

RSA – QoS: A Resource Loss Aware Scheduling Algorithm for Enhancing the Quality of Service in Mobile Networks

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Abstract

Adaptive Multi-Rate Codec is one of the codecs which is used for making voice calls. It helps to connect people who are scattered in various geographical areas. It adjusts its bit-rate according to the user's channel conditions. It plays a vital role in providing an improved speech quality of voice connection in Long Term Evolution (LTE). There are some constraints which need to be addressed in providing this service profitably. Quality of Service (QoS) is the dominant mechanism which determines the quality of the speech in communication. On several occasions, number of users are trying to access the same channel simultaneously by standing in a particular region for a longer period of time. It refers to Multi-user channel sharing problem which leads to resource loss very often. The main aim of this paper is to develop a novel RSA – QoS scheduling algorithm for reducing the Resource Loss Ratio. Eventually, it increases the throughput. The simulation result shows that the RSA – QoS increases the number of users for accessing the resources better than the existing algorithms in terms of resource loss and throughput. Ultimately, it enhances the QoS in Mobile Networks.

Keywords: Adaptive Multi-Rate, Voice over Internet Protocol, Resource Block, Throughput, Resource Loss Ratio and Long Term Evolution.

1. Introduction

LTE is a technology of Fourth Generation [1] network which is an enhanced version of the third generation network. It was launched in the year of 2004 by 3rd Generation Partnership Project (3GPP) [2]. In High Mobility, it provides 100 Mbps and 50 Mbps for downloading and uploading respectively. In Low Mobility, it offers 1 Gbps data transfer speed along with quicker smooth handoff. LTE has 20 MHz wider in channel bandwidth for accessing purpose including control signals. It supports a vast number of multimedia applications [3] like Skype, GTalk, IMO, Online Gaming, etc. The quality of voice has been expanded and enriched in this Fourth Generation network. Circuit Switching has been replaced by Packet switching in order to reduce the connection establishment delay of 4G network. 3GPP has introduced the concept of Voice over Long Term Evolution (VoLTE) [4] to facilitate voice service in LTE. VoIP has become more prominent since it provides voice call in a cheaper cost than VoLTE. It is very sensitive to QoS by having the Quantitative parameters like Delay, Jitter, Packet Loss, Packet Delivery Ratio, Throughput, etc., eventhough, it provides lower cost. eNodeB acts as a scheduler for resource allocation by means of Medium Access Control (MAC) [5]. It plays a reasonable role in assigning resources to the users and terminates the users connection as soon as the service is over. LTE works by invoking the Evolved Universal Terrestrial Radio Access Network (E-UTRAN) architecture for providing the seamless communication. It adapts all types of networks like GSM, CDMA, UMTS, WLAN, etc. Generally, the data transmission is carried out in the form of frames in this network. Fig. 1 shows the structure of the LTE Frame [6] in detail.

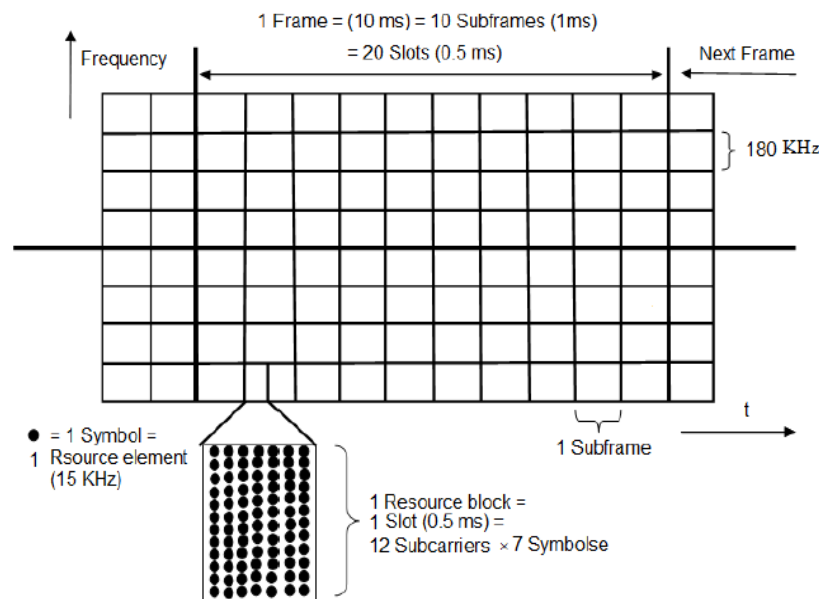


Fig. 1. LTE Frame Structure [6]

One Frame is composed of ten sub frames. Each sub frame lasts for one Transmission Time Interval (TTI). One sub frame is made up of two individual slots where each has 0.5ms. Based on the channel bandwidth, the number of Resource Blocks differs. One Resource Block is the combination of 12 subcarriers and 7 symbols in normal prefix mode and 6 Symbols in an extended mode. So, totally 1 RB contains 168 (12 X 7 X 2) Resource Elements (RE) or extended contains 144 REs (12 X 6 X 2). Each RE can carry 2, 4 and 6 bits in QPSK, 16 QAM

and 64 QAM modulations respectively. Mobile users compute and send their CQI values as a feedback to the eNodeB by using the reference symbols. Thereafter, the eNodeB decides the number of resources that the user can get access for transmission through downlink channel. For downlink scheduling, LTE uses Orthogonal Frequency Division Multiple Access (OFDMA) scheme and Adaptive Modulation and Coding (AMC) adjusts the coding rate for determining the modulation. Fig. 2 explains the flow of LTE downlink scheduling model.

