

A NOTE ON SYMMETRIC DIFFERENCES OF ORTHOMODULAR LATTICES

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ABSTRACT. There exist two distinct symmetric differences in a non Boolean orthomodular lattice. Let L be an orthomodular lattice. Then L is a Boolean algebra if and only if one symmetric difference is equal to the other. An orthomodular lattice L is Boolean if and only if one of two symmetric differences of L is associative.

1. Introduction

An *orthocomplementation* on a bounded poset P is a *unary operation* $'$ on P which satisfies the following properties: (1) if $x \leq y$ then $y' \leq x'$; (2) $x'' = x$; (3) $x \vee x' = 1$ and $x \wedge x' = 0$. We call a bounded poset P with an orthocomplementation an *orthoposet*. Two elements x, y of an orthoposet are *orthogonal*, written $x \perp y$, in case $x \leq y'$. An *ortholattice* (abbreviated by OL) is an orthoposet which is also a lattice.

An *orthomodular lattice* (abbreviated by OML) is an ortholattice L which satisfies *the orthomodular law*: if $x \leq y$, then $y = x \vee (x' \wedge y)$ [5]. A *Boolean algebra* B is an ortholattice satisfying the *distributive law*: $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$. The *commutator* of a and b of an OML L is denoted by $a*b$, and is defined by $a*b = (a \vee b) \wedge (a \vee b') \wedge (a' \vee b) \wedge (a' \vee b')$.

For elements a, b of an ortholattice, we say a *commutes with* b , in symbols $a \mathbf{C} b$, if $a = (a \wedge b) \vee (a \wedge b')$. If L is an OML, then the relation \mathbf{C} is symmetric [p.22, 5].

One of our most important computational tools is the Foulis-Holland Theorem: *Let a, b, c be elements in an OML L such that one of them commutes with the other two. Then the sublattice generated by $\{a, b, c\}$ is distributive* [p.25, 5].

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A ring is called *Boolean* if all of its elements are *idempotent* (i.e., $a^2 = a$). For any x and y in a Boolean algebra B , the *symmetric difference* of x and y (in symbols $x + y$) is defined by $x + y = (x \wedge y') \vee (x' \wedge y)$. We can prove that the equivalence of a Boolean algebra and a Boolean ring using the symmetric difference in a Boolean algebra [4]. In this paper, we will study symmetric differences of orthomodular lattices.

The symmetric difference of a Boolean algebra has the following properties.

PROPOSITION 1.1. *The symmetric difference of x and y , $x + y$, of a Boolean algebra B has the following properties for all x, y, z and w in B .*

- (1) $x + x = 0$
- (2) $x + 1 = x'$
- (3) $x + y = y + x$
- (4) $x + 0 = x$
- (5) $(x + y) + z = x + (y + z)$
- (6) $(x + y) \wedge z = (x \wedge z) + (y \wedge z)$
- (7) $x \vee y = (x + y) + (x \wedge y)$
- (8) $x + y = x' + y'$
- (9) $x \wedge y' \leq x + y \leq x \vee y$
- (10) $x \perp y \iff x + y = x \vee y$
- (11) $x \leq y \iff x + y = x' \wedge y$
- (12) $(x + y) \perp (x \wedge y)$
- (13) $x + y = (x \vee y) \wedge (x' \vee y')$
- (14) $x = y \iff x + y = 0$
- (15) $x + y = y + z \iff x = z$
- (16) $(x + y) \vee (y + z) = (x \vee y \vee z) \wedge (x \wedge y \wedge z)'$
- (17) $(x + y) \vee (y + z) \vee (z + w) = (x \vee y \vee z \vee w) \wedge (x \wedge y \wedge z \wedge w)'$
- (18) $x \wedge y = (x + y') \wedge x = x + (x \wedge y')$.

PROOF. Each simple calculation is omitted. □

2. Symmetric differences of orthomodular lattices

DEFINITION 2.1. Let L be an OML, and $+$ be a binary operation on L . We call $+$ a *symmetric difference* on L if $+$ satisfies the following properties for all x, y in L :

- (1) $x + x = 0$
- (2) $x + 1 = x'$

$$(3) \ x + y = y + x.$$

It is known that the OML F_2 generated by two distinct elements is isomorphic to $2^4 \times MO2$ [2]. Therefore we have the following theorem.

