>\mathbf{0}$. From (4.1), we obtain $M(\widehat{\mathcal{M}})^{-1} \widehat{\mathcal{N}} \mathbf{x}^{m-1} \leq \lambda \mathbf{x}^{[m-1]}$. By lemma 3.2 in [13], we have

$$
\rho\left(\widehat{\mathcal{T}}_{r, w}\right) \leq \lambda=\rho\left(\mathcal{T}_{r, w}\right)
$$

When $r=w=1$,

$$
\rho\left(\widehat{\mathcal{T}}_{1,1}\right)=\lim _{r \rightarrow 1^{-}} \rho\left(\widehat{\mathcal{T}}_{r, 1}\right) \leq \lim _{r \rightarrow 1^{-}} \rho\left(\mathcal{T}_{r, 1}\right)=\rho\left(\mathcal{T}_{1,1}\right)<1
$$

Case 2: $\mathcal{A}$ is reducible. By the lemma 2.3 , there exists a positive number $\varepsilon$ such that $\mathcal{A}(\varepsilon)$ is a irreducible $\mathcal{M}$-tensor. According to the proof above, we obtain

$$
\rho\left(\widehat{\mathcal{T}}_{r, w}\right)=\lim _{\varepsilon \rightarrow 0} \rho\left(\widehat{\mathcal{T}}_{r, w}(\varepsilon)\right) \leq \lim _{\varepsilon \rightarrow 0} \rho\left(\mathcal{T}_{r, w}(\varepsilon)\right)=\rho\left(\mathcal{T}_{r, w}\right)<1
$$

Thus, $\rho\left(M(\widehat{\mathcal{M}})^{-1} \widehat{\mathcal{N}}\right) \leq \rho\left(M(\mathcal{M})^{-1} \mathcal{N}\right)<1$. The proof is completed.

## 3. Modified AOR iterative method

We give the splitting $\widehat{\mathcal{A}}=\overline{\mathcal{U}}-\overline{\mathcal{L}}$, where

$$
\overline{\mathcal{U}}=\left\{\begin{array}{cc}
\hat{a}_{i i_{2} \cdots i_{m}}, & i_{2}, \cdots, i_{m} \geq i \\
0, & \text { otherwise }
\end{array}\right.
$$

We give the definition: $\mathcal{A}_{i}=\left(a_{i i_{2} \cdots i_{m}}\right)_{i_{2}, i_{3}, \cdots, i_{m}=1}^{n}$, then

$$
\mathcal{A} \mathbf{x}^{m-1}=\left(\mathcal{A}_{1} \mathbf{x}^{m-1}, \mathcal{A}_{2} \mathbf{x}^{m-1}, \cdots, \mathcal{A}_{n} \mathbf{x}^{m-1}\right)^{T}
$$

$i=1,2, \cdots, n$. Then we propose modified AOR iterative method as follows.
Algorithm 3.1:
(1) We give $k_{\text {max }}$ as the maximum iteration steps and the precision $\varepsilon$ as the termination conditions. Then we take a positive initial vector $\mathbf{x}_{0}=\left(x_{1}^{(0)}, x_{2}^{(0)}, \cdots x_{n}^{(0)}\right)^{T}$ and let $k=1$;
(2) While $k \leq k_{\max }$;
(3) $a_{i i \cdots i}\left[x_{i}^{(k)}\right]^{m-1}=a_{i i \cdots i}\left[x_{i}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{i}\left(\mathbf{x}_{k, i-1}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left(\widehat{b}_{i}-\widehat{\mathcal{A}}_{i} \mathbf{x}_{k-1}^{m-1}\right)$,
$i=1,2, \cdots, n, k=1,2, \cdots$, where $\mathbf{x}_{k, i}=\left(x_{1}^{(k)}, \cdots x_{i}^{(k)}, x_{i+1}^{(k-1)}, \cdots x_{n}^{(k-1)}\right)^{T}$,
$\mathbf{x}_{k, 0}=\mathbf{x}_{k-1}, \mathbf{x}_{k, n}=\mathbf{x}_{k}$;
(4) If $\left\|\mathcal{A} \mathbf{x}_{k}^{m-1}-b\right\|_{2}<\varepsilon$, break and output $\mathbf{x}_{k}$;
(5) $k=k+1$ and back to step (2). Before giving the convergence theorem, we give following two lemmas.

