Non-periodic Subway Scheduling that Minimizes Operational Cost and Passenger Waiting Time

YunWoo Hong*, Yerim Chung**, YunHong Min***

Abstract

Subway metro scheduling is one of the most important problems impacting passenger convenience today. To operate efficiently, the Seoul metro uses regular, periodic schedules for all lanes, both north and southbound. However, many past studies suggest that non-periodic scheduling would better optimize costs. Since the Seoul metro is continuously facing a deficit, adopting a non-periodic schedule may be necessary. Two objectives are presented: the first, to minimize the average passengers’ waiting time, and the second, to minimize total costs, the sum of the passenger waiting time, and the operational costs. In this paper, we use passenger smart card data and a precise estimation of transfer times. To find the optimal time-table, a genetic algorithm is used to find the best solution for both objectives. Using Python 3.5 for the analysis, for the first objective, we are able to reduce the average waiting time, even when there are fewer trains. For the second objective, we are able to save about 4.5 thousand USD with six fewer trains.

Keyword: Transfer Passenger, Subway Scheduling, Timetabling, Optimization, Genetic Algorithm, Metro Operating Cost

I. Introduction

In Seoul, Korea, more than 170 thousand people, on average, are using the Seoul metro line number 8 each weekday[1][2]. With such a large demand, the metro is facing various operating problems. One of the biggest is its endless deficit. According to Money News[3], the Seoul metro line lost more than about 370 million USD in 2014 and this has continued to increase to date. Seoul metro has pointed to several reasons for this deficit: the first is the low price of the tickets, the second is the large number of people who ride for free (legally and illegally), and the third is the decrease in its advertising earnings. In order to overcome this situation, the line has raised ticket prices continuously. The ticket fare has increased 150% over its 2000 price. Moreover, the line has cracked down on illegal train riders. Currently, Seoul metro is considering shortening the subway schedule. Ironically, all these actions taken to save costs have incurred other costs; not only monetary costs but also non-monetary costs such as passenger inconvenience. Thus, true cost saving should be considered from better schedule efficiency.

The Seoul metro schedule is fairly regular and periodic. On line 8, most of the headways are fixed at every five minutes in rush hour, and every eight minutes
otherwise. This periodic schedule is easy and convenient to operate. However, it may increase passenger waiting time and decrease resource efficiency. According to Barrena et al. [4], a regular schedule does not always minimize passenger waiting time. The Seoul Statistics reveal that line 8 has the lowest demand level and among the lowest congestion, as shown in figure 1. Before noon, demand for line 8 reaches only 71% of its capacity, while Seoul metro has stated that it can handle up to 150% of its capacity. This offers evidence that its operational efficiency is low. Thus, Seoul metro needs to adjust its operating system, the timetable, and number of subways. To increase efficiency, the subway schedule should better meet passenger demand. Thus, a non-periodic schedule based on passenger demand appears to be necessary.

III. Literature Review

Public transportation scheduling as a problem was initiated in 1971 by Gordon Newell[6]. He tried to find the optimal schedule to minimize passenger waiting time. After his mathematical modeling, scholars tried to solve related problems. Many tried to minimize costs or passenger waiting time. Guan et al. [7] and Assis and Milani[8] used exact algorithms. On the other hand, Li [9] used simulation methods to solve a similar problem. Meta heuristic methods have also been used ([10][11][12]). Among them, Hu et al. [10] and Niu and Zhou [11] used genetic algorithms to find the optimal schedule to minimize passenger waiting time. Other than general scheduling problems, many scholars have tried to solve timetable synchronization problems (TTSPs) to minimize transfer passenger waiting time. Many methodologies are also used in TTSPs. Wong and Leung[13] used mixed integer programming; Feng et al.[14] used a simulation; and Poorjafari et al.[15] used genetic algorithms and simulated annealing. In passenger flow problems, Hong et al. [16] and Park and Lee[17] used smart card data for their predictions. Since smart card data mainly contain departure station (O), entry time at gate (OT), arrival station (D), and exit time at gate (DT), they could estimate passenger flow precisely.

