

ON THE M -SOLUTION OF THE FIRST KIND EQUATIONS

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1. Introduction

Let K be a bounded linear operator from Hilbert space H_1 into Hilbert space H_2 . When numerically solving the first kind equation $Kf = g$, one usually picks n orthonormal functions $\phi_1, \phi_2, \dots, \phi_n$ in H_1 which depend on the numerical method and on the problem, see Varah [12] for more details. Then one finds the unique minimum norm element $f_M \in M$ that satisfies $\|Kf_M - g\| = \inf\{\|Kf - g\| : f \in M\}$, where M is the linear span of $\phi_1, \phi_2, \dots, \phi_n$. Such a solution $f_M \in M$ is called the M -solution of $Kf = g$. Some methods for finding the M -solution of $Kf = g$ were proposed by Banks [2] and Marti [9, 10]. See [5, 6, 8] for convergence results comparing the M -solution of $Kf = g$ with f_0 , the least squares solution of minimum norm (LSSMN) of $Kf = g$.

Throughout this paper, it is assumed that H_1 and H_2 are Hilbert spaces and M is a finite dimensional subspace of H_1 . Let $B(H_1, H_2)$ denote the space of all bounded linear operators from H_1 to H_2 , and for $K \in B(H_1, H_2)$ let K_N denote the restriction of K to a subspace N of H_1 , i.e., $K_N = K|_N : N \rightarrow H_2$ is an operator such that $K_N x = Kx$ for all $x \in N$. For $K \in B(H_1, H_2)$, K^* denotes the adjoint operator in $B(H_2, H_1)$ and $K^\dagger : R(K) + R(K)^\perp \rightarrow H_1$ denotes the Moore-Penrose generalized inverse of K [4], where $R(K)$ refers to the range space of K and $R(K)^\perp$ refers to the orthogonal complement of $R(K)$. Let (\cdot, \cdot) denote an inner product and $\|\cdot\|$ denote a corresponding norm, and let $\langle \phi_1, \phi_2, \dots, \phi_n \rangle$ denote the closed linear span of $\phi_1, \phi_2, \dots, \phi_n$.

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It is well-known that the problem of solving the first kind equation $Kf = g$ is ill-posed in that arbitrarily small perturbations in the data g or K may cause an arbitrarily large perturbation in the solution f [1, 11]. A typical example of such an ill-posed problem is the Fredholm integral equation of the first kind. The purpose of this paper is to develop a new method producing the M -solution to $Kf = g$ which is well-posed under perturbations in both K and g .

This paper is organized as follows. First, it is shown that the M -solution of $Kf = g$ is *well-posed* under perturbation in g , but is *not well-posed* under perturbations in both K and g . Then, a generalized Gram-Schmidt (GGS) method for finding the M -solution of $Kf = g$ is proposed, and it is shown that a *modified version* of the GGS method, called the $\text{GGS}(\delta)$, provides the M -solution to $Kf = g$ which is well-posed under perturbations in both K and g . Lastly, numerical implementation of the $\text{GGS}(\delta)$ method is described in detail, and then some numerical results obtained by applying this method to the first kind Fredholm integral equations are reported.

2. Numerical method and its well-posedness

We begin this section by giving a precise definition for the well-posedness of the M -solution to $Kf = g$.

DEFINITION 2.1. Let $K \in B(H_1, H_2)$. The M -solution of the first kind equation $Kf = g$ is *well-posed under perturbations in both K and g* if for each $\epsilon > 0$, there exist $\delta_1 > 0$ and $\delta_2 > 0$ such that for all $\hat{g} \in H_2$ and $\hat{K} \in B(H_1, H_2)$ with $\|g - \hat{g}\| < \delta_1$ and $\|K - \hat{K}\| < \delta_2$, $\|f_M - \hat{f}_M\| < \epsilon$, where f_M and \hat{f}_M are the M -solutions of $Kf = g$ and $\hat{K}\hat{f} = \hat{g}$, respectively. That is, if the M -solution of $Kf = g$ depends continuously on the data K and g , then the M -solution is said to be well-posed under perturbations in both K and g .