Lemma 3.1. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. If $0<r \leq w \leq 1(w \neq 0)$, under the conditions of lemma 2.1, the vector sequence $\left\{\mathbf{x}_{k}\right\}$ generated by the algorithm 3.1 is increasing for any positive initial vector $\mathbf{x}_{0}>\mathbf{0}$ with $\mathbf{0}<\mathcal{A} \mathbf{x}_{0}^{m-1} \leq \mathbf{b}$.

Proof. By lemma 2.1, $\widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor. Assuming that $\mathbf{x}_{k-1}>\mathbf{0}$ and $\widehat{\mathcal{A}} \mathbf{x}_{k-1}^{m-1} \leq \widehat{\mathbf{b}}$, when $i=1$, it is obvious that

$$
\begin{aligned}
& a_{11 \cdots 1}\left[x_{1}^{(k)}\right]^{m-1} \\
= & a_{11 \cdots 1}\left[x_{1}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{1}\left(\mathbf{x}_{k, 0}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left(\widehat{b}_{1}-\widehat{\mathcal{A}}_{1} \mathbf{x}_{k-1}^{m-1}\right) \\
= & a_{11 \cdots 1}\left[x_{1}^{(k-1)}\right]^{m-1}+w\left(\widehat{b}_{1}-\widehat{\mathcal{A}}_{1} \mathbf{x}_{k-1}^{m-1}\right) \\
\geq & a_{i i \cdots i}\left[x_{1}^{(k-1)}\right]^{m-1} .
\end{aligned}
$$

When $i=j$, assuming that $\left[x_{t}^{(k-1)}\right]^{m-1} \leq\left[x_{t}^{(k)}\right]^{m-1}, t=1,2, \cdots, j$. When $i=j+1$, it is obvious that $\mathbf{x}_{k, j} \geq \mathbf{x}_{k, 0}=\mathbf{x}_{k-1}$. We have

$$
\begin{aligned}
& a_{j+1, j+1, \cdots, j+1}\left[x_{j+1}^{(k)}\right]^{m-1} \\
= & a_{j+1, j+1, \cdots, j+1}\left[x_{j+1}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{j+1}\left(\mathbf{x}_{k, j}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left[b_{j+1}-\mathcal{A}_{j+1} \mathbf{x}_{k-1}^{m-1}\right] \\
\geq & a_{j+1, j+1, \cdots, j+1}\left[x_{j+1}^{(k-1)}\right]^{m-1}+w\left[b_{j+1}-\mathcal{A}_{j+1} \mathbf{x}_{k-1}^{m-1}\right] \\
\geq & a_{j+1, j+1, \cdots, j+1}\left[x_{j+1}^{(k-1)}\right]^{m-1} .
\end{aligned}
$$

Therefore, $\left[x_{i}^{(k-1)}\right]^{m-1} \leq\left[x_{i}^{(k)}\right]^{m-1}$ for $i=1,2, \cdots, n$, we obtain $\mathbf{x}_{k} \geq \mathbf{x}_{k-1}$. Noticing that $0<r \leq w \leq 1$, then