As all of the studies are related, these factors should be considered at the same time. From this point of view, these past studies contain some limitations. In the regular scheduling problems, transfer passengers were not considered and neither were non-transfer passengers in the TTSP. On the Seoul metro, more than half of the passengers are transferring passengers. Especially on line 8, where more than 70% are transferring passengers. Thus, both non-transfer and transfer passengers should be considered. In this study, we precisely estimate the flow of the transfer passengers to find precise metro schedules that minimize passenger average waiting time and total cost while considering both types of passengers. We use the Seoul metro line 8 as the subject of the experiment. However, we
only consider the northbound lane until 11:30 AM.

To consider both non-transfer and transfer passengers, we use smart card data to estimate their flow and real data to estimate the transfer time. Next, with a genetic algorithm, we try to identify the optimal solution for both objectives presented earlier.

IV. Method

1. Assumption

In this study, we assume that all passengers obey the first in first out (FIFO) principle and that they take the first possible train. Therefore, they take the train, within its capacity limit, in order. Those assumptions are necessary to solve the problem mathematically with smart card data. Without this assumption, the efficiency of smart card data decreases gradually, and it is nearly impossible to compute using a mathematical model. The last assumption is that all passengers choose the shortest paths for their trips. This assumption is needed to estimate the transfer passenger paths. Thus, this is a key assumption for data reformulation of transfer passenger data. Those three assumptions are widely used in past studies to solve various research questions related to transportation systems.

2. Data Reformulation

2.1 Determining the transfer station

As we explained, smart card data contain O, OT, D, and DT, which can also be called the O-D pair. Using O and OT, we can determine the number of passengers boarding and waiting, and D helps us determine the remaining capacity of the train for future boarding passengers. DT is not necessarily considered since it changes depending on the subway schedule. With this O-D pair, we can divide passengers into different types as follows:

- **Non-Transfer Passenger**: Passengers without transfer
- **Transfer Passenger 1-1**: Start the trip with line 8, transfer once
- **Transfer Passenger 1-2**: Start the trip with line 8, transfer twice
- **Transfer Passenger 2-1**: End the trip with line 8, transfer once
- **Transfer Passenger 2-2**: End the trip with line 8, transfer twice
- **Transfer Passenger 3**: Did not start and end the trip with line 8, transfer twice

In this study, we did not consider passengers with three or more transfers since their population is about 1.5% of the total trips on the Seoul metro (Shin et al.[18]). Since we are considering line 8, we need data that relate to line 8 only. However, for transfer passengers, their O or D, or both, will not be on line 8. Additionally, OT is the time when the passenger passes the gate and not the actual platform arrival time. Therefore, data reformulation is necessary. More than the O-D pair, we need smart card data containing the transfer line ID and its direction. With such data and our second assumption (passenger chooses the shortest path), we can determine the transfer station.

In order to update the OT for non-transfer passengers, we need to determine the walking speed of the passengers. According to [19], the average walking speed in a low population density situation is 1.31 m/sec with a standard deviation of 0.22, and in a high population density situation, the average speed is 1.17 m/sec with a standard deviation of 0.23. Since rush hour is more likely to have higher population density, we use 1.17 m/sec for passengers during rush hour and 1.31 m/sec at other times. Moreover, according to [20], the average speed for walking up stairs is 0.67 m/sec, while the average speed walking down is 0.77 m/sec. According to Seoul metro, the average distance between the gates to the platform is 100m and one stair level down. Thus, using such data, we estimate the non-transfer passenger platform arrival time.

Figure 2 presents the line 8 route map from the official

Fig. 2. Seoul Metro Line 8 Route Map
Seoul metro website. There are five transfer stations, Moran, Bokjeong, Garak Market, Jamsil, and Cheonho. Moran and Bokjeong offer transfers to the Bundang Line, Garak Market to Line 3, Jamsil to Line 2, and Cheonho to Line 5. According to these data, except for the Bundang line, the other transfer lines have exactly one transfer station. Thus, we can determine the transfer passenger’s transfer station (who does not use the Bundang Line) without any problem. For Bundang line users, there are two transfer stations, Moran and Bokjeong. According to Seoul metro, transferring at Bokjeong takes less time for passengers heading further than Bokjeong on line 8 (for the northbound lane) than transferring ahead at Moran. With the earlier assumption, we now can determine their transfer station.

