

## AN UPPER BOUND ON THE CHEEGER CONSTANT OF A DISTANCE-REGULAR GRAPH

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**ABSTRACT.** We present an upper bound on the Cheeger constant of a distance-regular graph. Recently, the authors found an upper bound on the Cheeger constant of distance-regular graph under a certain restriction in their previous work. Our new bound in the current paper is much better than the previous bound, and it is a general bound with no restriction. We point out that our bound is explicitly computable by using the valencies and the intersection matrix of a distance-regular graph. As a major tool, we use the discrete Green's function, which is defined as the inverse of  $\beta$ -Laplacian for some positive real number  $\beta$ . We present some examples of distance-regular graphs, where we compute our upper bound on their Cheeger constants.

### 1. Introduction

A notion of the Cheeger constant of a graph has an important geometric meaning in graph theory. The Cheeger constant of a graph is closely related to the problem of separating a graph into two large components by making a small edge-cut. In fact, the Cheeger constant of a connected graph is strictly positive. If the Cheeger constant of a connected graph is “small”, then it means that there are two large sets of vertices with “few” edges between them. On the other hand, if a graph has “large” Cheeger constant, then it indicates that there are two sets of vertices with “many” edges between these two subsets. In general, computation of the Cheeger constant of a graph is a hard task. Only limited research has been done for finding the Cheeger constant of a graph. We are interested in finding bounds of Cheeger constants of graphs.

We begin with introducing some definitions in graph theory. Let  $\Gamma = (V, E)$  be a simple and connected graph, where  $V$  is the vertex set of  $\Gamma$  and  $E$  is the

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edge set of  $\Gamma$ . Let  $S$  be a nonempty subset of  $V$ . The *edge boundary* of  $S$ , denoted by  $\partial S$ , is defined as follows:

$$\partial S = \{\{x, y\} \in E \mid x \in S \text{ and } y \in V - S\}.$$

The *volume* of  $S$ , denoted by  $\text{vol}(S)$ , is defined as follows:

$$\text{vol}(S) = \sum_{u \in S} k_u,$$

where  $k_u$  is the valency of  $u$  in  $\Gamma$ . The *Cheeger ratio* of  $S$ , denoted by  $h_S$ , is defined as

$$h_S = \frac{|\partial S|}{\min\{\text{vol}(S), \text{vol}(\Gamma) - \text{vol}(S)\}}.$$

The *Cheeger constant* of  $\Gamma$ , denoted by  $h_\Gamma$ , is defined as

$$h_\Gamma = \min\{h_S \mid S \subseteq V\}.$$

Recent developments in [3, 4, 14] regarding distance-regular graphs show that there is a close connection between the Cheeger constant and vertex (or edge) connectivity. From Propositions A, B, and C we see that, for a distance-regular graph, there are close connections between the Cheeger constant and vertex and edge connectivity.

**Proposition A** ([3]). *Let  $\Gamma$  be a distance-regular graph with more than one vertex. Then its edge-connectivity equals its valency  $k$ , and the only disconnecting sets of  $k$  edges are the sets of edges incident with a single vertex.*

**Proposition B** ([4]). *Let  $\Gamma$  be a non-complete distance-regular graph of valency  $k > 2$ . Then the vertex-connectivity  $\kappa(\Gamma)$  equals  $k$ , and the only disconnecting sets of vertices of size not more than  $k$  are the point neighbourhoods.*

**Proposition C** ([14]). *Let  $\Gamma = (V, E)$  be a simple graph with the vertex-connectivity  $\kappa(\Gamma)$  and the edge-connectivity  $\lambda(\Gamma)$ . Then*

$$\frac{2\kappa(\Gamma)}{|V|} \leq \frac{2\lambda(\Gamma)}{|V|} \leq \inf \frac{|\partial S|}{|S|} \leq \kappa(\Gamma) \leq \lambda(\Gamma),$$

where  $S$  is a subset of  $V$  with  $|S| \leq \frac{|V|}{2}$ .

The Cheeger constants [9, 15] are related to the eigenvalues of the Laplacians of distance-regular graphs, and their eigenvalues are also involved with the *intersection numbers* of distance-regular graphs [13, 16]. However, in general, it is a hard task to compute the Cheeger constant of a distance-regular graph. Distance-regular graphs introduced by Biggs [2] are connected with coding theory and design theory; well-known examples of distance-regular graphs are the Hamming graphs and the Johnson graphs. In [11, 12], by using the relationship between a discrete Green's function and the Cheeger constant, we obtain an upper bound on the Cheeger constant of a distance-regular graph under a certain condition.

We find a general upper bound on the Cheeger constant of a distance-regular graph with no additional condition. Furthermore, our bound is a much more improved one comparing with the bound in [12] under the same additional condition; in Example 7 and Example 8, we show that our bound is much more improved one comparing with the bound in [12] under the same additional condition:  $\beta vr_d^{(\beta)} > \frac{\lambda_1}{1+\lambda_1}$ .