$$
\begin{aligned}
& a_{i i \cdots i}\left[x_{i}^{(k)}\right]^{m-1} \\
= & a_{i i \cdots i}\left[x_{i}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{i}\left(\mathbf{x}_{k, i-1}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left[\widehat{b}_{i}-\widehat{\mathcal{A}}_{i} \mathbf{x}_{k-1}^{m-1}\right] \\
\leq & a_{i i \cdots i}\left[x_{i}^{(k-1)}\right]^{m-1}+\overline{\mathcal{L}}_{i}\left(\mathbf{x}_{k, i-1}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+\left[\widehat{b}_{i}-\widehat{\mathcal{A}}_{i} \mathbf{x}_{k-1}^{m-1}\right] \\
= & a_{i i \cdots i}\left[x_{i}^{(k-1)}\right]^{m-1}+\overline{\mathcal{L}}_{i} \mathbf{x}_{k, i-1}^{m-1}+\widehat{b}_{i}-\overline{\mathcal{U}}_{i} \mathbf{x}_{k-1}^{m-1} \\
\leq & a_{i i \cdots i}\left[x_{i}^{(k)}\right]^{m-1}+\overline{\mathcal{L}}_{i} \mathbf{x}_{k}^{m-1}+\widehat{b}_{i}-\overline{\mathcal{U}}_{i} \mathbf{x}_{k}^{m-1} \\
\leq & a_{i i \cdots i}\left[x_{i}^{(k)}\right]^{m-1}+\left[\widehat{b}_{i}-\widehat{\mathcal{A}}_{i} \mathbf{x}_{k}^{m-1}\right] .
\end{aligned}
$$

we get $\hat{\mathcal{A}} \mathbf{x}_{k}^{m-1} \leq \mathbf{b}$. Since $P>\mathbf{0}$ and $\mathbf{0}<\mathcal{A} \mathbf{x}_{0}^{m-1} \leq \mathbf{b}$, it is easy to know $\widehat{\mathcal{A}} \mathbf{x}_{0}^{m-1} \leq \widehat{\mathbf{b}}$. By the mathematical induction, we have $\mathbf{x}_{k}^{[m-1]} \geq \mathbf{x}_{k-1}^{[m-1]}$ $(k=1,2, \cdots)$. The vector sequence $\left\{\mathbf{x}_{k}\right\}$ is increasing.

Lemma 3.2. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. If $0 \leq r \leq w \leq 1(w \neq 0)$, under the conditions of lemma 2.1, the vector sequence $\left\{\mathbf{x}_{k}\right\}$ generated by the algorithm 3.1 is bounded above for any positive initial vector $\mathbf{x}_{0}>\mathbf{0}$ with $\mathbf{0}<\mathcal{A} \mathbf{x}_{0}^{m-1} \leq \mathbf{b}$.

Proof. By lemma 2.1, $\widehat{\mathcal{A}}$ is a strong $\mathcal{M}$-tensor. Let $\widehat{\mathcal{A}}=\widehat{\mathcal{D}}-\widehat{\mathcal{B}}$, where $\widehat{\mathcal{B}}$ is a nonnegative tensor and $\widehat{\mathcal{D}}=\widehat{D} \mathcal{I}_{m}$. Since $0<\mathcal{A} \mathbf{x}_{0}^{m-1} \leq b$, it is easy to know $\mathbf{0} \leq \widehat{\mathcal{A}} \mathbf{x}_{0}^{m-1} \leq \widehat{\mathbf{b}}$. Let

$$
\left\{\begin{array}{l}
\widehat{\mathcal{B}} \mathbf{x}_{0}^{m-1} \leq \alpha \widehat{\mathcal{D}} \mathbf{x}_{0}^{m-1}, \quad 0<\alpha<1 \\
\widehat{\mathbf{b}} \leq \beta \widehat{\mathcal{D}} \mathbf{x}_{0}^{m-1}
\end{array}\right.
$$

It is easy to know $\widehat{\mathcal{B}} \mathbf{x}_{0}^{m-1}+\mathbf{b} \leq(\alpha+\beta) \widehat{\mathcal{D}} \mathbf{x}_{0}^{m-1}$. By lemma 2.4, we can get $\mathbf{x}_{k}>\mathbf{0}$ and $\widehat{\mathcal{A}} \mathbf{x}_{k}^{m-1} \leq \widehat{\mathbf{b}}$ for $k=1,2, \cdots$. Therefore,

$$
a_{i i \cdots i}\left[x_{i}^{(k)}\right]^{m-1}=\widehat{\mathcal{D}}_{i} \mathbf{x}_{k}^{m-1} \leq \widehat{\mathcal{B}}_{i} \mathbf{x}_{k}^{m-1}+b_{i}
$$