After the transfer station is determined, the transfer passengers’ O and D now switches to stations on line 8. Then, for transfer passengers 1–1 and 1–2, we can add walking and going down stairs time to their OT. Next, non-transfer passenger data reformulation is completed. For the rest of the transfer passenger types, we need to determine new platform arrival times. As we explained, this is significant factor. Thus, estimating arrival time is an important issue in this study.

2.2. Passenger arrival time estimation

To estimate transfer time, we visited all of the transfer stations to measure various factors. The findings are shown in Figure 3.

Due to the existence of the passenger transfer and the nearest door for the transfer, we need to estimate passenger behavior. To determine the transfer method, we manually observed six different transfer spots to find the ratio of those choosing the stairs. For the more than 40 observations we collected, we tried to find the relation between the number of passengers and the choice of the stairs as a ratio by using a linear regression model. We used IBM SPSS Statistics version 23, and the results are given in Figure 4.

![Fig. 4. Linear Regression Results](image)

Since the Durbin-Watson value is close to 2, the model independence is approved. The R square value is 0.42, which indicates that the model explains 42% of the values, which is high enough to be used. Moreover, all of the values are less than 0.05, which indicates that the model is valid. Therefore, we can conclude that the relation between choosing the stairs and the number of passengers is as in equation (1) below.

\[
\text{Ratio of Stair Choice} = 10.967 + 0.231\times (\text{Number of Passengers})
\] (1)

To determine which door passengers use to exit to transfer, we constructed a survey. This factor is important because the walking distance will differ depending on which door the passenger uses. The survey included questions related to on-train or beforehand migration to faster exits after arrival using a five-point Likert scale. A total of 491 people participated. We sampled only those who use the subway more than four times a week, which represented a total of
325 people. The results are given in Figure 5.

Finally, we have to determine when the passengers arrive at the transfer station. For transfer passengers, their arrival time is dependent on their train arrival time. Not only to reduce computation time but also to build a framework for future studies, we assume all of the transfer lines have the same number of train arrivals in an hour. In addition, the number of trains in the hour is estimated by the average number of trains arriving at the transfer lines. Moreover, we assume that the distribution of transfer passenger flow is equivalent to the non-transfer passenger flow. With these findings and assumptions, we can estimate the transfer passengers' arrival times at the platforms. Now, all of the passengers have O-D pairs only related to the Seoul metro line 8.

3. Genetic Algorithm

A genetic algorithm is widely used in scheduling problems. The following related studies have used such algorithms previously.

<table>
<thead>
<tr>
<th>Author</th>
<th>Objective</th>
<th>Chromosome</th>
<th>Initial population</th>
<th>Reproduction</th>
<th>Crossover</th>
<th>Mutation</th>
<th>Selection</th>
</tr>
</thead>
<tbody>
<tr>
<td>Y. Li &amp; Y. Fang (2013)</td>
<td>Minimize Total Cost in Emergency Situation</td>
<td>Binary</td>
<td>Random</td>
<td>Single point crossover (rate=0.7)</td>
<td>Flip bit method (rate=0.01)</td>
<td>Roulette wheel</td>
<td></td>
</tr>
<tr>
<td>C. Rongwu et al. (2014)</td>
<td>Train Speed Optimization in Automatic System</td>
<td>m x n matrix (integer)</td>
<td>Random</td>
<td>Roulette wheel</td>
<td>Single point crossover (rate=0.8)</td>
<td>Random position changed to random feasible integer (rate=0.01)</td>
<td>Roulette wheel</td>
</tr>
<tr>
<td>P. Tormos et al. (2008)</td>
<td>Minimize Total Travel Time</td>
<td>array, unit-pair (t, s)</td>
<td>Regret-based random sampling</td>
<td>Single point crossover (rate=0.8)</td>
<td>re-positioning (rate=0.05)</td>
<td>2-tournament selection</td>
<td></td>
</tr>
<tr>
<td>W. Zhu &amp; H. Hu (2014)</td>
<td>Minimize Total Travel Time</td>
<td>Binary</td>
<td></td>
<td>Two point crossover</td>
<td>Flip bit method</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