THEOREM 2.2. *There exist two distinct symmetric differences in a non Boolean OML.*

PROOF. It is sufficient to show that there are only two polynomials on the OML F_2 which satisfy (1), (2), and (3) of Definition 2.1. Indeed, there are exactly three polynomials $(x \wedge y') \vee (x' \wedge y)$, $(x \vee y) \wedge (x' \vee y')$ and $(x \vee y) \wedge (x \vee y') \wedge (x' \vee y) \wedge (x' \vee y')$ in F_2 which satisfy (1) and (3) [pp.81-85, 1]. Finally, only two polynomials $(x \wedge y') \vee (x' \wedge y)$ and $(x \vee y) \wedge (x' \vee y')$ satisfy (2). \square

We will denote $(x' \wedge y) \vee (x \wedge y')$ and $(x \vee y) \wedge (x' \vee y')$ with $x +_1 y$ and $x +_2 y$, respectively. That is,

$$x +_1 y = (x' \wedge y) \vee (x \wedge y')$$

$$x +_2 y = (x \vee y) \wedge (x' \vee y').$$

The two symmetric differences $x +_1 y$ and $x +_2 y$ of x, y in an OML L have the following common properties.

PROPOSITION 2.3. *Two symmetric differences $x +_1 y$ and $x +_2 y$ have the following properties for all x, y in an OML and each $i = 1, 2$.*

- (1) $x +_i y = x' +_i y'$
- (2) $x +_i 0 = x$
- (3) $x \perp y \iff x +_i y = x \vee y$
- (4) $x \leq y \iff x +_i y = x' \wedge y$
- (5) $(x +_i y) \perp (x \wedge y)$
- (6) $(x \wedge y') \leq (x +_i y) \leq (x \vee y)$
- (7) $x +_i y = x' +_i y'$
- (8) $x +_i (x \wedge y') = x \wedge (x' \vee y)$

PROOF. Each simple calculation is omitted. \square

REMARK. For x, y in an ortholattice L , if we define $x +_1 y$ and $x +_2 y$ to be the same as in an OML, then $x +_1 y$ and $x +_2 y$ have all properties in Definition 2.1 and Proposition 2.3.

We have the following two properties on the symmetric differences of orthomodular lattices.

THEOREM 2.4. *Let L be an OML and $+$ be a symmetric difference on L . Then*

(1) $x+y = x+{}_1y$ if and only if $x+y = (x\wedge y')+(x'\wedge y) \quad \forall x, y \in L$
and

(2) $x+y = x+{}_2y$ if and only if $x\vee y = (x+y)+(x\wedge y) \quad \forall x, y \in L$.

PROOF. (1) The conclusion follows since $(x\wedge y')+(x'\wedge y) = (x\wedge y')\vee(x'\wedge y)$ by (3) of Proposition 2.3.

(2) Assume that $x+y = x+{}_2y = (x\vee y)\wedge(x'\vee y') \quad \forall x, y \in L$.

Then

$$\begin{aligned}
 (x+y)+(x\wedge y) &= (x+{}_2y)+{}_2(x\wedge y) \\
 &= (((x\vee y)\wedge(x'\vee y'))\vee(x\wedge y)) \\
 &\quad \wedge(((x'\wedge y')\vee(x\wedge y))\vee(x'\vee y')) \\
 &= ((x\vee y)\wedge(x'\vee y'))\vee(x\wedge y) \\
 &= ((x\vee y)\vee(x\wedge y))\wedge((x'\vee y')\vee(x\wedge y)) \\
 &= (x\vee y)\vee(x\wedge y) \\
 &= x\vee y.
 \end{aligned}$$

Conversely, assume that $x\vee y = (x+y)+(x\wedge y) \quad \forall x, y \in L$. Then

$$\begin{aligned}
 &x+{}_2y \\
 &= (x\vee y)\wedge(x'\vee y') \\
 &= ((x+y)+(x\wedge y))\wedge((x'+y')+(x'\wedge y')) \\
 &= ((x+y)\vee(x\wedge y))\wedge((x'+y')\vee(x'\wedge y')) \quad \text{since } (x+y)\perp(x\wedge y) \\
 &= ((x+y)\vee(x\wedge y))\wedge((x+y)\vee(x'\wedge y')) \quad \text{since } x+y = x'+y' \\
 &= (((x+y)\vee(x\wedge y))\wedge(x+y))\vee(((x+y)\vee(x\wedge y))\wedge(x'\wedge y')) \\
 &= (x+y)\vee(((x+y)+(x\wedge y))\wedge(x'\wedge y')) \\
 &= (x+y)\vee((x\vee y)\wedge(x'\wedge y')) \\
 &= x+y.
 \end{aligned}$$