In the above definition, if K is fixed and only g is allowed to vary, then the M -solution of $Kf = g$ is said to be *well-posed under perturbation in g* . Similarly, the well-posedness of the M -solution to $Kf = g$ under perturbation in K can be defined.

THEOREM 2.2. *The M -solution of $Kf = g$ is well-posed under perturbation in g , but it is not well-posed under perturbation in K , where $K \in B(H_1, H_2)$.*

Proof. Since M is finite dimensional, K_M has closed range. Thus K_M^\dagger is a bounded linear operator. For an element \hat{g} in H_2 , let f_M and \hat{f}_M denote the M -solutions of $Kf = g$ and $Kf = \hat{g}$, respectively. Then

$$\|\hat{f}_M - f_M\| = \|K_M^\dagger \hat{g} - K_M^\dagger g\| \leq \|K_M^\dagger\| \|\hat{g} - g\|.$$

Therefore, the M -solution is well-posed under perturbation in g .

To see that the M -solution is not well-posed under perturbation in K , let $H_1 = H_2$, ϕ_1 and ϕ_2 be orthonormal vectors in H_1 , $M = \langle \phi_1, \phi_2 \rangle$, and let P_i be the orthogonal projection of H_1 onto $\langle \phi_i \rangle$ for $i = 1, 2$. Then define $K = P_1$ and $K^{(n)} = P_1 + \frac{1}{n}P_2$, where n is a natural number. For $g = \alpha_1\phi_1 + \alpha_2\phi_2$ with $\alpha_2 \neq 0$, let f_M and $f_M^{(n)}$ denote the M -solutions of $Kf = g$ and $K^{(n)}f = g$, respectively. Then, we have

$$\begin{aligned} \|f_M^{(n)} - f_M\| &= \|(K_M^{(n)})^\dagger g - K_M^\dagger g\| \\ &= \|(\alpha_1\phi_1 + n\alpha_2\phi_2) - \alpha_1\phi_1\| \\ &= n|\alpha_2|. \end{aligned}$$

Since $\|K - K^{(n)}\| = \frac{1}{n}$, $K^{(n)}$ converges to K as $n \rightarrow \infty$. However, from the above equality, $f_M^{(n)}$ does not converge to f_M as $n \rightarrow \infty$. This completes the proof.

The above theorem implies that the M -solution of the first kind equation $Kf = g$ is not well-posed under perturbations in both K and g . We now introduce a generalized Gram-Schmidt (GGs) method for calculating the M -solution of $Kf = g$ which is based on the Gram-Schmidt orthogonalization procedure, where $K \in B(H_1, H_2)$. From now on, I denotes an index set $\{1, 2, \dots, n\}$. Let $\{\phi_i\}_{i \in I}$ be an orthonormal basis for M . The GGS method on K relative to M consists of two steps. The **first step** of the GGS method is as follows:

$$\begin{aligned} &\text{Compute } y_1 = K\phi_1 \\ &\text{Let } \psi_1 = \begin{cases} \frac{y_1}{\|y_1\|} & \text{if } y_1 \neq 0 \\ 0 & \text{if } y_1 = 0 \end{cases} \end{aligned}$$

For $i = 2, \dots, n$

$$y_i = K\phi_i - \sum_{j=1}^{i-1} (K\phi_i, \psi_j)\psi_j$$

$$\text{Let } \psi_i = \begin{cases} \frac{y_i}{\|y_i\|} & \text{if } y_i \neq 0 \\ 0 & \text{if } y_i = 0 \end{cases}$$

From this step, it can be easily seen that $\|\psi_i\| = \begin{cases} 1 & \text{if } y_i \neq 0 \\ 0 & \text{if } y_i = 0 \end{cases}$, $(\psi_i, \psi_j) = 0$ for $i \neq j$, and $\langle \psi_1, \dots, \psi_i \rangle = \langle K\phi_1, \dots, K\phi_i \rangle$ for each $i \in I$.