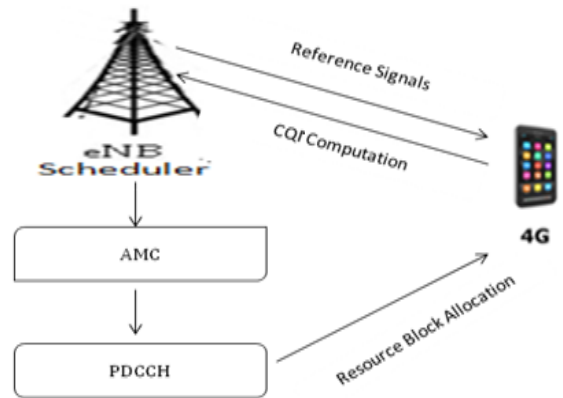


Fig. 2. Downlink Scheduling Model

This paper is structured as follows. Section 2 explains the Review of literature to prepare this article and Section 3 elucidates the proposed RSA – QoS scheduling algorithm to reduce the resource loss in detail. Section 4 discusses the Scenarios and Discussions of the proposed algorithm, Section 5 provides the Mathematical Model and Notations, and Section 6 shows the Performance Evaluation. Section 7 deals with the Conclusion and Finally, References is placed in Section 8.

2. Related Work

Initial research on scheduling, the channel bandwidth (Time/Frequency based resources) was not used in effective manner. This section shows the inconsistency of the scheduler pertaining to reduction of the resource loss ratio in increasing the system throughput. Capozzi et al. [7] investigated Heterogeneous network, the preliminary steps to develop a scheduling algorithm, its service architecture and the constraints needed to build a scheduling model. It enlists some research problems in scheduling with possible solutions. Round Robin (RR) [8] algorithm works without considering the channel conditions and assigns the resources to the users by First Come First Served basis. It lacks in two things. First, the QoS requirements of an application and second, their demand in resource allocation. It gives fine results in fairness because all the users get an equal opportunity in accessing the resources.

Best Channel Quality Indicator (BCQI) [9] algorithm keeps on assigning the resources to the users only by looking at the channel conditions of highest quality users. It fails to give the resources to the cell edge users. Therefore, fair allocation in BCQI is very poor when comparing with other scheduling algorithms. The resource allocation decision metric of i^{th} user value is calculated by using the Equation (1).

$$M_{i,k}^{\text{BCQI}} = r_{i,k}^{\text{th}}(t) \quad \dots (1)$$

Proportional Fair (PF) [10] scheduler is implemented for maintaining the fairness and using the spectrum efficiently. It considers both the channel quality and users past average throughput for scheduling the user and offers sufficient amount of resources to more number of cell edge users. Equation (2) shows the resource allocation decision metric of 'ith' user value by PF algorithm.

$$M_{i,k}^{PF} = r_i^k(t) / R_i^k(t) \quad \dots (2)$$

In Equation (1), $R_i(t)$ defines the past average throughput of ith user at tth interval and $r_i(t)$ defines the achieved throughput of the ith user at n-1th Transmit Time Interval (TTI). RR, BCQI and PF are conventional schedulers in LTE. Duy Huy Nguyen, et al. [11] discussed the MLDWF scheduler by using virtual token and introduced the new factor called Maximum Queue Size (MQS) in order to reduce the packet loss rate. Maximum queue fails to identify the user demand and cannot predict changing the user's data rate. Aswin et al. [12] proposed a new downlink scheduler by introducing two phases which helps to increase the throughput and balance the fairness moderately.

Chiapin et al. [13] developed a delay reduced scheduler that assigns the resources to the real-time users followed by non real-time users. It enhances the QoS of real-time application which is designed to increase the spectral efficiency and also fairness of non real-time applications. In [14] authors introduced Buffer Aware Scheduling algorithm which maximizes the throughput of video applications by fixing the weight based priority. In order to achieve it, buffer size and channel conditions are considered as main parameters. It suffers when greater number of users requesting to get the service in the network. The highest priority is assigned to users who has higher channel condition. It does not check the users demand during the resource assignment.

Nandini et al. [15] compared several scheduling algorithms in WiMAX technology and proposed one cross layer algorithm to increase the throughput of real-time users by assigning the highest priority in which NS-2 has been used in analyzing the results. Mohammad et al. [16] done a comparative study of several downlink scheduling algorithms in finding the strength and weakness of LTE by using the parameters such as mobility speed, frame sizes, number of intervals, etc., to the users. Mohamad et al. [17] proposed an algorithm to increase the average cell throughput with degradation of system fairness. It is an extension of proportional fair algorithm used log based metric for balancing the throughput and fairness. It did not consider demands of the users. Nilam et al. [18] explained proportional fair algorithm and its weakness in offering services to the real-time applications. QoS metric with a proportional fair algorithm classifies the high priority and low priority users for achieving the higher throughput of real-time users.

Dejan et al. [19] developed a novel scheduling algorithm to admit each user at most once in the corresponding locations for enhancing the cell capacity and fairness. Proposed algorithm is compared with MAX C/I algorithm and the results were also discussed. Some of the existing downlink scheduling algorithms are focusing on channel quality alone. It never cares about the user demands which leads to the users with Minimum data rate (real-time) and Maximum bandwidth. The users who are very close to the Base Station access more resources than the users who are at longer distances. It leads to starvation of cell edge users.

Calduwel et al. [20] developed an algorithm to differentiate each codec user from one another by preparing a ranking order. This ranking increases the throughput and maintains the bandwidth of ongoing calls without any wastage. So, bandwidth efficiency is improved using Analytic Hierarchy Process based Call Admission Control algorithm. Calduwel et al. [21] proposed a priority based scheduling algorithm to reduce the average waiting time of the real-time users in the cellular coverage area. Rebekka et al. [22] developed a GBR algorithm to convince more number of real-time users by setting the bandwidth threshold values. It limits the user satisfaction when more users are crossing the threshold values. It is also noted that, the PF and BCQI do not consider the user requirements.