We assume that $\mathbf{x}_{k-1}^{[m-1]} \leq\left(\alpha^{(k-1) n}+\alpha^{(k-1) n-1} \beta+\cdots+\alpha \beta+\beta\right) \mathbf{x}_{0}^{[m-1]}, k>1$. For the $k$-step iteration: when $i=1$, we have

$$
\begin{aligned}
& a_{11 \cdots 1}\left[x_{1}^{(k)}\right]^{m-1} \\
= & a_{11 \cdots 1}\left[x_{1}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{1}\left(\mathbf{x}_{k, 0}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left(\widehat{b}_{1}-\widehat{\mathcal{A}}_{1} \mathbf{x}_{k-1}^{m-1}\right) \\
\leq & \widehat{\mathcal{D}}_{1} \mathbf{x}_{k-1}^{m-1}+\left(\widehat{b}_{1}-\widehat{\mathcal{A}}_{1} \mathbf{x}_{k-1}^{m-1}\right) \leq \widehat{\mathcal{B}}_{1} \mathbf{x}_{k-1}^{m-1}+b_{1} \\
\leq & \left(\alpha^{(k-1) n}+\alpha^{(k-1) n-1} \beta+\cdots+\alpha \beta+\beta\right) \widehat{\mathcal{B}}_{1} \mathbf{x}_{0}^{m-1}+\widehat{b}_{1} \\
\leq & \left(\alpha^{(k-1) n+1}+\alpha^{(k-1) n} \beta+\cdots+\alpha \beta+\beta\right) a_{11 \cdots 1}\left[x_{1}^{(0)}\right]^{m-1} .
\end{aligned}
$$

we obtain $x_{1}^{(k)} \leq \sqrt[m-1]{\left(\alpha^{(k-1) n+1}+\alpha^{(k-1) n} \beta+\cdots+\alpha \beta+\beta\right)} x_{1}^{(0)}$.
When $i=j$, assuming that $x_{t}^{(k)} \leq \sqrt[m-1]{\left(\alpha^{(k-1) n+t}+\alpha^{(k-1) n+t-1} \beta+\cdots+\beta\right)} x_{t}^{(0)}$, $t=1,2, \cdots, j$.
When $i=j+1$, we have

$$
\begin{aligned}
& a_{j+1, \cdots, j+1}\left[x_{j+1}^{(k)}\right]^{m-1} \\
= & a_{j+1, \cdots, j+1}\left[x_{j+1}^{(k-1)}\right]^{m-1}+r \overline{\mathcal{L}}_{j+1}\left(\mathbf{x}_{k, j}^{m-1}-\mathbf{x}_{k-1}^{m-1}\right)+w\left[\widehat{b}_{j+1}-\widehat{\mathcal{A}}_{j+1} \mathbf{x}_{k-1}^{m-1}\right] \\
\leq & \widehat{\mathcal{D}}_{j+1} \mathbf{x}_{k-1}^{m-1}+\left(\widehat{b}_{j+1}-\widehat{\mathcal{A}}_{1} \mathbf{x}_{k, j}^{m-1}\right)=\widehat{\mathcal{B}}_{j+1} \mathbf{x}_{k, j}^{m-1}+\widehat{b}_{j+1} \\
\leq & \left(\alpha^{(k-1) n}+\alpha^{(k-1) n-1} \beta+\cdots+\alpha \beta+\beta\right) \widehat{\mathcal{B}}_{j+1} \mathbf{x}_{0}^{m-1}+\widehat{b}_{j+1} \\
\leq & \left(\alpha^{(k-1) n+1}+\alpha^{(k-1) n} \beta+\cdots+\alpha \beta+\beta\right) a_{j+1, \cdots, j+1}\left[x_{j+1}^{(0)}\right]^{m-1} .
\end{aligned}
$$

We obtain $x_{j+1}^{(k)} \leq \sqrt[m-1]{\alpha^{(k-1) n+j+1}+\alpha^{(k-1) n+j} \beta+\cdots+\alpha \beta+\beta} x_{j+1}^{(0)}$.

