

TWO-SIDED BEST SIMULTANEOUS APPROXIMATION

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ABSTRACT. Let $C_1(X)$ be a normed linear space over \mathbb{R}^m , and S be an n -dimensional subspace of $C_1(X)$ with spanned by $\{s_1, \dots, s_n\}$. For each ℓ -tuple vectors F in $C_1(X)$, the two-sided best simultaneous approximation problem is

$$\min_{s \in S} \max_{i=1}^{\ell} \{\|f_i - s\|_1\}.$$

A $s \in S$ attaining the above minimum is called a two-sided best simultaneous approximation or a Chebyshev center for $F = \{f_1, \dots, f_\ell\}$ from S . This paper is concerned with algorithm for calculating two-sided best simultaneous approximation, in the case of continuous functions.

1. Introduction

We assume that X is a compact subset of \mathbb{R}^m satisfying $X = \overline{\text{int} X}$, S is an n -dimensional subspace of $C_1(X)$, and μ is any 'admissible' measure on X . For any ℓ -tuple f_1, \dots, f_ℓ in $C_1(X)$, we present a algorithm to the solution of our problem of finding a $s^* \in S$ such that

$$\max_{1 \leq i \leq \ell} \|f_i - s^*\|_1 \leq \max_{1 \leq i \leq \ell} \|f_i - s\|_1$$

for any $s \in S$.

DEFINITION 1.1. Suppose that K is a nonempty subset of a normed linear space W . Given any bounded subset $F \subset W$, define

$$d(F, K) := \inf_{k \in K} \sup_{f \in F} \|f - k\|_1.$$

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An element $k^* \in K$ is said to be a two-sided best simultaneous approximation or a Chebyshev center for the set F from K if

$$d(F, K) = \sup_{f \in F} \|f - k^*\|_1.$$

That mean that a ball of center k^* with radius $d(F, K)$ is the smallest circle containing the set F .

For each positive integer n , define the set

$$\bar{F}_n := \{(\bar{\lambda}_n, \bar{f}_n) : \bar{\lambda}_n = (\lambda_1, \dots, \lambda_n), \bar{f}_n = (f_1, \dots, f_n), f_i \in F,$$

$$\lambda_i \geq 0 \ (i = 1, \dots, n), \ \Sigma_{i=1}^n \lambda_i = 1\}.$$

Then, for any compact set F , \bar{F}_n is a compact set in the product space and we note that for any $k \in K$,

$$\max_{f \in F} \|f - k\|_1 = \max_{\bar{F}_n} \Sigma_{i=1}^n \lambda_i \|f_i - k\|_1.$$

So $d(F, K) = d(\text{con}(F), K)$ where $\text{con}(F)$ denotes the convex hull of F . Furthermore, as a consequence of the above equality we have a remark.

REMARK 1.2. For any compact set F and K is a nonempty subset of a normed linear space W , k^* is a best simultaneous approximation for F from K if and only if k^* is a best simultaneous approximation for $\text{con}(F)$ from K .

In the next theorem, we are able to determine criteria for when a finite dimensional subspace. It is the existence of a best simultaneous approximation. For a closed convex set, this problem is much more difficult.

THEOREM 1.3. [7] *If K is a finite dimensional subspace of a normed linear space W , then for any compact subset $F \subset W$, there exists a best simultaneous approximation from K .*

COROLLARY 1.4. [8] *Suppose that K is a closed convex subset of a finite dimensional subspace of a normed linear space W . For any compact subset $F \subset W$, there exists a best simultaneous approximation from K .*

2. Two-sided best simultaneous approximation

This algorithm is based on a discretization of our problem. For each positive integer m , let $x_1^m, \dots, x_m^m \in X$ and $\delta_1^m, \dots, \delta_m^m$ be strictly positive numbers such that

$$\lim_{m \rightarrow \infty} \sum_{i=1}^m \delta_i^m f(x_i^m) = \int_X f d\mu$$

for each $f \in C_1(X)$ where μ is any 'admissible' measure on X . For convenience only, we assume that $\sum_{i=1}^m \delta_i^m = \mu(X)$ for all m . The points $\{x_i^m\}_{i=1}^m$ become dense in X . Before proving the convergence of the algorithm, we need to pursue some technical facts.