Chao He et al. [23] enhanced the QoS of real-time users by giving awareness to each user's QoS requirements independently. It sets a regulating factor for balancing the packet loss ratio with the help of MOS value which is being noted in every iteration of user changes. It is quite difficult to fix the exact threshold value in high mobility scenarios which results to unpredicted application requirements. By setting the few thresholds, this proposed work fails to estimate the user demand which increases the loss ratio and also reduces the throughput. Kanmani et al. [24] introduced blind channel estimation algorithm for reducing the cost value and Bit Error Rate in OFDMA. Generally, several channel estimation algorithms function without having the channel knowledge and its information which lead to resource damage. In [25] authors developed channel estimation algorithm for analysing in mountain environment using the pilot symbols based insertion. It is simulated for COST 207 channel model with 1024 subcarrier and reduces the noise rate. It poorly performs in increasing the throughput. Complexity reduced Least Square channel estimation algorithm is proposed by [26]. It reduces the BER by prefixing the channel conditions along with pilot symbol arrangement.

Reasons for Developing the Proposed Algorithm

Literatures reveal that, many schedulers are focusing only on the channel conditions, not considering into the user's speed along with their varying data rate. Most of the previous works consider the urgency packets with HoL delay, buffer size and etc., it does not occur at all the times. Whenever the number of users gets increased with same data rate which might result in varying users demand and mobility. It is difficult to predict their need since there is a possibility that multiple channels may have the same channel gain condition with different data rate. Some of the existing algorithms consider only the buffer size and application type.

When it fails to estimate an accurate value of channel condition, the scheduler assigns the resources wrongly, it turns into improper link adaptation. Finally, scheduler is unable to find the user demand correctly. In order to offer MIMO based service, base station should receive the correct channel conditions from the user's report. Inaccurate CQI values affect the communication between the Base Station and the User.

Channel condition should be in proper method in order to provide more reliable service. Some of the blind channel estimation algorithms fail to do this. So, there is a need to create scheduler which identifies the user requirement along with mobility prediction, henceforth the proposed RSA – QoS is introduced. The proposed work correctly identifies the variable bit-rate from the user device itself. Therefore, an innovative scheduler is required to reduce the loss ratio as well as increase the number of user's access. In general, every scheduling algorithm should be capable of identifying the user's expectations and their channel quality as well. To address all these issues, the following algorithm is proposed to enhance the QoS by

reducing the resource loss during resource allocation.

3. RSA – QoS: A PROPOSED ALGORITHM

RSA – QoS is developed for assigning the required amount of resources to VoIP users. So that, the users can be increased in number. In order to achieve this, the dynamic prediction is applied in each interval for finding the suitable share of every user by fixing certain limitations. It assigns resources to the users based on the demand by using CQI values to improve system capacity. The users of the network who are using the required resources inside the coverage area through which the total cell capacity can be determined. There are some existing algorithms which assign the resources to the users without considering their application types. Sometimes, Real-time applications may consume a constant amount of bandwidth unlike Non Real-time application takes variable amount of bandwidth frequently.

Input: Number of Users, Number of RBs, Number of Intervals and Types of AMR bit rates

Initialization:

$N_{RB} = \Phi$; $N_{UE} = \Phi$; $Tot_RB = \Phi$; $D_{UE} = \Phi$; $PD = 20ms$; $t = \Phi$; $t-1 = \Phi$; $bit_rate = \Phi$;

In RSA – QoS, the scheduler collects information N_{UE} and D_{UE} to assign the resources to the users. These metrics are used to develop (support) the scheduler which enhances the QoS in Mobile Networks.

Proc RSA_QoS (bit_rate, D_{UE})

Begin

$Tot_RB = N_{RB} \times PD$... (3)

for ($i \leftarrow 0; i < N_{UE}; i++$)

for ($j \leftarrow 0; j < N_{RB}; j++$)

$UE_i = [CQI_{1,1} \dots \dots \dots, CQI_{N,P}]$... (4)

$UE_i[D_{UE}] = (PD \times bit_rate) + UE_i[OH]$... (5)

If ($bit_rate[UE_i]$ in 't' time == $bit_rate[UE_i]$ in 't-1' time) *Then*

Set as Static_UE_i[bit_rate]

Else If ($bit_rate[UE_i]$ in 't' time != $bit_rate[UE_i]$ in 't-1' time) *Then*

Set as Dynamic_UE_i[bit_rate]

End If

End for

CALL RB_Assignment(CQI);

End

Proc RB_Assignment(CQI)

Begin

Divide the shares for both Static and Dynamic users according to the utility values

Distribute the resources by using the following

for ($i \leftarrow 0; i < N_{UE}; i++$)

```

    If ( $UE_i$  CQI is Low) Then // CQI 1-6
        Grant 3 RBs to each  $UE_i$ 
    Else If ( $UE_i$  CQI is Med) Then // CQI 7-9
        Grant 2 RBs to each  $UE_i$ 
    Else If ( $UE_i$  CQI is High) Then // CQI 10-15
        Grant 1 RBs to each  $UE_i$ 
    End If
End for
End

```

Output: placing the users one by one according to the channel condition and arrival time. Hence, Resource Loss is reduced by which QoS in Mobile networks can be enhanced. The used mathematical notations and symbols are listed in **Table 1**.

Table 1. Notations and their Meaning

Notation	Meaning
Bit_rate	AMR codec bit rate
t, t-1	Current and Previous time of request
N_{RB}	Number of RBs
N_{UE}	Number of UEs
Tot_RB	Total Number of Available RBs
D_{UE}	Demand of UE
PD	Packetization Delay (20ms)
OH	Overhead in Bytes
$CQI_{N,P}$	Channel Quality Indicator in ' N^{th} ' Interval and ' P^{th} ' RB

The total number of resources present in the scheduler is depicted in Equation (3). Equation (4) shows that the user CQI calculation for ' i^{th} ' user which can be repetitively done for all users. Equation (5) explains the ' i^{th} ' user's previous $(t-1)^{th}$ and current $(t)^{th}$ bandwidth. Similarly, due to user mobility and/or sudden changes in the network, the user faces congestion from one bandwidth range to another which might either be increased or decreased. In this situation, it is quite difficult to identify the user's demand and scheduler starts assigning the resources poorly. RSA – QoS estimates user needs and offers resources according to their requests. RSA – QoS is developed to allocate the resources efficiently rather than the existing algorithms like ANINTECH, EPF and GBR. RR does not consider user demand which leads to resource saturation. RSA – QoS addresses this issue and assigns the resource blocks in increasing the users in number. It also increases their average throughput and gives more opportunity to the users. It divides the shares equally, even when the user moves from one range into another range while the number of users is less than the other intervals. Generally, the AMR codec produces the speech (audio packets) exactly for every 20ms and consider the

overhead (OH) in bits which is constantly taken by RSA – QoS. It can be applicable to all ranges of AMR codecs in common used for simulation. Most of the codecs would be in 8 – 15 CQI range while calculating the total number of bits.