To avoid a convergence of solutions, the solutions are selected without replacement. Moreover, to increase improvement efficiency, the solutions that are better than the original solutions are automatically selected.
first, we remove the lowest demand train; second, we remove a randomly selected train. Those options are also presented in [4]. Both are considered because this study focuses on two different objectives and we are willing to assess which method is more effective to solve the problems.

V. Results and Discussion

To determine the schedule efficiency, this study classifies passengers into two different types. Type I passengers are passengers who take the subway without facing any capacity issues and Type II passengers are those who do. In other words, Type II passengers will have waited for other trains due to oversaturated train conditions. We use Python 3.5 in Intel® Core™ i3-6100 CPU @ 3.70GHz with 8.00GB RAM.

1.1 Problem 1 and Removing the lowest demand level train

Compared to the original solution, the passenger average waiting time decreased by 10 seconds while one less train was used. Moreover, since an average of 11 more passengers rode each train with no Type II passengers, operational efficiency improved significantly. On the other hand, the maximum waiting time, the waiting time for passengers who had to wait the longest, increased by two minutes. This was expected since we adopted a non-periodic schedule without increasing the number of trains.

1.2 Problem 1 and Removing the randomly selected train

Although the number of trains used was the same, the average waiting time decreased by six seconds. There

<table>
<thead>
<tr>
<th></th>
<th>Average Waiting Time</th>
<th>Type I Average Waiting Time</th>
<th>Number of Type I</th>
<th>Number of Type II</th>
<th>Max Demand</th>
<th>Min Demand</th>
<th>Average Demand</th>
<th>Max Waiting Time</th>
<th>Average Waiting Time in Transfer Station</th>
<th>Average Waiting Time in Non-Transfer Station</th>
</tr>
</thead>
<tbody>
<tr>
<td>Original</td>
<td>182.53</td>
<td>182.53</td>
<td>24,547</td>
<td>0</td>
<td>51</td>
<td>1102</td>
<td>597.61</td>
<td>500</td>
<td>184.77</td>
<td>181.37</td>
</tr>
<tr>
<td>Improved</td>
<td>176.05</td>
<td>176.05</td>
<td>34,547</td>
<td>0</td>
<td>51</td>
<td>1210</td>
<td>598.64</td>
<td>719</td>
<td>169.81</td>
<td>174.66</td>
</tr>
</tbody>
</table>

Fig. 8. Result of Problem 1 and Removing the Lowest Demand Level Train

Fig. 9. Schedules at the Jamsil Station in Problem 1 and Removing the Lowest Demand Level Train

Fig. 10. Result of Problem 1 and Removing Randomly Selected Train
were no Type II passengers as well, which shows that the operating efficiency improved. Due to the characteristics of the non-periodic schedule, the maximum waiting time increased by 1.5 minutes.

For both these methods where trains were removed, the average waiting time decreased. As shown in figures 8 and 10, the waiting time at the transfer station decreased a larger amount compared to that at the non-transfer station. This may because of the characteristics of the genetic algorithm. The average

number of boarding passengers at the transfer stations was greater than at the non-transfer stations (an average of 540 passengers). Since the genetic algorithm seeks the solution that fits better, it chooses to decrease the larger wait time at the transfer station. The number of trains in the second method stayed the same, while it decreased in the first method. This may because of the objective. The objective of problem 1 was very focused on the average waiting time. Thus, the objective would not improve if we removed trains without considering passenger flow. Notably, the first method removing the train gave us a better result than the second method. With one less train used, the average waiting time was smaller in the first method. This may because of the limitation of randomness. Usually, the genetic algorithm provides better results with higher randomness. Thus, fewer trains give a wider range of scheduling because each train not only has running time but also has platform wait time and headway. However, this result may change if we run a genetic algorithm infinitely. In problem 1, removing the lowest demand train shows a better result because of our waiting time focused objective.