LEMMA 2.1. [4] *There exists a natural number M_1 where $\{x_1^m, \dots, x_m^m\}$ is such that for all $m \geq M_1$ we have $\dim S|_{\{x_1^m, \dots, x_m^m\}} = n$.*

For any f_1, \dots, f_ℓ in $C_1(X)$, let $F = (f_1, \dots, f_\ell)$ and $\|F\|_1 = \|(f_1, \dots, f_\ell)\|_1 = d(F, \{0\})$. By the remark 1.0.2, we denote $\|F\|_1 = \max_{a \in A} \|\sum_{i=1}^\ell a_i f_i\|_1$ where $A = \{a = (a_1, \dots, a_\ell) : \sum_{i=1}^\ell a_i = 1, a_i \geq 0, i = 1, \dots, \ell\}$. Throughout this article, we assume that $F = (f_1, \dots, f_\ell)$ in $C_1(X)$ are given and S is an n -dimensional subspace of $C_1(X)$. We want to consider the problem of approximating these functions simultaneously by elements in S . In other words, we want to find $s^* \in S$ to minimize

$$\|(f_1 - s^*, \dots, f_\ell - s^*)\|_1.$$

For each positive integer m , we denote by

$$\sigma_m := \min_{s \in S} \max_{j=1}^\ell \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s(x_i^m)|.$$

Let $s^m \in S$ be a solution to satisfy the above equality.

We first prove the following lemma.

LEMMA 2.2. *There exists a positive integer M_2 such that the set $\{s^m\}_{m \geq M_2}$ are uniformly bounded.*

Proof. By the lemma 2.0.5, there exists a natural number M_1 , for all $m \geq M_1$, $\dim S|_{\{x_1^m, \dots, x_m^m\}} = n$. Since, for each $f \in C_1(X)$, $\lim_{m \rightarrow \infty} \sum_{i=1}^m \delta_i^m f(x_i^m) = \int_X f d\mu$, and $\|F\|_1 = \|(f_1, \dots, f_\ell)\|_1 = \max_{a \in A} \|\sum_{i=1}^\ell a_i f_i\|_1 = \max_{1 \leq j \leq \ell} \|f_j\|_1$ there exists a natural number M_2 , for all $m \geq M_2$,

$$\max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m)| \leq \|F\|_1 + 1.$$

Let $M = \max\{M_1, M_2\}$. Then, for all $m \geq M$ and any $j \in \{1, \dots, \ell\}$,

$$\sum_{i=1}^m \delta_i^m |s^m(x_i^m)| \leq \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^m(x_i^m)| + \sum_{i=1}^m \delta_i^m |f_j(x_i^m)| \leq 2\|F\|_1 + 2,$$

since

$$\max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^m(x_i^m)| \leq \max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s(x_i^m)|$$

for all $s \in S$, that is

$$\max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^m(x_i^m)| \leq \max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m)|.$$

Hence there exists a constant $C > 0$ such that

$$\sum_{i=1}^m \delta_i^m |s^m(x_i^m)| \leq C$$

for all $m \geq M$.

Let $s^m = \sum_{h=1}^n a_h^m s_h$, where s_1, \dots, s_n is a basis for S . If the $\{s^m\}$ are not uniformly bounded, there exists a subsequence $\{m_k\}$ on which

- (a) $|a_r^{m_k}| = \max\{|a_h^{m_k}| : h = 1, \dots, n\}$,
- (b) $\lim_{k \rightarrow \infty} |a_r^{m_k}| = \infty$.

Let $b_h^{m_k} = a_h^{m_k} / a_r^{m_k}$, $h = 1, \dots, n$. Then the sequence $\{b_h^{m_k}\}$ is bounded, so there exists a subsequence of $\{m_k\}$, again denoted by $\{m_k\}$, on which

$$\lim_{k \rightarrow \infty} b_h^{m_k} = b_h, \quad h = 1, \dots, n.$$

Set $\sum_{h=1}^n b_h s_h = v$. Since $|b_h^{m_k}| \leq 1 = |b_r^{m_k}|$, $h = 1, \dots, n$, $v \neq 0$. Thus