3.1 User Demand Estimation

User demand is calculated using the bit-rate and packetization delay. For example, AMR 7.4 bit-rate codec can produce 204 bits (Bandwidth X20ms+OH=148+56=204) for transmission. If a user is using the AMR codec (4.75 bit-rate) and the value of CQI 10 can be satisfied by giving one RB. According to [Table. 2 \[27\]](#), one RB in QPSK Modulation, the user can transmit 104 bits.

Table 2. Performance of Reference Channel [\[19\]](#)

Reference Channel	A3-1	A3-2	A3-3	A3-4	A3-5	A3-6	A3-7
Allocated resource blocks	1	6	15	25	50	75	100
DFT-OFDM Symbols per sub frame	12	12	12	12	12	12	12
Modulation	QPSK	QPSK	QPSK	QPSK	QPSK	QPSK	QPSK
Code rate	1/3	1/3	1/3	1/3	1/3	1/3	1/3
Payload size (bits)	104	600	1544	2216	5160	6712	10296
Transport block CRC (bits)	24	24	24	24	24	24	24
Code block CRC size (bits)	0	0	0	0	0	24	24
Number of code blocks – C	1	1	1	1	1	2	2
Coded block size including 12bits trellis termination (bits)	396	1884	4716	6732	15564	10188	15564
Total number of bits per sub-frame	288	1724	4320	7200	14400	21600	28800
Total symbols per sub-frame	144	864	2160	3600	7200	10800	14400

Codec based technique [\[28\]](#) adapts this demand for reducing the transmission delay. The RSA-QoS considers some of the AMR-NB Codecs like, 4.75, 6.6, 7.4, 7.95, 8.85 and 10.2 for analysis. A general assumption is made that the users are initially fixed in a position who are moving with little change in speed after a certain time interval. RSA – QoS is applied in all formats of 16 QAM and 64 QAM reference channel parameters to find the number of resource blocks based on the user demand, it assigns the resources to all of them. Other existing algorithms are providing services to less number of users comparing with the RSA – QoS.

CQI value is set to 6 from 1 and fixed a Region of Interest (RoI) value in the simulator for representing all 4.75 Codec users are located in cell edge area. After 20ms, the users are slowly moving towards base station, some users can access high bandwidth (6.6kbps) and their channel quality could also be high. These changes cause the dynamic mode to be active and user demand keeps fluctuating. It needs to differentiate from one codec type to another which is carried out in the proposed RSA – QoS by changing the mobility model.

3.2 RB Assignment

The RSA – QoS analyses the capacity of various CQIs during transmission of bits and found that, most of the AMR codec can be satisfied by giving one RB in CQI 10 to 15 ranges. So, RSA – QoS assigns one RB to the user who produces CQI value between 10 and 15. Similarly, for other ranges both Low (1-6) and Medium (7-9), RSA – QoS offers 3 RBs, 2RBs respectively. It increases the throughput of cell edge users because of more RBs are assigned to them if they appeared in greater numbers.

4. Scenarios and Discussions

Initially, the network scenario is configured by having 10 users, 3 MHz channel bandwidth, 15 RBs, prepared to run for two ms intervals (2TTIs) and shown in Fig. 3. Here, U_1 demands 2 RBs in AMR 23.85 codec bit rate, U_2 demands 2 RBs in AMR 19.85 codec bit rate, U_3 demands 2 RBs in AMR 18.25 codec bit rate, U_4 demands 2 RBs in AMR 15.85 codec bit rate, U_5 demands 3 RBs in AMR 12.6 codec bit rate, U_6 demands 3 RBs in AMR 12.2 codec bit rate, U_7 demands 3 RBs in AMR 14.25 codec bit rate, U_8 demands 4 RBs in AMR 10.2 codec bit rate, U_9 demands 4 RBs in AMR 5.5 codec bit rate and U_{10} demands 5 RBs in AMR 7.4 codec bit rate in standing at different locations.