We point out that our bound is explicitly computable by using the valencies and the *intersection matrix* of a distance-regular graph; first, our bound is expressed in terms of  $q$ -numbers, and in general, it is not easy to compute the  $q$ -numbers. For resolving this problem, we obtain an alternative expression of our bound using the valencies and the intersection matrix of a distance-regular graph. In Example 10 and Example 11, we compute the upper bound on the Cheeger constant using the alternative expression in Theorem 2 and Remark 6. As a major tool, we use the discrete Green's function, which is defined as the inverse of  $\beta$ -Laplacian for some positive real number  $\beta$ . We present some examples which show our upper bound on the Cheeger constant for some distance-regular graphs.

We discuss our main result in more detail for the rest of this section. In this paper, we study distance-regular graphs. Let  $\Gamma = (V, E)$  be a distance-regular graph of order  $v$ , diameter  $d$  and valency  $k$ . Let  $A_1$  be the adjacency matrix of  $\Gamma$  and  $P$  be the transition probability matrix of  $\Gamma$ . Two adjacent vertices  $x, y$  are denoted by  $x \sim y$ . For a function  $f : V \rightarrow \mathbb{R}$ , we define a Laplace operator  $\Delta$  by  $\Delta f(x) = \frac{1}{k} \sum_{y \sim x} (f(x) - f(y))$ . Then  $\Delta = I - \frac{1}{k} A_1$ . Let  $\mathcal{L}_\beta$  be the  $\beta$ -normalized Laplacian  $\beta I + \Delta$ . For  $\beta > 0$ , let  $\mathcal{G}_\beta$  be a discrete Green's function denoted by the symmetric matrix which satisfies  $\mathcal{L}_\beta \mathcal{G}_\beta = I$ ; that is,  $\mathcal{G}_\beta$  is defined as the inverse of the  $\beta$ -Laplacian  $\mathcal{L}_\beta$  [5, 6, 7]. As in [11], for any positive real number  $\beta$ , let  $r_i^{(\beta)}$  ( $i = 0, 1, \dots, d$ ) denote the components of a Green's function  $\mathcal{G}_\beta$ . We define  $\alpha_i$  to be the limit of a sequence  $\{\alpha_i^{(\beta)}\}$  as  $\beta$  goes to  $0^+$ , where

$$(1) \quad \alpha_i^{(\beta)} = \frac{\beta^2 vr_i^{(\beta)}}{1 - \beta vr_i^{(\beta)}} \quad (i = 0, 1, \dots, d).$$

In fact, we can express  $\alpha_i$ 's by the eigenvalues  $\lambda_j$  of the Laplacian  $\mathcal{L}_\beta$  and the  $q$ -numbers  $q_j(i)$  of the P-polynomial scheme [1, 8, 10, 11]:

$$(2) \quad \alpha_i = \frac{1}{-q_1(i) \frac{1}{\lambda_1} - \dots - q_d(i) \frac{1}{\lambda_d}}.$$

We also see that  $0 < \alpha_d < \alpha_{d-1} < \dots < \alpha_e < \lambda_1$  for some  $e$ .

The authors obtain the following result [12] on an upper bound on the Cheeger constant of a distance-regular graph with a certain restricted condition as follows.

**Theorem A** ([12]). *Let  $\Gamma$  be a distance-regular graph with diameter  $d$  and  $\beta vr_d^{(\beta)} > \frac{\lambda_1}{1+\lambda_1}$  for  $\beta \leq \alpha_d$ . Let  $\lambda_1$  be the smallest eigenvalue of the Laplacian.*

Then we have

$$\lambda_1 h_\Gamma < \alpha_d < \alpha_{d-1} < \dots < \alpha_e < \lambda_1$$

for  $e \in \mathbf{C}'_\beta$ .

Main results of this paper are the following Theorem 1, Theorem 2 and Corollary 3. Theorem 1 presents an upper bound on the Cheeger constant of a distance-regular graph. In Theorem 2, we find an explicit expression for the bound given in Theorem 1 by using the valencies  $k_j$  and the basis of nullspace  $\mathcal{N}(L_{sub}^{(\alpha_d)})$ . This shows that our new bound is a computable bound using the valencies and the intersection matrix of a distance-regular graph. Corollary 3 shows that our generalized bound in Theorem 1 and Theorem 2 improves the bound given in Theorem A [12] under the same additional condition.

**Theorem 1.** *Let  $\Gamma$  be a distance-regular graph of diameter  $d$ . Then we have the following upper bound:*

$$h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}},$$

where  $\alpha_d = \lim_{\beta \rightarrow 0^+} \frac{\beta^2 vr_d^{(\beta)}}{1 - \beta vr_d^{(\beta)}}$  and  $\alpha_d^{(\alpha_d)} = \frac{\alpha_d^2 vr_d^{(\alpha_d)}}{1 - \alpha_d vr_d^{(\alpha_d)}}$ .

**Theorem 2.** *Let  $\Gamma$  be a distance-regular graph of order  $v$  and diameter  $d$ . Let  $(u_0^{(\alpha_d)}, u_1^{(\alpha_d)}, \dots, u_d^{(\alpha_d)})$  be a basis of  $\mathcal{N}(L_{sub}^{(\alpha_d)})$  with  $u_d^{(\alpha_d)} = 1$  as in Lemma 4. Then we have*

$$h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} = \alpha_d \left( \frac{1}{v} \sum_{j=0}^d k_j u_j^{(\alpha_d)} - 1 \right),$$

where  $k_j$  are valencies as in Lemma 4 and  $\alpha_d$  is the same as in Lemma 5.