□

PROPOSITION 2.5. *Let L be an OML and x, y in L . Then $(x+{}_1y)' = x'+{}_2y$ and $x+{}_1y \leq x+{}_2y$.*

PROOF. The conclusions follow by two definitions of $x+{}_1y$ and $x+{}_2y$. □

We know that $x * y = (x+{}_1y)' \wedge (x'+{}_1y)'$ and $x * y = (x+{}_2y) \wedge (x'+{}_2y)$. Now, we are ready to prove the following theorem.

THEOREM 2.6. *Let L be an OML and x, y in L . Then the following conditions are equivalent.*

- (1) $x \mathbf{C} y$
- (2) $x +_1 y = x +_2 y$
- (3) $x +_2 y = (x \wedge y') +_2 (x' \wedge y)$
- (4) $x +_1 y = (x \vee y) +_1 (x \wedge y)$
- (5) $x = (x \wedge y) +_i (x \wedge y')$ for each $i = 1, 2$.

PROOF. (1) \iff (2)

Assume that $x \mathbf{C} y$. Then

$$\begin{aligned} x +_1 y &= (x' \wedge y) \vee (x \wedge y') \\ &= ((x' \wedge y) \vee x) \wedge ((x' \wedge y) \vee y') \\ &= ((x' \vee x) \wedge (y \vee x)) \wedge ((x' \vee y') \wedge (y \vee y')) \\ &= (y \vee x) \wedge (x' \vee y') \\ &= x +_2 y. \end{aligned}$$

Conversely, assume that $x +_1 y = x +_2 y$. Then $(x \wedge y') \vee (x' \wedge y) = (x \vee y) \wedge (x' \vee y')$. Thus

$$\begin{aligned} x * y &= ((x \vee y) \wedge (x' \vee y')) \wedge (x' \vee y) \wedge (x \vee y') \\ &= (((x \wedge y') \vee (x' \wedge y)) \wedge (x' \vee y)) \wedge (x \vee y') \\ &= (((x \wedge y') \wedge (x' \vee y)) \vee ((x' \wedge y) \wedge (x' \vee y))) \wedge (x \vee y') \\ &= (x \wedge y') \wedge (x' \vee y) \wedge (x \vee y') \\ &= (x' \wedge y) \wedge (x \vee y') \\ &= 0. \end{aligned}$$

Therefore $x \mathbf{C} y$ since $x * y = 0$ if and only if $x \mathbf{C} y$ [3].

(2) \iff (3)

$$x +_2 y = x +_1 y = (x \wedge y') \vee (x' \wedge y) = (x \wedge y') +_2 (x' \wedge y).$$

(2) \iff (4)

$$\begin{aligned} x +_1 y &= x +_2 y \\ &= (x \vee y) \wedge (x' \vee y') \\ &= ((x' \wedge y') \vee (x \wedge y))' \\ &= ((x' \wedge y') +_2 (x \wedge y))' && \text{by (3) of Proposition 2.3} \\ &= (x' \wedge y') +_1 (x \wedge y) && \text{by Proposition 2.5} \\ &= (x \vee y) +_1 (x \wedge y) \end{aligned}$$

(1) \iff (5)

$$x \mathbf{C} y \iff x = (x \wedge y) \vee (x \wedge y') \iff x = (x \wedge y) +_i (x \wedge y') \quad \square$$

$\forall i = 1, 2$.

The following three corollaries follow.

COROLLARY 2.7. *Let L be an OML. Then L is a Boolean algebra if and only if $x +_1 y = x +_2 y \quad \forall x, y \in L$.*

PROOF. Suppose that L is a Boolean. Then $x +_1 y = x +_2 y \quad \forall x, y \in L$ by (13) of Proposition 1.1.