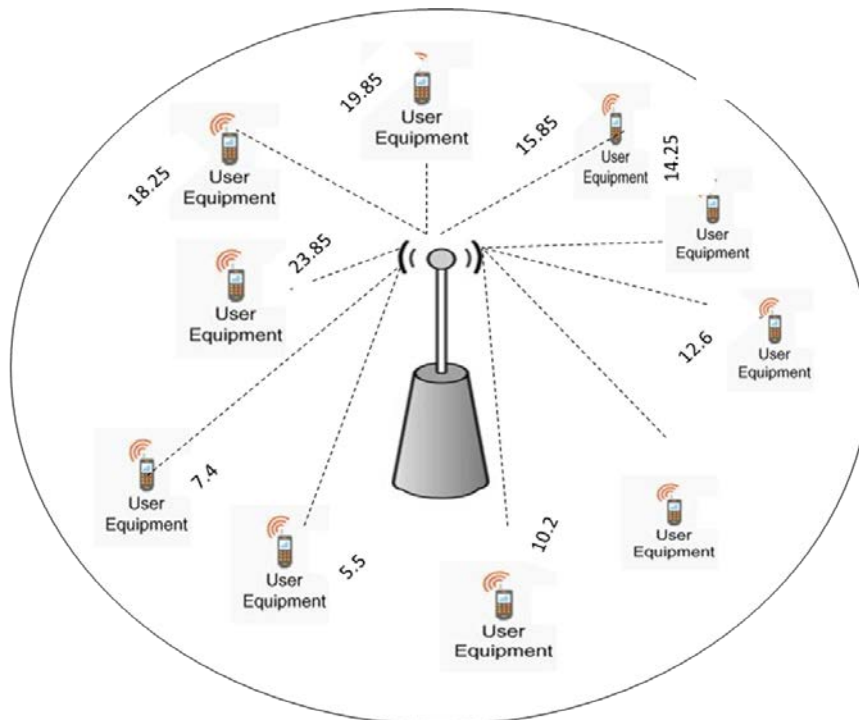


Fig. 3. Example Network Scenario

The total bit transmission (capacity) of each CQI is calculated by using this calculation: $1 \text{ RB} = 12 \text{ Subcarrier} \times 14 \text{ Symbols} = 168 \text{ REs}$, where PDCCH = 48 REs and Data = 140 REs. CQI 15 uses 64 QAM, efficiency = $0.948/102 = 0.926$, $1 \text{ RE} = 6 \times 0.926 = 5.5567$, therefore, $120 \times 5.5567 = 666 \text{ bits}$. It proves that, CQI 15 can carry 666 bits in one RB. An AMR codec

23.85 bit rate user is having CQI 12 can produce $(23.85 \times 20\text{ms}) + 56 = 533$ bits and can be satisfied by having only 2 RBs. The user in CQI 15 can transmit 666 bits and the same user in CQI 12 can transmit 468 bits, therefore for carrying 533 bits 2 RBs are enough. In such a way that, RSA – QoS predicts the demand and grants the resources to the users. Fig. 4 Shows the Resource Loss Ratio Comparison of existing ANINTECH, EPF, GBR and proposed RSA – QoS algorithm performance results in detail.

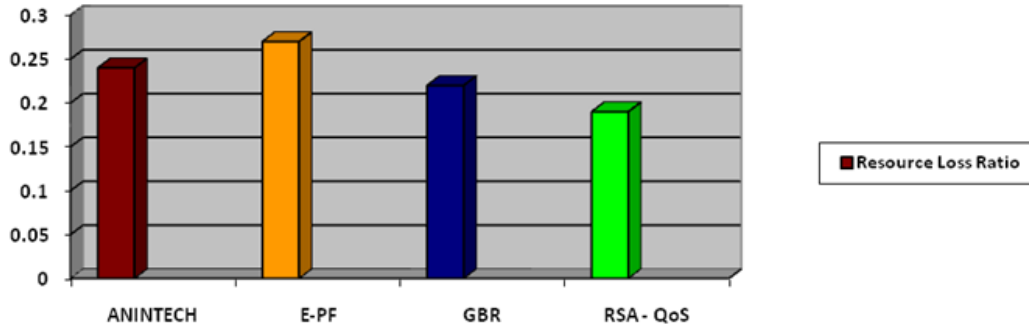


Fig. 4. Resource Loss Ratio Comparisons of Existing and Proposed Algorithms

5. Mathematical Model and Notations

The mathematical model of the RSA - QoS algorithm is discussed which adopts set theory. For users, resource allocation and classification of static – dynamic user identification. It concludes the number of resources which can be assigned to each user. The ultimate aim of this model is to bring out the working principle of the proposed algorithm in reducing the resource loss and increasing the throughput.

5.1 System Model

The proposed work considers the entire network as Set. It assumes that, set $'U' = \{u_1 \dots u_n\}$ which contains $'n'$ number of elements. Each element of the Set $'U'$ signifies a mobile user, where u_1 and u_n denote the first and last elements of the Set respectively. Scheduling is done by collecting the user's feedback in matrix format from each user. Then, the total number of resource blocks is calculated by multiplying the number of available resources per interval with the total number of intervals (i.e. No. of TTI \times No. of Resources per TTI). Equation (6) shows the bandwidth from the first user to the last user and particularly shown for $'i^{\text{th}}$ user.

$$u_i = \{b \text{ for } i = 1 \text{ to } n \quad \dots (6)$$

where $'u_i'$ refers to the $'i^{\text{th}}$ user bandwidth

Identify the bandwidth consumption of each user is dynamically changing or not. If all the users use a constant amount of bandwidth which is represented by $'b'$ and $'n'$ denotes the number of users in the static scenario and defined in Equation (7).

$$u_i = \begin{cases} b & \text{for } i = 1 \text{ to } r \\ c & \text{for } i = r + 1 \text{ to } n \end{cases} \quad \dots (7)$$

Equation (7) represents, after the certain interval of time, some of the users may be accessing the variable bandwidth due to its multi-rate feature of codec. Here, few users (1 to r) still are consuming the bandwidth 'b' and some other users (r+1 to n) are switching to another bandwidth 'c' where 'c' refers to the bandwidth of dynamic users and 'b' refers to the bandwidth of static users. Resource balancing decision is done by taking the utility values of both static and dynamic users who are residing in the coverage area. The total resource is equally distributed according to the utility value by applying Equations (8) and (9).

$$\text{Static user's share} = \frac{r}{n} mp \quad \dots (8)$$

$$\text{Dynamic user's share} = \frac{n-r}{n} mp \quad \dots (9)$$

After the share distribution for every user's CQI is taken to grant the number of resource blocks of each user. For example, resource allocation decision is defined in Equation (10).

$$\text{Resource Block of 'i' user} = \begin{cases} 1 \text{ RB, if } 10 \leq \text{CQI} \leq 15; \\ 2 \text{ RB, if } 7 \leq \text{CQI} \leq 9; \\ 3 \text{ RB, if } 1 \leq \text{CQI} \leq 6; \end{cases} \quad \dots (10)$$

5.2 Model Comparison

The corresponding resource block assignment in scheduling is defined by C_{ij} variable. It holds the binary relation {0 or 1} which is shown in Equation (11). The Number of Lost and No Lost users list is shown in Equations (12) and (13).