To describe the second step of the GGS, we need the following theorem which underlies this step. Define the $n \times n$ matrix $A = (a_{ij})$, where $a_{ij} = (K\phi_j, \psi_i)$, and the vector $b = (b_i) \in R^n$, where $b_i = (g, \psi_i)$. Note that $a_{ij} = 0$ for $i > j$ and $a_{ii} = (K\phi_i, \psi_i) = (y_i, \psi_i) = \|y_i\|$. Recall that the elements of M can be written as $\sum_{j=1}^n \beta_j \phi_j$.

THEOREM 2.3. $\sum_{j=1}^n \beta_j \phi_j$ is the M -solution of $Kf = g$ if and only if the vector $z = (\beta_j) \in R^n$ is the minimum 2-norm solution of $Ax = b$, i.e., $z = A^\dagger b$, where A^\dagger is the Moore-Penrose generalized inverse of the matrix A [7].

Proof. Let P denote the orthogonal projection of H_2 onto $R(K_M)$. Suppose that $\tilde{f} = \sum_{j=1}^n \beta_j \phi_j$ is a solution to $Kf = Pg$. Since $(g - Pg) \in \langle K\phi_1, \dots, K\phi_n \rangle^\perp$, by the construction of $\{\psi_i\}$ $(g - Pg) \in \langle \psi_1, \dots, \psi_n \rangle^\perp$.

Thus, for each $i \in I$

$$0 = (g - Pg, \psi_i) = (g - K\tilde{f}, \psi_i) = (g, \psi_i) - (K\tilde{f}, \psi_i).$$

It follows that for each $i \in I$

$$b_i = (g, \psi_i) = \sum_{j=1}^n \beta_j (K\phi_j, \psi_i) = \sum_{j=1}^n a_{ij} \beta_j.$$

Hence the vector $z = (\beta_j)$ is a solution of $Ax = b$.

Suppose that $z = (\beta_j)$ is a solution to $Ax = b$. Let $\tilde{f} = \sum_{j=1}^n \beta_j \phi_j$. Since $Az = b$, for each $i \in I$ $b_i = \sum_{j=1}^n a_{ij} \beta_j$ and so $(g, \psi_i) = (K\tilde{f}, \psi_i)$.

Thus, $(g - K\tilde{f}, \psi_i) = 0$ for each $i \in I$. This implies that $(g - K\tilde{f}, K\phi_i) = 0$ for each $i \in I$ and hence $(g - K\tilde{f}) \in R(K_M)^\perp$. Therefore, $\tilde{f} = \sum_{j=1}^n \beta_j \phi_j$ is a solution to $Kf = Pg$.

Note that for $\tilde{f} = \sum_{j=1}^n \beta_j \phi_j$ and $z = (\beta_j)$, $\|\tilde{f}\|^2 = \sum_{j=1}^n |\beta_j|^2 = \|z\|_2^2$. Since the M -solution of $Kf = g$ is the minimum norm element in M satisfying $Kf = Pg$, from the above arguments the theorem holds.

From this theorem, we can see that the **second step** of the GGS method is to find the minimum norm solution $z = (\beta_j)$ of $Ax = b$ and then form the M -solution $\sum_{j=1}^n \beta_j \phi_j$. Notice that the GGS method does not require the linear independence of $K\phi_1, \dots, K\phi_n$, whereas the Banks method [2] does require the linear independence of $K\phi_1, \dots, K\phi_n$.