According to the initial conditions, we can get

$$
\mathbf{x}_{1}^{[m-1]} \leq\left(\alpha^{n}+\alpha^{n-1} \beta+\cdots+\alpha \beta+\beta\right) \mathbf{x}_{0}^{[m-1]}
$$

by calculating. By the mathematical induction, we have

$$
\mathbf{x}_{k}^{[m-1]} \leq\left(\alpha^{k n}+\alpha^{k n-1} \beta+\cdots+\alpha \beta+\beta\right) \mathbf{x}_{0}^{[m-1]}
$$

for $k=1,2, \cdots$. Let $k \rightarrow \infty$, then $\lim _{k \rightarrow \infty}\left(\alpha^{k n}+\alpha^{k n-1} \beta+\cdots+\alpha \beta+\beta\right)=\frac{\beta}{1-\alpha}$. Therefore, the vector sequence $\left\{\mathbf{x}_{k}\right\}$ is bounded above.

According to the two lemmas, we have the following convergence theorem.
Theorem 3.3. Let $\mathcal{A}$ be a strong $\mathcal{M}$-tensor. If $0 \leq r \leq w \leq 1(w \neq 0)$, under the conditions of lemma 2.1, then the vector sequence $\left\{\mathbf{x}_{k}\right\}$ generated by the algorithm 3.1 converges to the only positive limit for any positive initial vector $\mathbf{x}_{0}>\mathbf{0}$ with $\mathbf{0}<\mathcal{A} \mathbf{x}_{0}^{m-1} \leq \mathbf{b}$.

Proof. By the lemma 3.1 and lemma 3.2, the vector sequence $\left\{\mathbf{x}_{k}\right\}$ is increasing and bounded above. Consequently, $\left\{\mathbf{x}_{k}\right\}$ converges to the only positive limit.

## 4. Numerical examples

All numerical examples will be done in MATLAB R2018b on a personal computer with $\operatorname{Intel}(\mathrm{R})$ Core (TM) i5-7300HQ CPU @2.50GHz and 8.00GB RAM. In the section, "IT" and "CPU" denote the number of iteration steps and the CPU time, respectively. In the numerical examples, we set the maximum number of iterative steps to 1000 and the precision to $10-11$. $w$ is from 0.1 to 2 and the interval is 0.1 , then $(w, r)_{\text {opt }}$ denotes the optimal parameters of AOR and modified AOR method. Considering the following preconditioner:

$$
P=I+U_{\beta}=\left[\begin{array}{ccccc}
1 & -\beta_{1} a_{12 \cdots 2} & -\beta_{1} a_{13 \cdots 3} & \cdots & -\beta_{1} a_{1 n \cdots n} \\
0 & 1 & -\beta_{2} a_{23 \cdots 3} & \cdots & -\beta_{2} a_{2 n \cdots n} \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & 0 & \cdots & -\beta_{n-1} a_{n-1, n \cdots n} \\
0 & 0 & 0 & \cdots & 1
\end{array}\right]
$$

Just for convenience, let $\beta_{1}=\beta_{2}=\cdots=\beta_{n-1}=\beta$.
Example 1. Let $\mathcal{B} \in \mathbb{R}^{[3, n]}$ be a nonnegative tensor with $b_{i_{1} i_{2} i_{3}}=\left|\sin \left(i_{1}+i_{2}+i_{3}\right)\right|$. By [7], $\mathcal{A}=n^{2} \mathcal{I}_{m}-\mathcal{B}$ is a strong $\mathcal{M}$-tensor.

Let $\mathbf{b}=\mathbf{1}$ and initial vector $\mathbf{x}_{0}=\mathbf{0}$. We take the parameters $\beta$ from 0 to 4 and the interval is 0.3 . The numerical results are shown in table 1 .