**Corollary 3.** *Let  $\Gamma$  be a distance-regular graph, and  $h_\Gamma$  be a Cheeger constant of  $\Gamma$ . If  $\beta vr_d^{(\beta)} > \frac{\lambda_1}{1 + \lambda_1}$  for  $\beta \leq \alpha_d$ , then we have*

$$h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} < \frac{\alpha_d}{\lambda_1},$$

where  $\lambda_1$  is the smallest positive eigenvalue of the Laplacian and  $\alpha_d$  is the same as in (1).

In Section 2, we introduce some notations and facts about distance-regular graph and some properties of the Green's function  $\mathcal{G}_\beta$ . In Section 3, we find a new upper bound on the Cheeger constant of a distance-regular graph. We also obtain an alternative expression of our upper bound by using the valencies  $k_j$  and the basis of nullspace  $\mathcal{N}(L_{sub}^{(\alpha_d)})$ . Finally, in Section 4, we present some examples about our upper bound on the Cheeger constant of some distance-regular graphs.

### 2. Preliminaries and Green's function

We introduce definitions of the distance-regular graphs and the  $P$ -polynomial schemes. A connected graph  $\Gamma$  with diameter  $d$  is called a *distance-regular graph* if there are constants  $c_i, a_i, b_i$  such that for all  $i = 0, 1, \dots, d$ , and all vertices  $x$  and  $y$  at distance  $i = d(x, y)$ , among the neighbors of  $y$ , there are  $c_i$  at distance  $i - 1$  from  $x$ ,  $a_i$  at distance  $i$ , and  $b_i$  at distance  $i + 1$ . It follows that  $\Gamma$  is a regular graph with valency  $k = b_0$ , and that  $c_i + a_i + b_i = k$  for all  $i = 0, 1, \dots, d$ . By these equations, the intersection numbers  $a_i$  can be expressed in terms of the others, and it is a standard to put these others in the so-called *intersection array*  $(b_0, b_1, \dots, b_{d-1}; c_1, c_2, \dots, c_d)$ . We describe the relations by its adjacency matrices  $A_i$  ( $i = 0, 1, \dots, d$ ) which are  $v \times v$  matrices defined by

$$(A_i)_{x,y} = \begin{cases} 1 & \text{if } (x, y) \in R_i, \\ 0 & \text{otherwise,} \end{cases}$$

[1, 8]. Let  $X$  be a nonempty finite set and  $R = \{R_0, R_1, \dots, R_d\}$  be a family of relations defined on  $X$ . We say that the pair  $(X, R)$  is a symmetric association scheme with  $d$  classes if it satisfies the following conditions.

- (1)  $A_0 = I$  (identity matrix).
- (2)  $A_0 + A_1 + \dots + A_d = J$  (all 1 matrix).
- (3)  $A_i A_j = \sum_{k=0}^d p_{ij}^k A_k$ , where  $p_{ij}^k$  is the number of  $z \in X$  such that  $(x, z) \in R_i$  and  $(z, y) \in R_j$ .
- (4)  $A_j^t = A_j$ .
- (5)  $A_i A_j = A_j A_i$ .

A symmetric association scheme  $\mathfrak{X} = (X, R)$  is called a  $P$ -polynomial scheme with respect to the ordering  $R_0, R_1, \dots, R_d$ , if there exist some complex coefficient polynomials  $v_i(x)$  of degree  $i$  ( $i = 0, 1, \dots, d$ ) such that  $A_i = v_i(A_1)$ , where  $A_i$  is the adjacency matrix with respect to  $R_i$ .

It is known [1, 8] that a distance-regular graph is equivalent to a  $P$ -polynomial scheme  $\mathfrak{X}$  with respect to some relations  $R_0, R_1, \dots, R_d$  on a vertex set  $V$  with  $|V| = v$ . Thus, we can define the Green's function over a  $P$ -polynomial scheme, and then by using the Green's function we will obtain an upper bound on the Cheeger constant of a distance-regular graph.

The first intersection matrix  $B_1$  of a distance-regular graph is a tridiagonal matrix with non-zero off diagonal entries:

$$B_1 = \begin{pmatrix} 0 & k & 0 & 0 & \cdots & 0 \\ 1 & a_1 & b_1 & 0 & \cdots & 0 \\ 0 & c_2 & a_2 & b_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ & & & c_{d-1} & a_{d-1} & b_{d-1} \\ 0 & \cdots & \cdots & 0 & c_d & a_d \end{pmatrix} \quad (b_i \neq 0, c_i \neq 0).$$

Let  $\mathcal{A}$  be the algebra spanned by the adjacency matrices  $A_0, A_1, \dots, A_d$ . Then  $\mathcal{A}$  is called the *Bose-Mesner algebra* of  $\mathfrak{X}$ , and  $\mathcal{A}$  has two distinguished

bases  $\{A_i\}$  and  $\{E_i\}$ , where the latter consists of primitive idempotent matrices. For  $A_i$  and  $E_i$ , we can express one in terms of the other as the following:

$$A_j = \sum_{i=0}^d p_j(i)E_i, \quad E_j = \frac{1}{|X|} \sum_{i=0}^d q_j(i)A_i$$

for  $j = 0, 1, \dots, d$ . The  $(d + 1) \times (d + 1)$  matrix  $\mathbf{P} = (p_j(i))$  (respectively,  $\mathbf{Q} = (q_j(i))$ ) is called *the first eigenmatrix* (respectively, the second eigenmatrix) of the P-polynomial scheme  $\mathfrak{X}$ , where  $p_j(i)$  (respectively,  $q_j(i)$ ) is a  $p$ -number (respectively, a  $q$ -number).