When one implements the GGS method on a computer with finite precision arithmetic, it is very unlikely to have an exact zero for y_i . This fact is a main motivation for developing a *modified version* of the GGS method which is called the $\text{GGS}(\delta)$ method from now on, where $\delta > 0$ is a fixed small constant. The choice of a suitable constant δ is discussed later. The $\text{GGS}(\delta)$ method on K relative to M also consists of two steps. The **first step** of the $\text{GGS}(\delta)$ method is as follows:

Compute $y_1 = K\phi_1$

$$\text{Let } \psi_1 = \begin{cases} \frac{y_1}{\|y_1\|} & \text{if } \|y_1\| \geq \delta \\ 0 & \text{if } \|y_1\| < \delta \end{cases}$$

For $i = 2, \dots, n$

$$y_i = K\phi_i - \sum_{j=1}^{i-1} (K\phi_i, \psi_j) \psi_j$$

$$\text{Let } \psi_i = \begin{cases} \frac{y_i}{\|y_i\|} & \text{if } \|y_i\| \geq \delta \\ 0 & \text{if } \|y_i\| < \delta \end{cases}$$

From the above step, it is easy to see that $\|\psi_i\| = \begin{cases} 1 & \text{if } \|y_i\| \geq \delta \\ 0 & \text{if } \|y_i\| < \delta \end{cases}$, $(\psi_i, \psi_j) = 0$ for $i \neq j$, and $\langle \psi_1, \dots, \psi_i \rangle \subset \langle K\phi_1, \dots, K\phi_i \rangle$ for each $i \in I$ which is not the same as the corresponding one of the GGS. The **second step** of the $\text{GGS}(\delta)$ is to compute $A^\dagger b$ and then form the *approximate*

M -solution $\sum_{j=1}^n (A^\dagger b)_j \phi_j$ to $Kf = g$, where A and b are defined the same as in the GGS method. From now on, μ refers to the positive real number $\min\{\|y_i\| : \|y_i\| > 0, i \in I\}$.

If δ is chosen so that $0 < \delta < \mu$, then $\|y_i\| < \delta$ is equivalent to $y_i = 0$, and hence the GGS(δ) method is mathematically the same as the GGS method. For this δ , it is clear that $\langle \psi_1, \dots, \psi_i \rangle = \langle K\phi_1, \dots, K\phi_i \rangle$ for each $i \in I$. Therefore, if $0 < \delta < \mu$, then by Theorem 2.3 the approximate M -solution $\sum_{j=1}^n (A^\dagger b)_j \phi_j$ obtained from the GGS(δ) becomes the (exact) M -solution to $Kf = g$ (it is of course assumed that all operations are carried out using infinite precision arithmetic). When using finite precision arithmetic, choosing a constant δ a priori so that $\sum_{j=1}^n (A^\dagger b)_j \phi_j$ is the M -solution to $Kf = g$ is not an easy problem, but numerical experiments show that such a δ may be chosen as a real number which is slightly greater than the unit roundoff for the specific computer to be used (see Section 3).

THEOREM 2.4. *Let $K \in B(H_1, H_2)$. If $\delta > 0$ is a fixed constant chosen sufficiently small, then the GGS(δ) method provides the M -solution of $Kf = g$ which is well-posed under perturbations in both K and g . Actually, the assumption of this theorem means that δ may be chosen to be any real number such that $0 < \delta < \mu$.*

For the proof of this theorem, we will use the following notations. Let $K^\epsilon \in B(H_1, H_2)$ and $g^\epsilon \in H_2$ such that $\|K - K^\epsilon\| \leq \epsilon$ and $\|g - g^\epsilon\| \leq \epsilon$. Let $\{\phi_i\}_{i \in I}$ be an orthonormal basis for $M \subset H_1$. Let y_i 's and ψ_i 's be generated from the GGS(δ) on K relative M , and let y_i^ϵ 's and ψ_i^ϵ 's be generated from the GGS(δ) on K^ϵ relative to M . The symbols a_{ij} , b_i , A , and b are defined the same as before. Similarly, we define $a_{ij}^\epsilon = (K^\epsilon \phi_j, \psi_i^\epsilon)$, $b_i^\epsilon = (g^\epsilon, \psi_i^\epsilon)$, $A^\epsilon = (a_{ij}^\epsilon)$, and $b^\epsilon = (b_i^\epsilon)$. From the construction of the GGS(δ), the following Lemma 2.5 can be easily shown.