Table 1. Numerical results of Example 1

| PAOR |  |  |  |  | modified PAOR |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: |
| $\beta$ | IT | CPU | $(w, r)_{\text {opt }}$ | IT | CPU | $(w, r)_{\text {opt }}$ |  |
| 0 | 27 | $2.055 \mathrm{e}-04$ | $(1.4,1.4)$ | 15 | $1.379 \mathrm{e}-04$ | $(1.5,1.4)$ |  |
| 0.3 | 26 | $1.422 \mathrm{e}-04$ | $(1.4,1.3)$ | 15 | $1.253 \mathrm{e}-04$ | $(1.5,1.4)$ |  |
| 0.6 | 23 | $1.226 \mathrm{e}-04$ | $(1.5,1.5)$ | 15 | $1.248 \mathrm{e}-04$ | $(1.5,1.4)$ |  |
| 0.9 | 23 | $1.153 \mathrm{e}-04$ | $(1.4,1.4)$ | 15 | $1.187 \mathrm{e}-04$ | $(1.5,1.4)$ |  |
| 1.2 | 22 | $1.113 \mathrm{e}-04$ | $(1.4,1.3)$ | 14 | $9.230 \mathrm{e}-05$ | $(1.5,1.4)$ |  |
| 1.5 | 21 | $1.132 \mathrm{e}-04$ | $(1.4,1.3)$ | 14 | $9.430 \mathrm{e}-05$ | $(1.5,1.4)$ |  |
| 1.8 | 20 | $1.131 \mathrm{e}-04$ | $(1.4,1.2)$ | 13 | $8.840 \mathrm{e}-05$ | $(1.5,1.4)$ |  |
| 2.1 | 18 | $9.450 \mathrm{e}-05$ | $(1.4,1.4)$ | 13 | $8.820 \mathrm{e}-05$ | $(1.5,1.4)$ |  |
| 2.4 | 17 | $9.140 \mathrm{e}-05$ | $(1.4,1.3)$ | 12 | $7.980 \mathrm{e}-05$ | $(1.5,1.4)$ |  |
| 2.7 | 16 | $8.090 \mathrm{e}-05$ | $(1.4,1.4)$ | 12 | $8.020 \mathrm{e}-05$ | $(1.3,1.2)$ |  |
| 3.0 | 16 | $7.540 \mathrm{e}-05$ | $(1.4,1.4)$ | 12 | $7.910 \mathrm{e}-05$ | $(1.3,1.2)$ |  |
| 3.3 | 16 | $8.990 \mathrm{e}-05$ | $(1.3,1.3)$ | 11 | $7.130 \mathrm{e}-05$ | $(1.3,1.2)$ |  |
| 3.6 | 15 | $1.343 \mathrm{e}-04$ | $(1.3,1.3)$ | 11 | $7.360 \mathrm{e}-05$ | $(1.3,1.2)$ |  |
| 3.9 | 16 | $8.540 \mathrm{e}-05$ | $(1.3,1.2)$ | 11 | $7.380 \mathrm{e}-05$ | $(1.3,1.2)$ |  |

Example 2 [24]. Let $\mathcal{A} \in \mathbb{R}^{[3, n]}$ and $\mathbf{b} \in \mathbb{R}^{n}$ with:

$$
\left\{\begin{array}{lr}
a_{111}=(2+n) / 2, & \\
a_{n n n}=1, & i=2,3, \cdots, n-1, \\
a_{i i i}=2, & i=2,3, \cdots, n-1, \\
a_{1 i i}=-1 / 2, & i=2,3, \cdots, n-1, \\
a_{i i i-1}=-1 / 2, & i=2,3, \cdots, n-1, \\
a_{i i-1 i-1}=-1 / 2, & i=2,3, \cdots, n-1 \\
a_{i i+1 i+1}=-1 / 2, &
\end{array}\right.
$$

and

$$
\left\{\begin{array}{l}
b_{1}=c_{0}^{2} \\
b_{i}=a /(n-1)^{2}, \quad i=2,3, \cdots, n-1 \\
b_{n}=c_{1}^{2}
\end{array}\right.
$$

Let initial vector $\mathbf{x}_{0}=(1,1, \cdots, 1)^{T}, c_{0}=1 / 2, c_{1}=1 / 3$ and $a=2$. When $\mathrm{n}=10$, we take the parameters $\beta$ from 0 to 2 and the interval is 0.2 . The numerical results are shown in table 2 .