As defined in Section 1, the Green's function  $\mathcal{G}_\beta$  is the inverse of the  $\beta$ -Laplacian  $\mathcal{L}_\beta$ . For  $\beta > 0$ , we thus have  $\mathcal{G}_\beta(\beta I + I - P) = I$ . Therefore, we get

$$\mathcal{L}_\beta = \sum_{j=0}^d \left( \beta + 1 - \frac{1}{k_1} p_1(j) \right) E_j, \quad \mathcal{G}_\beta = \sum_{j=0}^d \left( \frac{k_1}{(\beta + 1)k_1 - p_1(j)} \right) E_j.$$

Since  $E_j = (1/v) \sum q_j(i)A_i$  and  $\lambda_j = 1 - p_1(j)/k_1$ , we get

$$\begin{aligned} \mathcal{G}_\beta &= \sum_{j=0}^d \left( \frac{k_1}{(\beta + 1)k_1 - p_1(j)} \right) \sum_{i=0}^d q_j(i) \\ &= \sum_{i=0}^d \sum_{j=0}^d (1/v) \left( \frac{k_1}{(\beta + 1)k_1 - p_1(j)} \right) q_j(i) A_i \\ &= \sum_{i=0}^d \sum_{j=0}^d (1/v) \left( \frac{1}{\beta + 1 - \frac{p_1(j)}{k_1}} \right) q_j(i) A_i \\ &= \sum_{i=0}^d \sum_{j=0}^d (1/v) \left( \frac{1}{\beta + \lambda_j} \right) q_j(i) A_i. \end{aligned}$$

That is,  $\mathcal{G}_\beta$  is a linear combination of adjacency matrices  $A_i$  as follows:

$$\mathcal{G}_\beta = r_0^{(\beta)} A_0 + r_1^{(\beta)} A_1 + \dots + r_d^{(\beta)} A_d,$$

where  $r_i^{(\beta)} = \frac{1}{v} \left( \frac{1}{\beta} + q_1(i) \frac{1}{\beta + \lambda_1} + \dots + q_d(i) \frac{1}{\beta + \lambda_d} \right)$  ( $i = 0, 1, \dots, d$ ).

In [10, 11], a  $d \times (d + 1)$  matrix  $L_{sub}^{(\beta)}$  is introduced as a matrix obtained by the removal of the first row of  $B_1 - k_1(\beta + 1)I$ .

**Lemma 4** ([11]). *For  $\beta > 0$ , let  $\mathcal{G}_\beta = r_0^{(\beta)} A_0 + r_1^{(\beta)} A_1 + \dots + r_d^{(\beta)} A_d$  be the Green's function of a distance-regular graph  $\Gamma$  of order  $v$ . Then we have*

- (a)  $\mathcal{G}_\beta$  can be expressed as  $\mathcal{G}_\beta = t u_0^{(\beta)} A_0 + t u_1^{(\beta)} A_1 + \dots + t u_d^{(\beta)} A_d$  for some nonzero  $t \in \mathbb{R}$ , where  $(u_0^{(\beta)}, u_1^{(\beta)}, \dots, u_d^{(\beta)})$  is the unique basis of the nullspace  $\mathcal{N}(L_{sub}^{(\beta)})$  of  $L_{sub}^{(\beta)}$  with  $u_d^{(\beta)} = 1$ .

- (b)  $k_0 r_0^{(\beta)} + k_1 r_1^{(\beta)} + \dots + k_d r_d^{(\beta)} = \frac{1}{\beta}$ , where  $k_j$  is the valency of  $A_j$  for  $j = 0, 1, \dots, d$ .
- (c)  $r_0^{(\beta)} > r_1^{(\beta)} > \dots > r_d^{(\beta)} > 0$ .
- (d)  $\lim_{\beta \rightarrow 0^+} |r_i^{(\beta)} - r_j^{(\beta)}| = 0$  for  $0 \leq i, j \leq d$ .

**3. A new improved bound on the Cheeger constant**

In this section we prove Theorem 1 and Theorem 2. We need the following lemma for the proof of Theorem 2 and Theorem 3. We consider a set  $\mathbf{C}_\beta = \{i \mid \frac{1}{\beta} - v r_i^{(\beta)} > 0\}$  as a subset of  $\{0, 1, 2, \dots, d\}$ ; then  $\mathbf{C}_\beta$  is a non-empty set by Lemma 4. When  $\beta$  is sufficiently close to  $0^+$ , we consider a set  $\mathbf{C}'_\beta = \{i \mid \beta v r_i^{(\beta)}(\beta + \lambda_1) < \lambda_1\}$ ; then  $\mathbf{C}'_\beta$  is a subset of  $\mathbf{C}_\beta$ .