1.3 Problem 2 and Removing the lowest demand level train

With a different objective, minimizing total cost, the improved schedule saves two million won (1.8 thousand USD). The average passenger waiting time decreases by 1.5 seconds and the number of trains decreases by one. As maximum demand decreases and minimum demand increases, the schedule is well distributed to the passenger flow.

1.4 Problem 2 and Removing the Randomly Selected Train

With the method removing the selected train randomly, six fewer trains are used to operate the metro. This results in 5,337,596 won cost saving, which is about 4.8 thousand USD. Without any Type II passengers, average waiting time increases by 20 seconds.
A big difference between the two methods is the change in average waiting time and the number of trains. The first method decreases the average waiting time, while the second method increases it. However, as shown in the total cost change, randomly removing the train shows a better result with a fewer number of trains used. This may be because the number of trains has a bigger influence on the objective 2, minimizing total cost. Moreover, in the first train removal method, if the result does not improve after the removal, the genetic algorithm will not decrease the number of trains at all. Therefore, decreasing passenger waiting time is the only way to improve the solution. Decreasing waiting time, however, has a smaller impact than decreasing the number of trains. In the random removal method, it keeps trying to decrease the number of trains randomly. Therefore, for objective 2, the second method of removing trains shows a better result due to the higher chance of removing a train.

VI. Conclusion and Future Study

In all of the experiments, using a non-periodic schedule shows a much better result, not only in terms of passenger convenience but also in operating cost. Thus, we can conclude that flexible scheduling is more effective than standard scheduling. Moreover, there are no type II passengers in any solution. Thus, we can also conclude that subway line 8 currently being under utilized is an issue. In this research, we reflect not only on the transferring passengers but also on the non-transferring passengers. Moreover, this study suggests precise estimation of transfer time. Although those topics are not new, we combine related studies into one topic to find more precise solutions. Seoul metro periodically updates its schedule; in the case of line 8, the latest update was four or five years ago. In the next update, our findings can be applied so that the metro can consider both types of passengers as well as the adoption of a non-periodic time table. Moreover, Seoul metro faced a huge subway union strike in 2016. During that time, six or less trains were used before noon. Even with fewer trains and highly saturated conditions, Seoul metro continued to use regular scheduling. As this study proves, non-periodic schedule can give a better result and the metro should consider a more flexible scheduling in various situations.

This study has some limitations as well. We only consider trains traveling one way and with a cutoff time due to computation time. Thus, we should improve this to include both ways and a full day schedule so that methods and results can be closer to the real world. For future study, real time scheduling dependent on dynamic capacity and passengers incoming will be one possible field of research. Since real time scheduling can reflect real time or stochastic passenger data, it can be flexible in responding to any kind of situation such as an emergency or an unexpected oversaturation of the number of passengers. Additionally, updating the schedule of the first and last train by considering passenger demand could be another field of research. This study should also consider the cost for people who are unable to ride the train due to the shift in schedule.

REFERENCES

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Yunwoo Hong received the M.S. degree in Business Administration from Yonsei University in 2017. He is interested in transportation scheduling problems.

Yerim Chung received the B.S. degree in Business Administration from Yonsei University, Korea, in 2000. She received the M.S. and Ph.D. degree in Applied Mathematics and Computer Science from Paris 1 University, France, in 2004 and 2010, respectively. Dr. Chung joined the faculty of Business School at Yonsei University, Seoul, Korea, in 2011. She is interested in inverse optimization, transportation science and network effects.

YunHong Min received the B.S. degree in Industrial & Management Engineering from POSTECH, and Ph.D. degree in Industrial Engineering from Seoul University in 2006 and 2012, respectively. Dr. Min joined the faculty of Graduate School of Logistics, Incheon National University in 2018. He is interested in mathematical programming and machine learning.