It is well known that we can get the Jacobi, the Gauss-Seidel and the successive overrelaxation (SOR) iteration methods by choosing certain values. The results of two numerical examples show that the AOR method is more effective than the GS and SOR method when we take the optimal parameter $(w, r)$. We compare AOR with modified AOR iteration method in example 1 and example 2. It can be seen from Table 1 and table 2 that modified AOR iteration method
requires less iterative steps and CPU time.

TABLE 2. Numerical results of Example 2

|  | PAOR |  |  |  | modified PAOR |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: |
| $\beta$ | IT | CPU | $(w, r)_{\text {opt }}$ | IT | CPU | $(w, r)_{\text {opt }}$ |  |
| 0 | 45 | 0.0034 | $(1.2,1.2)$ | 24 | 0.0011 | $(1.3,1.2)$ |  |
| 0.2 | 41 | 0.0404 | $(1.4,1.4)$ | 22 | 0.0008 | $(1.3,1.2)$ |  |
| 0.4 | 37 | 0.0365 | $(1.4,1.4)$ | 21 | 0.0007 | $(1.3,1.0)$ |  |
| 0.6 | 32 | 0.0303 | $(1.4,1.4)$ | 21 | 0.0007 | $(1.3,1.0)$ |  |
| 0.8 | 29 | 0.0014 | $(1.4,1.1)$ | 21 | 0.0007 | $(1.3,0.8)$ |  |
| 1.0 | 28 | 0.0279 | $(1.4,1.4)$ | 21 | 0.0007 | $(1.1,1.0)$ |  |
| 1.2 | 27 | 0.0009 | $(1.2,1.2)$ | 20 | 0.0008 | $(1.1,1.0)$ |  |
| 1.4 | 24 | 0.0009 | $(1.2,1.2)$ | 19 | 0.0007 | $(1.1,1.0)$ |  |
| 1.6 | 22 | 0.0026 | $(1.2,1.2)$ | 21 | 0.0007 | $(1.3,1.0)$ |  |
| 1.8 | 25 | 0.0232 | $(1.2,1.2)$ | 33 | 0.0012 | $(1.3,1.2)$ |  |
| 2.0 | 28 | 0.0007 | $(1.0,1.0)$ | 69 | 0.0024 | $(1.5,1.4)$ |  |

## 5. Conclusions

In this paper, we propose two preconditioned AOR iteration methods to solve multi-linear systems with $\mathcal{M}$-tensor. We prove that the two proposed methods are convergent when the preconditioner satisfies certain conditions. Moreover, the comparison theorem shows that the AOR iteration method with preconditioner converges faster than the method without preconditioner. A new preconditioner is given in the numerical examples. The results of the numerical examples illustrate the effectiveness of these methods. In the future, we will consider how to choose the best preconditioner.