**Lemma 5** ([11]). *For  $\beta > 0$ , let  $\Gamma$  be a distance-regular graph of order  $v$ , and let  $\mathcal{G}_\beta = r_0^{(\beta)} A_0 + r_1^{(\beta)} A_1 + \dots + r_d^{(\beta)} A_d$  be a Green's function of  $\Gamma$ . We recall that  $\alpha_i^{(\beta)} := \frac{\beta^2 v r_i^{(\beta)}}{1 - \beta v r_i^{(\beta)}} (i = 0, 1, \dots, d)$  as given in Eq. (1). Then for  $i \in \mathbf{C}'_\beta$ , we have the following:*

- (a)  $\lim_{\beta \rightarrow 0^+} \beta v r_i^{(\beta)} = 1^-$ ,  $\lim_{\beta \rightarrow 0^+} \beta^2 v r_i^{(\beta)} = 0^+$ .
- (b)  $\alpha_i^{(\beta)}$  is decreasing in  $i \in \mathbf{C}'_\beta$ .
- (c) There exists  $i \in \mathbf{C}'_\beta$  such that  $\lim_{\beta \rightarrow 0^+} \alpha_i^{(\beta)} = \alpha_i < \lambda_1$ .
- (d)  $\alpha_i^{(\beta)}$  is decreasing in  $\beta > 0$ .

**Proof of Theorem 1.** Let  $S$  be a subset of the vertex set  $V$  of  $\Gamma$  with  $\text{vol}(S) \leq \text{vol}(V)/2$  and  $\frac{|\partial S|}{\text{vol}(S)} \neq h_\Gamma$ . Let  $S'$  be a subset of  $V$  with  $\frac{|\partial S'|}{\text{vol}(S')} = h_\Gamma$ . Then there exists some  $\beta$  with  $0 < \beta < 1$  such that

$$(3) \quad \frac{|\partial S| \beta}{\text{vol}(S)} = \frac{|\partial S'|}{\text{vol}(S')}.$$

We first note that for a positive integer  $n$ ,

$$(4) \quad \beta < \frac{\alpha_d^{(\alpha_d \beta^n)}}{\alpha_d^{(\alpha_d)}};$$

this follows immediately from  $\frac{\alpha_d^{(\alpha_d)}}{\alpha_d} < \frac{\alpha_d^{(\alpha_d \beta^n)}}{\alpha_d \beta}$ , which is clear since  $\alpha_d^{(x)}$  is decreasing in  $x$  by Lemma 5.

We claim that for any  $\epsilon > 0$ , there exists a positive integer  $N_\epsilon$  such that

$$(5) \quad \alpha_d^{(\alpha_d \beta^n)} < \alpha_d^2 + \epsilon$$

for any  $n \geq N_\epsilon$ ; we use Lemma 5 for the proof as follows. Let  $f_n = \alpha_d \beta^n v r_d^{(\alpha_d \beta^n)}$ . Then  $\lim_{n \rightarrow \infty} f_n = 1^-$  by Lemma 5(a). Thus, for sufficiently large positive integer  $n$ , we obtain the following approximations:

$$f_n(\beta^n + \alpha_d) \approx \alpha_d$$

$$\begin{aligned} \Rightarrow f_n \beta^n &\approx \alpha_d (1 - f_n) \\ \Rightarrow \alpha_d^{(\alpha_d \beta^n)} &= \frac{f_n \beta^n \alpha_d}{1 - f_n} \approx \alpha_d^2; \end{aligned}$$

so our claim in Eq. (5) follows.

From Eq. (4) and Eq. (5), we thus have that

$$(6) \quad \beta < \frac{\alpha_d^2 + \varepsilon}{\alpha_d^{(\alpha_d)}}.$$

Taking  $\varepsilon > 0$  to be such that  $\varepsilon < (\frac{\text{vol}(S)}{|\partial S|} - 1) \alpha_d^2$  (noting that the right hand side of this inequality is positive), we obtain

$$h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}};$$

this is because from Eq. (3) and Eq. (6), we get the following:

$$h_\Gamma = \frac{|\partial S| \beta}{\text{vol}(S)} < \frac{|\partial S|}{\text{vol}(S)} \left( \frac{\alpha_d^2 + \varepsilon}{\alpha_d^{(\alpha_d)}} \right) < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}}.$$

Consequently, the result follows as desired.  $\square$

**Proof of Theorem 2.** From Theorem 1, we have

$$h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}}.$$

By Lemma 5, we have

$$\alpha_d = \lim_{\beta \rightarrow 0^+} \frac{\beta^2 v r_d^{(\beta)}}{1 - \beta v r_d^{(\beta)}}, \quad \alpha_d^{(\alpha_d)} = \frac{\alpha_d^2 v r_d^{(\alpha_d)}}{1 - \alpha_d v r_d^{(\alpha_d)}}.$$