$$C_{ij} = \begin{cases} 1 \text{ if } i^{\text{th}} \text{ user is got assignment to } j^{\text{th}} \text{ RB} \\ 0 \text{ if } i^{\text{th}} \text{ user is not assigned to } j^{\text{th}} \text{ RB} \end{cases} \quad \dots (11)$$

$$\text{Let } R_L = \{ (i, j) / C_{ij} = 0 \} \quad \dots (12)$$

$$\text{Let } R_{NL} = \{ (i, j) / C_{ij} = 1 \} \quad \dots (13)$$

Where, $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$

Similarly, the total number of users and their assignment classification can be done by having the Equation (14).

$$n(R_L) + n(R_{NL}) = m \times n \quad \dots (14)$$

Where, $n(R_L)$ is the number of users unassigned (i.e. lost) and $n(R_{NL})$ is the number of users got assignment. Comparison is shown in Equations (15) and (16) uniquely.

$$n(R_{NL(P)}) > n(R_{NL(E)}) \quad \dots (15)$$

$$n(R_{L(P)}) < n(R_{L(E)}) \quad \dots (16)$$

The resource loss ratio comparison of both existing and RSA – QoS is described using the Equation (15) which shows that the RSA-QoS outperforms than the existing algorithm. That is, it is having a greater number of assigned resources to users than unassigned users due to loss ratio. Similarly, RSA – QoS yields less loss ratio than the existing algorithm which is described using the Equation (16). Hence, the RSA – QoS eliminates the users who are wasting the resources frequently and it is proven using the mathematical Equations (15) and (16).

Where, CQI – Channel Quality Indicator

i – i^{th} user Feedback

m – Number of Scheduling Intervals (1 to 20ms)

N – Total Number of Users for Resource Loss

n – Number of Users in the network

b – Bandwidth of users in static case

c – Bandwidth of users in dynamic case

r – Total number of Static Users

p – Number of Resources (i.e. depends on the channel bandwidth it gets varied)

u_i – Bandwidth of the ' i^{th} ' user where, $i=1,2,\dots,n$

T_1 and T_2 – Switching Time Intervals

$u_i^{(t)}$ – ' i^{th} ' user bandwidth on ' t^{th} ' time

$u_i^{(t-1)}$ – ' i^{th} ' user bandwidth on ' $t-1^{\text{th}}$ ' time

R – Binary variable for Resource Loss Identification

$n(R_{NL(P)})$ – Number of No Lost user in the Proposed (RSA-QoS) Algorithm

$n(R_{NL(E)})$ – Number of No Lost user in the Existing Algorithm

$n(R_{L(P)})$ – Number of Lost user in the Proposed (RSA-QoS) Algorithm

$n(R_{L(E)})$ – Number of Lost user in the Existing Algorithm

6. Performance Evaluation

A Downlink System Level Simulator [29] is used for configuring in both Single-Input Single-Output (SISO) and Multi-Input Multi-Output (MIMO) based network. A Single Base Station with a fixed number of UEs and Mobility management model are taken into consideration for forming the network. Users finite buffer based scenario takes less time whereas, infinite buffer based assumption consumes much time for estimating user's variable demand in short span of time. Initially, the network is configured by taking the parameters such as 1.44 MHz Channel Bandwidth, 20 ms Intervals and 100 Users in Free space path loss model of trial sector tilted environment. At very first time, the UEs are distributed randomly within the cell edge area in simulation. An assumption is made that, every user reports the CQI value between 1 and 6 in each interval. Aftermath, the scheduler assigns the resources by looking at their CQI range and codec bit-rate.

For Case I, each user is assumed to use AMR 4.75 bit-rate codec by having lower CQI values in all the intervals during the resource distribution. Simulation is carried out by choosing the Base Station '13' and users are placed constantly. In simulation, only the VoIP traffic is enabled by setting '1' in traffic model method and it is repeatedly done for Five different cases. Existing algorithms are incapable in finding the user demand and misplaces the resources frequently. BCQI offers RBs to the highest channel quality user again and again. It increases the individual user throughput and not the overall throughput. In EPF, an adjustable value is added into the PF and it produces increased average throughput than BCQI. In GBR, the threshold is set by this work then observed the number of users and throughput. ANINTECH offers the resources excessively to only limited number of users and achieves higher throughput in individual users. RSA – QoS shares the resource effectively to 40 users only by assigning 3 RBs and provides better results than the existing algorithms in both rural and urban path loss environments. **Table 3** shows the parameters that are used in this simulation of LTE network.

Table 3. Simulation Parameters and Values

Parameter	Values
Channel Bandwidth	1.44 MHz
Frequency	2.1 GHz
Sector Used	Tri Sector Tilted
Base Station	13
Scheduler	ANINTECH, EPF, GBR and RSA-QoS
Simulation TTIs	20
UEs per eNodeB	100
Path Loss Models	TS36942, Free Space
Path Loss Environment	Urban and Rural
Transmit Power	40dB
Channel Model	Winner+
UE distribution	Constant UEs per Cell
UE speed/Mobility	15/3.6 (1.38ms)
Application Type	VoIP (AMR Codec)

In Case II, users in 75% of QPSK (AMR 4.75 codec) type in the same position and the remaining 25% of QAM users in 6.6 bit-rate are considered. Also user speed is changed for this mobility. That is, QAM users were moving towards base station that allows for accessing more bandwidth. ANINTECH provides more number of resources to 6.6 codec type user than 4.75 codec users. GBR, EPF gives resources to 4.75 and 6.6 codec users somehow equally. RSA – QoS distributes half amount of resources to both codecs. Case III, users in 50% of QPSK where using AMR 4.75 codec type, 25% of the users 16 QAM in AMR 6.6 codec and remaining 25% of 64 QAM with 6.6 codecs are taken. RSA- QoS gives more resources for 6.6 codec users than others. Case IV, users of 25% are kept as in AMR 4.75 codec and the rest of the 75% of the users are in AMR 6.6 codec with 50% of 64 QAM and 25% of 16 QAM. More opportunity is given to users who resides near the Base Station in all scheduling algorithms. In this case, RSA – QoS produces more throughput in 6.6 codec rate than other algorithms. Case V is taken by considering all users in AMR 6.6 codec with 50% of 16 QAM and 50% 64 QAM modulations. 95% of the users got full amount of resources and is given by the RSA – QoS.