LEMMA 2.5. *If $\delta > 0$, then $a_{ii} = \begin{cases} \|y_i\| & \text{if } \|y_i\| \geq \delta \\ 0 & \text{if } \|y_i\| < \delta \end{cases}$ for each $i \in I$. Moreover, if $0 < \delta < \mu$, then $a_{ij} = 0$ for $i > j$.*

LEMMA 2.6. *Suppose $\delta > 0$ is chosen so that $\delta \neq \|y_i\|$ for all $i \in I$. Then, for each $i \in I$, $\lim_{\epsilon \rightarrow 0} y_i^\epsilon = y_i$ and $\lim_{\epsilon \rightarrow 0} \psi_i^\epsilon = \psi_i$.*

Proof. We proceed by mathematical induction on the index i . For

$i = 1$, $y_1 = K\phi_1$ and $y_1^\epsilon = K^\epsilon\phi_1$. Then

$$(1) \quad \|y_1^\epsilon - y_1\| = \|K^\epsilon\phi_1 - K\phi_1\| \leq \|K^\epsilon - K\| \leq \epsilon.$$

From the inequality (1), $\lim_{\epsilon \rightarrow 0} y_1^\epsilon = y_1$. If $\|y_1\| < \delta$, then $\|y_1^\epsilon\| < \delta$ for sufficiently small ϵ and so $\psi_1^\epsilon = \psi_1 = 0$. Clearly, $\lim_{\epsilon \rightarrow 0} \psi_1^\epsilon = \psi_1$ for $\|y_1\| < \delta$. Since $\delta \neq \|y_1\|$, we consider the case $\|y_1\| > \delta$. Since $\lim_{\epsilon \rightarrow 0} y_1^\epsilon = y_1$, $\|y_1^\epsilon\| > \delta$ for sufficiently small ϵ and hence $\psi_1^\epsilon = \frac{y_1^\epsilon}{\|y_1^\epsilon\|}$. Note that $\|\psi_1 - \psi_1^\epsilon\| \leq \frac{2}{\|y_1\|} \|y_1 - y_1^\epsilon\|$. Therefore, $\lim_{\epsilon \rightarrow 0} \psi_1^\epsilon = \psi_1$ for the case of $\|y_1\| > \delta$. Next, we suppose that $\lim_{\epsilon \rightarrow 0} y_i^\epsilon = y_i$ and $\lim_{\epsilon \rightarrow 0} \psi_i^\epsilon = \psi_i$ for all $i < k$. Then,

$$(2) \quad \begin{aligned} \|y_k - y_k^\epsilon\| &= \|K\phi_k - \sum_{j=1}^{k-1} (K\phi_k, \psi_j)\psi_j - K^\epsilon\phi_k + \sum_{j=1}^{k-1} (K^\epsilon\phi_k, \psi_j^\epsilon)\psi_j^\epsilon\| \\ &\leq \|K^\epsilon - K\| + \sum_{j=1}^{k-1} \|(K^\epsilon\phi_k, \psi_j^\epsilon)\psi_j^\epsilon - (K\phi_k, \psi_j)\psi_j\|. \end{aligned}$$

On the other hand, for each $j < k$,

$$(3) \quad \begin{aligned} &\|(K^\epsilon\phi_k, \psi_j^\epsilon)\psi_j^\epsilon - (K\phi_k, \psi_j)\psi_j\| \\ &\leq |(K^\epsilon\phi_k, \psi_j^\epsilon)| \|\psi_j^\epsilon - \psi_j\| + |(K^\epsilon\phi_k, \psi_j^\epsilon) - (K\phi_k, \psi_j)| \|\psi_j\| \\ &\leq \|K^\epsilon\| \|\psi_j^\epsilon - \psi_j\| + |(\phi_k, (K^\epsilon)^*\psi_j^\epsilon - K^*\psi_j)| \\ &\leq \|K^\epsilon\| \|\psi_j^\epsilon - \psi_j\| + \|(K^\epsilon)^*\psi_j^\epsilon - K^*\psi_j\| \\ &\leq \|K^\epsilon\| \|\psi_j^\epsilon - \psi_j\| + \|(K^\epsilon)^*\| \|\psi_j^\epsilon - \psi_j\| + \|(K^\epsilon - K)^*\| \\ &= 2\|K^\epsilon\| \|\psi_j^\epsilon - \psi_j\| + \|K^\epsilon - K\|. \end{aligned}$$