Thus,

$$(7) \quad \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} = \frac{1 - \alpha_d v r_d^{(\alpha_d)}}{v r_d^{(\alpha_d)}} = \frac{1}{v r_d^{(\alpha_d)}} - \alpha_d,$$

where  $v r_d^{(\alpha_d)} = \frac{1}{\alpha_d} + q_1(d) \frac{1}{\alpha_d + \lambda_1} + \cdots + q_d(d) \frac{1}{\alpha_d + \lambda_d}$ . From Lemma 4, we have

$$\sum_{j=0}^d k_j r_j^{(\alpha_d)} = \frac{1}{\alpha_d},$$

and this implies that

$$r_d^{(\alpha_d)} \sum_{j=0}^d k_j u_j^{(\alpha_d)} = \frac{1}{\alpha_d};$$

so we get

$$\alpha_d \sum_{j=0}^d k_j u_j^{(\alpha_d)} = \frac{1}{r_d^{(\alpha_d)}}.$$



It thus follows that

$$(8) \quad \frac{1}{vr_d^{(\alpha_d)}} - \alpha_d = \alpha_d \left( \frac{1}{v} \sum_{j=0}^d k_j u_j^{(\alpha_d)} - 1 \right).$$

The result follows immediately by combining Eq. (7) with Eq. (8).  $\square$

The following remark shows that  $u_j^{(\beta)}$  can be expressed by a determinant of a submatrix  $L_j^{(\beta)}$  of  $\mathcal{N}(L_{sub}^{(\beta)})$ , and  $\alpha_d$  can be expressed in terms of a basis  $(u_0^{(\beta)}, u_1^{(\beta)}, \dots, u_d^{(\beta)})$  of  $\mathcal{N}(L_{sub}^{(\beta)})$  and the valencies  $k_j$ 's as in Lemma 4.

*Remark 6.* (1) [10, 11] For  $\beta > 0$ , let  $L_0^{(\beta)}$  be the  $d \times d$  matrix obtained by the removal of the first column of  $L_{sub}^{(\beta)}$  as in Lemma 4. Let  $L_j^{(\beta)}$  be the  $(d-j) \times (d-j)$  matrix obtained by the removal from the first row (respectively, column) to the  $j$ -th row (respectively, column) of  $L_0^{(\beta)}$ , and let  $(u_0^{(\beta)}, u_1^{(\beta)}, \dots, u_d^{(\beta)})$  be a basis of  $\mathcal{N}(L_{sub}^{(\beta)})$  with  $u_d^{(\beta)} = 1$ . Then we have

$$u_j^{(\beta)} = (-1)^{d-j} \frac{\det(L_j^{(\beta)})}{c_{j+1}c_{j+2} \cdots c_d}, \quad j = 0, 1, \dots, d-1,$$

where  $\det(L_d^{(\beta)}) = 1$ .

(2) [11] Let  $\Gamma$  be a distance-regular graph of order  $v$ , and let  $\mathcal{G}_\beta = r_0^{(\beta)} A_0 + r_1^{(\beta)} A_1 + \cdots + r_d^{(\beta)} A_d$  be a Green's function of  $\Gamma$  for  $\beta > 0$ . Then we have

$$\alpha_d = \lim_{\beta \rightarrow 0^+} \frac{\beta v}{\sum_{j=0}^d k_j u_j^{(\beta)} - v} \quad \text{and} \quad \alpha_d < \alpha_d^{(\beta)} + \beta.$$

**Proof of Corollary 3.** Since  $\beta vr_d^{(\beta)} > \frac{\lambda_1}{1+\lambda_1}$  for  $\beta \leq \alpha_d$ , we have

$$\lambda_1 < \frac{\beta vr_d^{(\beta)}}{1 - \beta vr_d^{(\beta)}} = \frac{\alpha_d^{(\beta)}}{\beta}.$$

Letting  $\beta = \alpha_d$ , we get  $\frac{\alpha_d}{\alpha_d^{(\alpha_d)}} < \frac{1}{\lambda_1}$ . It thus follows

$$\frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} < \frac{\alpha_d}{\lambda_1}. \quad \square$$

#### 4. Examples

In this section we present some examples regarding our upper bound on the Cheeger constants for some distance-regular graphs. In particular, Example 7 and Example 8 show that our bound is much more improved one comparing with the bound in [11, 12] under the same additional condition.

Theorem 1 shows an upper bound on the Cheeger constant  $h_\Gamma$  in terms of  $\alpha_d$  and  $\alpha_d^{(\beta)}$ . From Equations (2) and (3), if we know the  $q$ -numbers of the given  $P$ -polynomial scheme, then we can find  $\alpha_d$  and  $\alpha_d^{(\beta)}$  immediately.

In the Hamming scheme  $H(d, q)$  (respectively, Johnson scheme  $J(m, d)$ ), the  $p$ -number  $p_j(i)$  is defined by the Krawtchouk polynomial (respectively, the Eberlein polynomial) [1]. Since  $\mathbf{PQ} = v\mathbf{I}$ , we can obtain the  $q$ -numbers  $q_j(i)$  of the Hamming scheme  $H(d, q)$  and the Johnson scheme  $J(m, d)$ . We present the following two examples for showing this case.