ANINTECH changes the priority order and gives more important to VoIP users. For this, it is edited in traffic model and is set to 0.5, then the rest of the applications are collectively put it as 0.5. It skips more number of users than EPF algorithm because, it assigns all the resources nearby tower users and did not consider both the farthest and medium distance users. Similarly, EPF performs slightly better than ANINTECH when comparing the number of users and an edge user throughput. It produces the higher loss ratio than ANINTECH. Round Robin assigns an equal amount of shares to all of them which performs poorly in throughput for the highest channel quality users. RSA – QoS shows an increased throughput when compared with the other existing algorithms. **Table 4** shows the total number of satisfied users in both the existing and RSA – QoS algorithm.

Table 4. Number of Contented Users in AMR

Test Cases	AMR 4.75 Mode				AMR 6.6 Mode				Total No. UEs			
	ANIN TEC H	E PF	GB R	RSA - QoS	ANIN TEC H	E PF	GB R	RSA - QoS	ANIN TEC H	E PF	GB R	RSA - QoS
Case I – T ₁ (100UE)	23	29	32	40	-	-	-	-	23	29	32	40
Case II – T ₁ & T ₂ (75UE) and (25UE)	8	14	18	30	17	14	20	15	25	28	38	45
Case III - T ₁ & T ₂ (50UE) and (50UE)	12	13	22	25	19	25	28	35	31	38	50	60
Case IV - T ₁ & T ₂ (25UE) and (75UE)	9	24	12	10	43	33	50	65	52	57	62	75
Case V- T ₂ (100UE)	-	-	-	-	71	70	76	95	71	70	76	95

Performance Metrics

The performance metrics which are used to evaluate the proposed RSA – QoS algorithm. RSA-QoS is compared with ANINTECH, EPF and GBR based on the various performance metrics like resource loss ratio and throughput.

Resource Loss Ratio (RLR)

The Resource Loss Ratio is calculated by taking the ratio of Number of Lost Resources into the Total Number of Received Resources and shown in Equation (17).

$$RLR = \frac{\text{No.of Lost Resources}}{\text{Total No.of Received Resources}} \times 100 \quad \dots (17)$$

6.1 Experimental Results and Analysis

The RSA – QoS and its comparisons are shown in Fig. 5. and it did not consider the users in High Mobility. Generally, users never wish to make a call while they were in high speed. So that, a little mobility is assumed and applied in the simulator for analyzing the results.

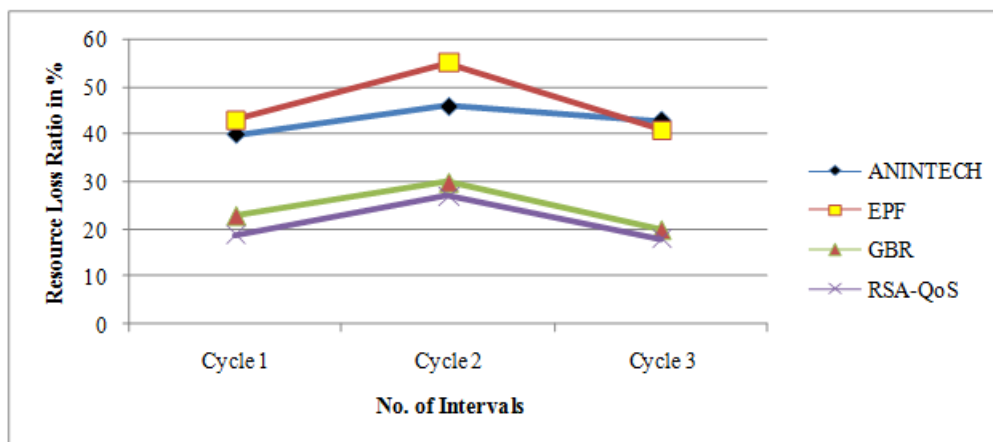


Fig. 5. Resource Loss Ratio Comparison for ANINTECH, EPF, GBR & RSA – QoS of 100 Users

RSA-QoS increases the throughput by reducing the resource loss even when channel bandwidth and the number of users get increased. Fig. 6 shows the throughput result of users for 3 cycles. Cycle 1 shows the RB assignment of 40 users, Cycle 2 shows another 40 users and Cycle 3 shows the RB assignment remaining users respectively. In this way, the whole 100 users can get resources appropriately. RSA-QoS is simulated to find the throughput performance of the schedulers in QPSK modulation for 3 cycles.

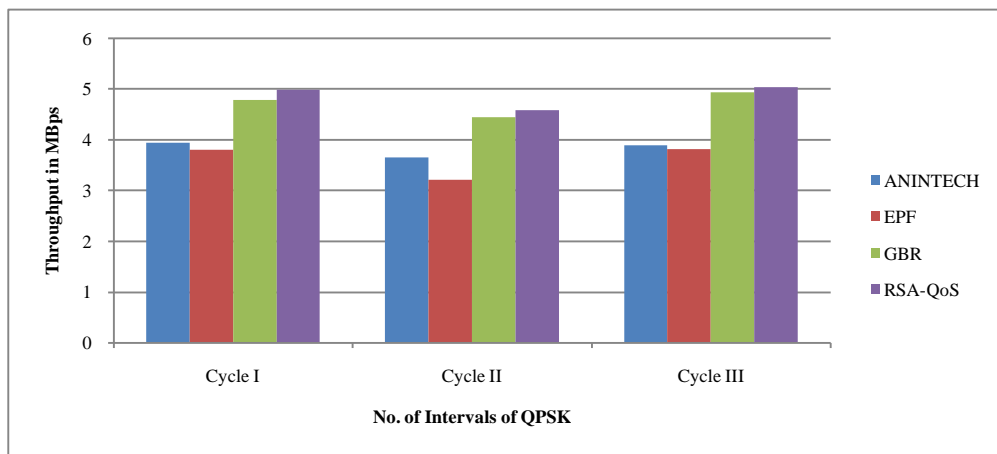


Fig. 6. Throughput Results of ANINTECH, EPF, GBR and RSA-QoS for QPSK

The same setup is used to test in another modulation 16 QAM for different intervals shown in Fig. 7. The throughput is increased gradually because RSA – QoS saves more number of resources.