Conversely, if $x +_1 y = x +_2 y \quad \forall x, y \in L$, then $x \mathbf{C} y \quad \forall x, y \in L$ by (1) and (2) of Theorem 2.6. This completes the proof. \square

COROLLARY 2.8. *Let L be an OML and $x, y \in L$. Then $x \leq y$ if and only if $x +_1 y = x +_2 y = x' \wedge y$.*

PROOF. Assume that $x \leq y$. Then $x +_1 y = x +_2 y = x' \wedge y$ by (2) of Proposition 2.3.

Conversely, assume that $x +_1 y = x +_2 y$. Then $x \mathbf{C} y$ by Theorem 2.6. Thus $y \mathbf{C} x$ since $y \mathbf{C} x$ if and only if $x \mathbf{C} y$ [5]. Therefore

$$\begin{aligned} y &= (x \wedge y) \vee (x' \wedge y) \\ &= (x \wedge y) \vee ((x \vee y) \wedge (x' \vee y')) \quad \text{for } x +_1 y = x' \wedge y \\ &= ((x \wedge y) \vee (x \vee y)) \wedge ((x \wedge y) \vee (x' \vee y')) \\ &= x \vee y. \end{aligned}$$

Thus $y \geq x$. \square

COROLLARY 2.9. *Let L be an OML and $x, y \in L$. Then $x \perp y$ if and only if $x +_1 y = x +_2 y = x \vee y$.*

PROOF.

$$\begin{aligned} x \perp y &\iff x \leq y' \\ &\iff x +_1 y' = x +_2 y' = x' \wedge y' && \text{by Corollary 2.8} \\ &\iff (x +_1 y')' = (x +_2 y')' = x \vee y \\ &\iff x' +_2 y' = x' +_1 y' = x \vee y && \text{by Proposition 2.5} \\ &\iff x +_2 y = x +_1 y = x \vee y && \text{by (1) of Proposition 2.3.} \end{aligned}$$

\square

We have the following equivalent condition for an ortholattice to be an OML.

THEOREM 2.10. *Let L be an ortholattice. Then L is an OML if and only if $x \vee y = (x +_2 y) +_2 (x \wedge y) \quad \forall x, y \in L$.*

PROOF. Assume that L is an OML. Then

$$\begin{aligned} x \vee y &= (x \wedge y) \vee ((x \wedge y)' \wedge (x \vee y)) \\ &= (x \wedge y) \vee ((x' \vee y') \wedge (x \vee y)) \\ &= (x \wedge y) \vee (x +_2 y) \\ &= (x \wedge y) +_2 (x +_2 y) \\ &= (x +_2 y) +_2 (x \wedge y). \end{aligned}$$

Conversely, assume $x \vee y = (x +_2 y) +_2 (x \wedge y)$ and $x \leq y$. Then

$$\begin{aligned} y &= x \vee y \\ &= (x +_2 y) +_2 (x \wedge y) \\ &= (x' \wedge y) +_2 x \quad \text{by (4) and (3) of Proposition 2.3 and the Remark} \\ &= (x' \wedge y) \vee x \quad \text{since } (x' \wedge y) \perp x. \end{aligned}$$

Thus $x \leq y$ implies $y = (x' \wedge y) \vee x$. □

The associativity of symmetric differences in an OML has the following property [1].

THEOREM 2.11. *An OML L is Boolean if and only if one of two symmetric differences of L is associative.*

PROOF. We know that if an OML L is Boolean then two symmetric differences in L are associative by (5) and (13) of Proposition 1.1.

Conversely, suppose that $+_1$ is associative in L . Then $x = x +_1 0 = x +_1 (y + y) = (x +_1 y) +_1 y = (((x \wedge y') \vee (x' \wedge y)) \wedge y') \vee (((x' \vee y) \wedge (x \vee y')) \wedge y) = (x \wedge y') \vee (y \wedge (x \vee y'))$ since $(x \wedge y') \mathbf{C} y$ and $(x' \wedge y) \mathbf{C} y$. Thus $x \wedge y = ((x \wedge y') \vee (y \wedge (x \vee y'))) \wedge y = (x \wedge y' \wedge y) \vee (y \wedge (x \vee y') \wedge y) = y \wedge (x \vee y')$. Hence $x = (x \wedge y') \vee (y \wedge (x \vee y')) = (x \wedge y') \vee (x \wedge y)$. This means that $x \mathbf{C} y$. Therefore L is Boolean. Similarly, we can show that if $+_2$ is associative, then $x \mathbf{C} y$. □

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