**Example 7.** Let  $\Gamma$  be the graph of the Hamming scheme  $H(d, q)$  with respect to  $A_1$ . Then  $\Gamma$  is a distance-regular graph with  $q^d$  vertices, valency  $d(q-1)$  and  $d$  diameter. We consider two cases: (a)  $d = 5$ ,  $q = 4$  and (b)  $d = 7$ ,  $q = 3$ , and in each case, our upper bound on the Cheeger constant is as follows:

- (a)  $H(5, 4)$  :  $v = 1024$ ,  $k_1 = 15$ ,  $\lambda_1 = 4/15$ ,  $\alpha_d = 16/137$ .  
 And, for  $\beta \leq \alpha_d$ ,  $\beta vr_d^{(\beta)} \geq \alpha_d vr_d^{(\alpha_d)} \approx 0.41256 > \frac{\lambda_1}{1+\lambda_1} \approx 0.21053$ .  
 Thus,  $h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.166291 < \frac{\alpha_d}{\lambda_1} \approx 0.43795$ .
- (b)  $H(7, 3)$  :  $v = 2187$ ,  $k_1 = 14$ ,  $\lambda_1 = \frac{3}{14}$ ,  $\alpha_d = \frac{10}{121}$ .  
 And, for  $\beta \leq \alpha_d$ ,  $\beta vr_d^{(\beta)} \geq \alpha_d vr_d^{(\alpha_d)} \approx 0.404240 > \frac{\lambda_1}{1+\lambda_1} \approx 0.176471$ .  
 Thus,  $h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.12180 < \frac{\alpha_d}{\lambda_1} \approx 0.385675$ .

**Example 8.** Let  $\Gamma$  be a graph of the Johnson scheme  $J(m, d)$  with respect to  $A_1$ . Then  $\Gamma$  is a distance-regular graph with  $\binom{m}{d}$  vertices, valency  $d(m-d)$  and  $d$  diameter. We consider two cases: (a)  $m = 6$ ,  $d = 3$  and (b)  $m = 11$ ,  $d = 5$ . In each case, our upper bound on the Cheeger constant is as follows:

- (a)  $J(6, 3)$  :  $v = 126$ ,  $k_1 = 20$ ,  $\lambda_1 = 9/20$ ,  $\alpha_d = 252/1325$ .  
 And, for  $\beta \leq \alpha_d$ ,  $\beta vr_d^{(\beta)} \geq \alpha_d vr_d^{(\alpha_d)} \approx 0.415036 > \frac{\lambda_1}{1+\lambda_1} \approx 0.310345$ .  
 Thus,  $h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.268058 < \frac{\alpha_d}{\lambda_1} \approx 0.422642$ .
- (b)  $J(11, 5)$  :  $v = 462$ ,  $k_1 = 30$ ,  $\lambda_1 = 11/30$ ,  $\alpha_d = 11088/79091$ .  
 And, for  $\beta \leq \alpha_d$ ,  $\beta vr_d^{(\beta)} \geq \alpha_d vr_d^{(\alpha_d)} \approx 0.40805 > \frac{\lambda_1}{1+\lambda_1} \approx 0.268293$ .  
 Thus,  $h_\Gamma < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.203375 < \frac{\alpha_d}{\lambda_1} \approx 0.382344$ .

**Example 9.** Let  $\Gamma$  be a Taylor graph with intersection array  $(275, 112, 1; 1, 112, 275)$ . Then  $\Gamma$  is a distance-regular graph with vertices 552, valency 275 and 3 diameter. Also,  $\lambda_1 = 4/5$ ,  $\alpha_3 = 3864/12475$ ,  $\alpha_3^{(\alpha_3)} = 0.2265$ . Thus we have an upper bound on the Cheeger constant of  $\Gamma$  as follows:

$$h_\Gamma < \frac{\alpha_3^2}{\alpha_3^{(\alpha_3)}} \approx 0.42357.$$

In Theorem 2, we find an alternative upper bound, which is explicitly computable, by using  $\alpha_d$ , the valencies  $k_j$ , and the basis of nullspace  $\mathcal{N}(L_{sub}^{(\beta)})$ . In Example 10 and Example 11, we compute the upper bound on the Cheeger constant using the alternative expression in Theorem 2 and Remark 6.

**Example 10.** Let  $\Gamma$  be a graph with respect to  $A_1$  of a Johnson scheme  $J(8, 4)$ . Then  $\Gamma$  is a distance-regular graph with 70 vertices and valency 16. Also, the

valencies of  $J(8, 4)$  are 1, 16, 36, 16, 1 and

$$L_{sub}^{(\beta)} = \begin{pmatrix} 1 & 6 - 16(\beta + 1) & 9 & 0 & 0 \\ 0 & 4 & 8 - 16(\beta + 1) & 4 & 0 \\ 0 & 0 & 9 & 6 - 16(\beta + 1) & 1 \\ 0 & 0 & 0 & 16 & -16(\beta + 1) \end{pmatrix}.$$

Since

$$\alpha_d = \frac{1}{-q_1(d)\frac{1}{\lambda_1} - \dots - q_d(d)\frac{1}{\lambda_d}},$$

$(q_1(4), \dots, q_4(4)) = (-7, 20, -28, 14)$  and  $(\lambda_1, \dots, \lambda_4) = (\frac{8}{16}, \frac{14}{16}, \frac{18}{16}, \frac{20}{16})$ , we get  $\alpha_d = 315/1522$ . Let  $\beta = 315/1522$ . By Remark 6, a basis  $(u_0^{(\alpha_d)}, \dots, u_4^{(\alpha_d)})$  for  $\mathcal{N}(L_{sub}^{(\alpha_d)})$  is

$$\left( \frac{10692972602391}{335381132641}, \frac{3108779427}{881422162}, \frac{969476}{579121}, \frac{1837}{1522}, 1 \right).$$