Combining the inequalities (2) and (3), one obtains

$$(4) \quad \begin{aligned} \|y_k - y_k^\epsilon\| &\leq k\|K^\epsilon - K\| + 2\|K^\epsilon\| \sum_{j=1}^{k-1} \|\psi_j^\epsilon - \psi_j\| \\ &\leq k\epsilon + 2(\|K\| + \epsilon) \sum_{j=1}^{k-1} \|\psi_j^\epsilon - \psi_j\|. \end{aligned}$$

By the induction hypothesis and the inequality (4), $\lim_{\epsilon \rightarrow 0} y_k^\epsilon = y_k$. If $\|y_k\| < \delta$, then $\|y_k^\epsilon\| < \delta$ for sufficiently small ϵ and hence $\psi_k = \psi_k^\epsilon = 0$. If $\|y_k\| > \delta$, then $\|y_k^\epsilon\| > \delta$ for sufficiently small ϵ and hence $\psi_k^\epsilon = \frac{y_k^\epsilon}{\|y_k^\epsilon\|}$. Since $\lim_{\epsilon \rightarrow 0} y_k^\epsilon = y_k$, $\lim_{\epsilon \rightarrow 0} \psi_k^\epsilon = \psi_k$ for the case of $\|y_k\| > \delta$.

LEMMA 2.7. *If $\delta > 0$ is chosen so that $\delta \neq \|y_i\|$ for all $i \in I$, then $\lim_{\epsilon \rightarrow 0} \|A^\epsilon - A\|_2 = 0$.*

Proof. Notice that $A = (a_{ij})$ and $A^\epsilon = (a_{ij}^\epsilon)$, where $a_{ij} = (K\phi_j, \psi_i)$ and $a_{ij}^\epsilon = (K^\epsilon\phi_j, \psi_i^\epsilon)$. Then,

$$\begin{aligned}
 |a_{ij} - a_{ij}^\epsilon| &= |(K\phi_j, \psi_i) - (K^\epsilon\phi_j, \psi_i^\epsilon)| \\
 &= |(\phi_j, K^*\psi_i - (K^\epsilon)^*\psi_i^\epsilon)| \\
 (5) \qquad &\leq \|K^*\psi_i - (K^\epsilon)^*\psi_i^\epsilon\| \\
 &\leq \|K - K^\epsilon\| + \|K^\epsilon\| \|\psi_i - \psi_i^\epsilon\| \\
 &\leq \epsilon + (\|K\| + \epsilon)\|\psi_i - \psi_i^\epsilon\|.
 \end{aligned}$$

Since $\|A - A^\epsilon\|_\infty = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij} - a_{ij}^\epsilon|$, from the inequality (5)

$$\|A - A^\epsilon\|_\infty \leq n \max_{1 \leq i \leq n} (\epsilon + (\|K\| + \epsilon)\|\psi_i - \psi_i^\epsilon\|).$$

Since $\lim_{\epsilon \rightarrow 0} \psi_i^\epsilon = \psi_i$ from Lemma 2.6 and $n < \infty$, $\lim_{\epsilon \rightarrow 0} \|A - A^\epsilon\|_\infty = 0$. Note that $\|A - A^\epsilon\|_2 \leq \sqrt{n}\|A - A^\epsilon\|_\infty$. Therefore, $\lim_{\epsilon \rightarrow 0} \|A - A^\epsilon\|_2 = 0$.