## References

1. L.Q. Qi and Z.Y. Luo, Tensor Analysis: Spectral Theory and Special Tensors, M. Philadelphia: SIAM, 2017.
2. W. Ding and Y. Wei, Theory and Computation of Tensors, M. Math. London: Academic Press, 2016.
3. W.Y. Ding and Y.M. Wei, Solving multi-linear systems with M-tensors, J. Sci. Comput. 68 (2016), 689-715.
4. Z. Luo, L. Qi and N. Xiu, The sparsest solutions to Z-tensor complementarity problems, J. Optimi. Lett. 11 (2017), 471-482.
5. X.-T. Li and K.N. Michael, Solving sparse non-negative tensor equations: algorithms and applications, J. Front. Math. 10 (2015), 649-680.
6. S. Du, L. Zhang, C. Chen and L. Qi, Tensor absolute value equations, Sci. China Math. 61 (2018), 1695-1710.
7. Z.J. Xie, X.Q. Jin and Y.M. Wei, Tensor methods for solving symmetric-tensor systems, J. Sci. Comput. 74 (2018), 412-425.
8. M.L. Liang, B. Zheng and R.J. Zhao, Alternating iterative methods for solving tensor equations with applications, J. Numerical Algorithms 80 (2019), 1437-1465.
9. H. He, C. Ling and L. Qi, G. Zhou, A globally and quadratically convergent algorithm for solving ultilinear systems with M-tensors, J. J. Scient. Comput. 76 (2018), 1718-1741.
10. L. Han, A homotopy method for solving multilinear systems with m-tensors, J. Appl. Math. Lett. 69 (2017), 49-54.
11. X.Z. Wang, M.L. Che and Y.M. Wei, Neural networks based approach solving multi-linear systems with M-tensors, J. Neurocomputing on Science Direct 351 (2019), 33-42.
12. D.D. Liu, W. Li, Seak-WengVonga, The tensor splitting with application to solve multilinear systems, J. of Computational and Applied Mathematics 330 (2018), 75-94.
13. D.D. Liu, W. Li, Seak-WengVonga, Comparison results for splitting iterations for solving multi-linear systems, J. Applied Numerical Mathematics 134 (2018), 105-121.
14. L.B. Cui, M.H. Li and Y.S. Song, Preconditioned tensor splitting iterations method for solving multi-linear systems, J. Applied Mathematics Letters 96 (2019), 89-94.
15. W.Y. Ding, L.Q. Qi and Y.M. Wei, M-tensors and nonsingular M-tensors, J. Linear Algebra and its Applications 439 (2013), 3264-3278.
16. Kelly J. Pearson, Essentially Positive Tensors, International Journal of Algebra 4 (2010), 421-427.
17. Liqun Qi, Eigenvalues of a real supersymmetric tensor, Journal of Symbolic Computation 40 (2005), 1302-1324.
18. A. Hadjidimos, Accelerated Overrelaxation Method, J. Mathematics of Computation 32 (1978), 149-157.
19. H. Niki, K. Harada and M. Morimoto, M. Sakakihara, The survey of preconditioners used for accelerating the rate of convergence in the Gauss-Seidel method, J. of Computational and Applied Mathematics 164 (2004), 587-600.
20. H. Kotakemori, H. Niki and N. Okamotob, Accelerated iterative method for Z-matrices, Journal of Computational and Applied Mathematics 75 (1996), 87-97.
21. K.C. Chang, K. Pearson and T. Zhang, Perron-Frobenius Theorem for nonnegative tensors, J. Communications in Mathematical Sciences 6 (2008), 507-520.
22. G.H. Golub , C.F. Van Loan , Matrix Computations, M. Johns Hopkins University Press, Baltimore, 2013.
23. Y. Zhang, Q. Liu and Z. Chen, Preconditioned Jacobi type method for solving multi-linear systems with M-tensors, J. Applied Mathematics Letters 104 (2020), 104:106287.
24. L.B. Cui, X.Q. Zhang and S.L. Wu, A new preconditioner of the tensor splitting iterative method for solving multi-linear systems with $M$-tensors, J. Computational and Applied Mathematics 39 (2020), 1-16.
25. D.D. Liu, W. Li and Seak-WengVonga, A new preconditioned SOR method for solving multi-linear systems with an M-tensors, J. Calcolo 57 (2020), Doi:10.1007/s10092-020-00364-8.

Meng Qi received M.Sc. from Northeastern University. His research interests include iterative algorithm and numerical optimization.

Department of Mathematics, College of Sciences, Northeastern University, Shenyang 110819, P.R. China.
e-mail: 769100925@qq.com
XinHui Shao received M.Sc. from Northeast Normal University, and Ph.D. from Northeastern University. Her research interests are computational mathematics, iterative method and Applied statistical research.
Department of Mathematics, College of Sciences, College of Sciences, Northeastern University, Shenyang 110819, P.R. China.


[^0]:    Received January 26, 2021. Revised April 16, 2021. Accepted April 22, 2021. * Corresponding author
    ${ }^{\dagger}$ This work was supported by by Central University Basic Scientific Research Business Expenses Special Funds(N2005013).
    © 2021 KSCAM.