Thus, by Theorem 2, we have

$$h_\Gamma < (315/1522) \left( \frac{1}{70} \sum_{j=0}^4 k_j u_j^{(\alpha_d)} - 1 \right) \approx 0.292388.$$

**Example 11.** Let  $X$  be a set of  $d \times n$  matrices over  $GF(p^t)$  ( $d \leq n$ ). We define the  $i$ -th relation  $R_i$  on  $X$  by  $(x, y) \in R_i$  if and only if  $\text{rank}(x - y) = i$ . Then  $\mathfrak{X} = (X, \{R_i\})$  ( $0 \leq i \leq d$ ) is a  $P$ -polynomial scheme with respect to the ordering  $R_0, R_1, \dots, R_d$ . Let  $p = 2, t = 1, d = 4, n = 5$ . Then  $L_{sub}^{(\beta)}$  is obtained as follows:

$$\begin{pmatrix} 1 & 44 - 465(\beta + 1) & 420 & 0 & 0 \\ 0 & 6 & 123 - 465(\beta + 1) & 336 & 0 \\ 0 & 0 & 28 & 245 - 465(\beta + 1) & 192 \\ 0 & 0 & 0 & 120 & 345 - 465(\beta + 1) \end{pmatrix}.$$

We have  $|X| = v = 1048576$ ,  $k_0 = 1$ ,  $k_1 = 465$ ,  $k_2 = 32550$ ,  $k_3 = 390600$  and  $k_4 = 624960$  by using  $k_i = \frac{k_1 b_1 b_2 \dots b_{i-1}}{c_2 c_3 \dots c_i}$  ( $i = 2, 3, \dots, d$ ).

Let  $\beta = \frac{1}{100}$ . Then by Lemma 4 and Remark 6 (1), we obtain the unique basis of  $\mathcal{N}(L_{sub}^{(\beta)})$  as follows:

$$(u_0^{(\beta)}, u_1^{(\beta)}, u_2^{(\beta)}, u_3^{(\beta)}, u_4^{(\beta)}) = \left( \frac{3921317781669}{358400000}, \frac{486743013}{17920000}, \frac{661683}{448000}, \frac{831}{800}, 1 \right).$$

Thus, by Remark 6, we find

$$\frac{(1048576)\frac{1}{100}}{(1)u_0^{(\beta)} + (465)u_1^{(\beta)} + (32550)u_2^{(\beta)} + (390600)u_3^{(\beta)} + (624960)u_4^{(\beta)} - 1048576} \approx 0.195023.$$

Thus, we have  $\alpha_d < \widetilde{\alpha}_d \approx 0.195023 + 0.01 = 0.205023$ .

Let  $\beta = 0.205023$ . Then, we obtain the unique basis of  $\mathcal{N}(L_{sub}^{(\widetilde{\alpha}_d)})$  as follows:

$$u_0^{(\widetilde{\alpha}_d)} = \frac{227848229494208860060049660207}{5120000000000000000000000000000}, \quad u_1^{(\widetilde{\alpha}_d)} = \frac{2234200208459744136613}{2560000000000000000000000000000},$$

$$u_2^{(\widetilde{\alpha}_d)} = \frac{85453392459301}{6400000000000000}, \quad u_4^{(\widetilde{\alpha}_d)} = \frac{14355713}{8000000}, \quad u_4^{(\widetilde{\alpha}_d)} = 1.$$

Thus, from Theorem 2, we obtain

$$h_\Gamma < (\widetilde{\alpha}_d) \left( \frac{1}{1048576} \sum_{j=0}^4 k_j u_j^{(\widetilde{\alpha}_d)} - 1 \right) \approx 0.305557.$$

*Remark 12.* In general, it is a hard task to compute the Cheeger constant of the graph, and there is not much known about the actual value of the Cheeger constant of a graph. As far as we know, the only known case is the Cheeger constant of the Hamming graph  $H(d, q)$  with  $q$  even, which is  $\frac{q}{2n(q-1)}$ . For instance, we consider two cases  $H(5, 2)$  and  $H(5, 4)$  for comparing our bound with the actual Cheeger constant:

(a)  $H(5, 2) : v = 512, k_1 = 5, \lambda_1 = 2/5, \alpha_d = 24/137, \alpha_d^{(\alpha_d)} \approx 0.123033.$

Thus,  $h_\Gamma = \frac{1}{5} = 0.2 < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.249437.$

(b)  $H(5, 4) : v = 1024, k_1 = 15, \lambda_1 = 4/15, \alpha_d = 16/137, \alpha_d^{(\alpha_d)} \approx 0.083724.$

Thus,  $h_\Gamma = \frac{2}{15} = 0.133333 \dots < \frac{\alpha_d^2}{\alpha_d^{(\alpha_d)}} \approx 0.166291.$

As we can see from these examples, our bound is close to the Cheeger constant, but it is not sharp yet.

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