$$\begin{aligned} \sum_{i=1}^{m_k} \delta_i^{m_k} \left| \sum_{h=1}^n b_h^{m_k} s_h(x_i^{m_k}) \right| &= \sum_{i=1}^{m_k} \delta_i^{m_k} \left| \sum_{h=1}^n \frac{a_h^{m_k}}{a_r^{m_k}} s_h(x_i^{m_k}) \right| \\ &= \frac{1}{|a_r^{m_k}|} \sum_{i=1}^{m_k} \delta_i^{m_k} \left| \sum_{h=1}^n a_h^{m_k} s_h(x_i^{m_k}) \right| \\ &= \frac{1}{|a_r^{m_k}|} \sum_{i=1}^{m_k} \delta_i^{m_k} |s^{m_k}(x_i^{m_k})| \\ &\leq \frac{C}{|a_r^{m_k}|} \end{aligned}$$

for all $m_k \geq M$. Let $k \rightarrow \infty$. The right hand side of the inequality tends to zero, but the left hand side of the inequality tends to

$\int_X |\sum_{h=1}^n b_h s_h| d\mu$. Since $\sum_{h=1}^n b_h s_h \neq 0$, it is a contradiction, which implies the desired result. \square

The algorithm is of course given by finding σ^m and each solution s^m at each step and letting m tend to infinity. Next, we prove a main theorem.

THEOREM 2.3. *Let*

$$\sigma_0 = \min_{s \in S} \{ \max_{j=1}^{\ell} \|f_j - s\|_1 \}.$$

Then $\lim_{m \rightarrow \infty} \sigma_m = \sigma_0$. Furthermore, every convergent subsequence of the $\{s^m\}$ converges to a $s^ \in S$ satisfying $\max_{j=1}^{\ell} \|f_j - s^*\|_1 = \sigma_0$.*

Proof. By the theorem 1.0.3, there exists a $s^0 \in S$ satisfy $\max_{1 \leq j \leq \ell} \|f_j - s^0\|_1 = d(F, S) = \sigma_0$. Then

$$\max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^0(x_i^m)| \geq \max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^m(x_i^m)| = \sigma_m.$$

Furthermore

$$\lim_{m \rightarrow \infty} \max_{1 \leq j \leq \ell} \sum_{i=1}^m \delta_i^m |f_j(x_i^m) - s^0(x_i^m)| = \max_{1 \leq j \leq \ell} \|f_j - s^0\|_1 = \sigma_0.$$

Therefore $\overline{\lim_{m \rightarrow \infty} \sigma_m} \leq \sigma_0$. Assume that a subsequence $\{s^{m_k}\}$ satisfy $\lim_{k \rightarrow \infty} s^{m_k} = s^*$.

claim $\max_{1 \leq j \leq \ell} \|f_j - s^*\|_1 = \sigma_0$.

It is clear, by definition, $\max_{1 \leq j \leq \ell} \|f_j - s^*\|_1 \geq \sigma_0$. Since

$$\begin{aligned} \lim_{k \rightarrow \infty} \sigma_{m_k} &= \lim_{k \rightarrow \infty} \max_{1 \leq j \leq \ell} \{ \sum_{i=1}^{m_k} \delta_i^{m_k} |f_j(x_i^{m_k}) - s^{m_k}(x_i^{m_k})| \} \\ &= \max_{1 \leq j \leq \ell} \|f_j - s^*\|_1. \end{aligned}$$

Thus $\max_{1 \leq j \leq \ell} \|f_j - s^*\|_1 \leq \sigma_0$. So $\max_{1 \leq j \leq \ell} \|f_j - s^*\|_1 = \sigma_0$.

By claim, s^* is a best simultaneous approximation to F from S . Thus every subsequence $\{m_k\}$ for which $\{s^{m_k}\}$ converge, $\lim_{k \rightarrow \infty} \sigma_{m_k} = \sigma_0$. Hence, $\lim_{m \rightarrow \infty} \sigma_m = \sigma_0$. We have therefore proven. \square

Note that s^* is unique, then

$$\lim_{m \rightarrow \infty} s^m = s^*$$

and the convergence is uniform.

Since the best approximation problem is an almost totally general form of a linear programming problem. Many different algorithms exist for solving linear programming problems. The main theorem is bases

on discretization to solve a two-sided best simultaneous approximation. Now we will extend the algorithm that involves absolutely no discretization, the study of this problem leads us to the important concepts of gradients and subgradients.

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