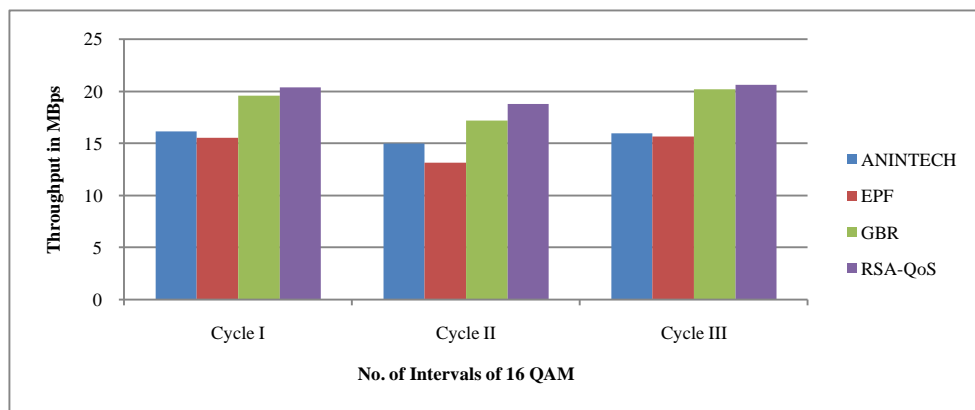


Fig. 7. Throughput Results of ANINTECH, EPF, GBR and RSA-QoS for 16 QAM

In 64 QAM assignment shows more number of user accessing the resource block without losing their intervals with the help of this demand based scheduler. Fig. 8 illustrates the throughput result of same users in 64 QAM.

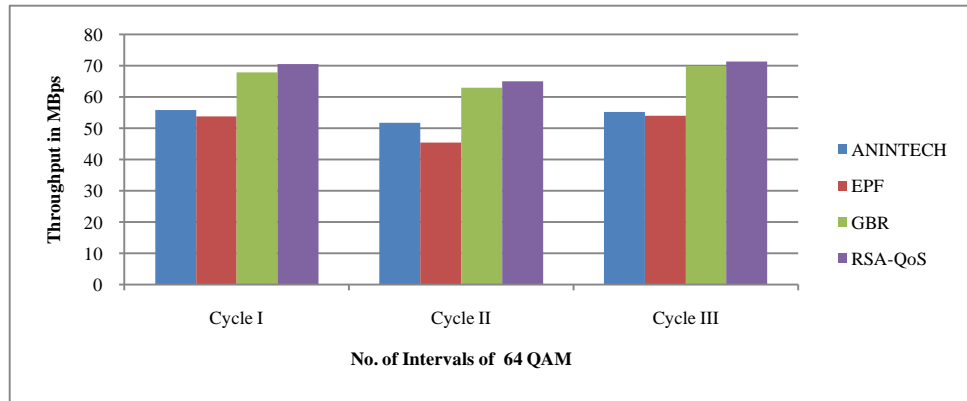


Fig. 8. Throughput Results of ANINTECH, EPF, GBR and RSA-QoS for 64 QAM

The RSA – QoS algorithm keeps on tracking the number of lost resources in every interval for offering to other users in very useful manner. **Fig. 9** displays the resource loss ratio performance in variable channel bandwidth, number of RBs in every (5ms) scheduling intervals.

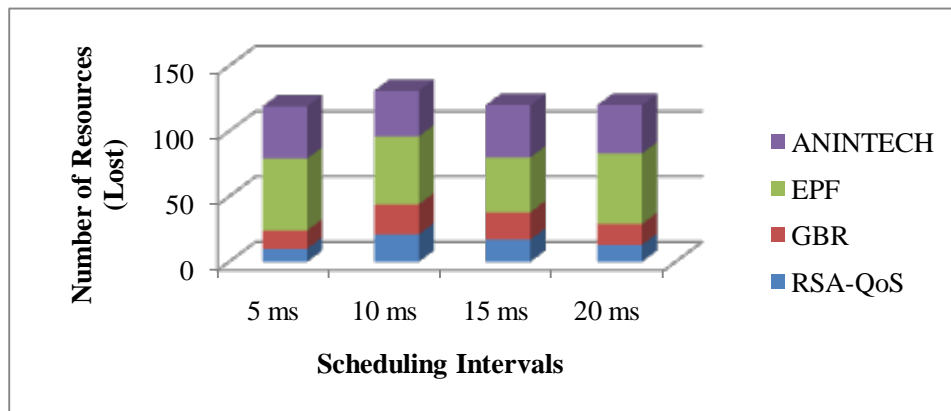


Fig. 9. Resource Loss Ratio for various Intervals

Resource loss ratio is measured by finding the ratio of the total number of missed resource into the total number of received resources of each user. RSA – QoS predicts the user demand and offers the resources according to them.

7. Conclusion

In several occasions, the real-time application takes the minimum data rate in transmission and may demand more number of resources with lesser delay residing in a particular sector. It creates a Multi-user channel sharing problem which limits the other user's access, those who are already existed and availed in that service area. If a scheduler assigns the resources to the users without estimating their demand, it would become ineffective. It leads to resource loss during transmission. Resource loss is a traditional issue which drops certain voice messages and finally, receiver can get only half of the message. This problem is addressed by the RSA – QoS algorithm. RSA – QoS allocates the resources to more number of

users even in various mobility scenarios. In RSA-QoS, High CQI gets only one RB which makes the way to save the resources and assigns them to the low CQI users. User's data rate is predicted at every interval and demand is measured to offer the resources appropriately. It is done by identifying the codec bit rate and its packetization delay of every user. The main aim of this paper is to reduce resource loss ratio and increases the throughput. Ultimately, it enhances the QoS in mobile networks.

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