LEMMA 2.8. *If B is an $n \times n$ matrix and $\|E\|_2 < \frac{1}{\|B^\dagger\|_2}$, then $\text{rank}(B) \leq \text{rank}(B + E)$. (See [7, Theorem 8.15] for the proof.)*

LEMMA 2.9. *If $0 < \delta < \mu$, $\text{rank}(A) = \text{rank}(A^\epsilon)$ for sufficiently small ϵ .*

Proof. By Lemma 2.7, $\lim_{\epsilon \rightarrow 0} \|A - A^\epsilon\|_2 = 0$. Therefore, from Lemma 2.8, $\text{rank}(A) \leq \text{rank}(A^\epsilon)$ for sufficiently small ϵ . Hence, if $\text{rank}(A) = n$, then $\text{rank}(A^\epsilon) = n$ for sufficiently small ϵ .

Suppose that $\text{rank}(A) < n$. It is easy to show that $a_{ii} = (K\phi_i, \psi_i) = (y_i, \psi_i) = 0$ if and only if $i \in I_0$, where $I_0 = \{i \in I \mid \|y_i\| < \delta\}$. Since

$0 < \delta < \mu$, $\text{rank}(A)$ equals the cardinality of $I - I_0$. From Lemma 2.6, we have $\lim_{\epsilon \rightarrow 0} y_i^\epsilon = y_i$ for all $i \in I$. Hence, for each $i \in I_0$, $\|y_i^\epsilon\| < \delta$ when ϵ is sufficiently small. If we let $I_0^\epsilon = \{i \in I \mid \|y_i^\epsilon\| < \delta\}$, we see that $I_0 \subset I_0^\epsilon$ for sufficiently small ϵ . If $i \in I_0^\epsilon$, then $a_{ij}^\epsilon = (K^\epsilon \phi_j, \psi_i^\epsilon) = 0$ for all $j \in I$. Hence, $\text{rank}(A^\epsilon)$ is not greater than the cardinality of $I - I_0^\epsilon$. Since $I - I_0^\epsilon \subset I - I_0$ for sufficiently small ϵ , we must have $\text{rank}(A^\epsilon) \leq \text{rank}(A)$. Therefore, $\text{rank}(A^\epsilon) = \text{rank}(A)$ for sufficiently small ϵ .

LEMMA 2.10. *Let B and B^ϵ be $n \times n$ matrices, and let c and c^ϵ be vectors in R^n . Suppose that $w = B^\dagger c$ and $w^\epsilon = (B^\epsilon)^\dagger c^\epsilon$. If $\|B - B^\epsilon\|_2 < \frac{1}{\|B^\dagger\|_2}$ and $\text{rank}(B^\epsilon) \leq \text{rank}(B)$, then $\text{rank}(B^\epsilon) = \text{rank}(B)$ and*

$$\begin{aligned} \|w - w^\epsilon\|_2 &\leq \|B^\dagger\|_2 \left[\frac{\|B - B^\epsilon\|_2 \|w\|_2 + \|c - c^\epsilon\|_2}{1 - \|B - B^\epsilon\|_2 \|B^\dagger\|_2} \right. \\ &\quad \left. + \frac{\|B - B^\epsilon\|_2 \|B^\dagger\|_2 \|c - Bw\|_2}{1 - \|B - B^\epsilon\|_2 \|B^\dagger\|_2} + \|B - B^\epsilon\|_2 \|w\|_2 \right]. \end{aligned}$$

(See [7, Theorem 9.7] for the proof.)

Proof of Theorem 2.4. Without the loss of generality, we can assume that $0 < \delta < \mu$. Let $z = A^\dagger b = (\beta_i)$ and $z^\epsilon = (A^\epsilon)^\dagger b^\epsilon = (\beta_i^\epsilon)$. Then, $f_M = \sum_{i=1}^n \beta_i \phi_i$ is the M -solution to $Kf = g$ and $f_M^\epsilon = \sum_{i=1}^n \beta_i^\epsilon \phi_i$ is the approximate M -solution to $K^\epsilon f = g^\epsilon$. Since $\{\phi_i\}_{i \in I}$ is an orthonormal set, $\|f_M - f_M^\epsilon\| = \|z - z^\epsilon\|_2$. By Lemmas 2.7 and 2.9, $\lim_{\epsilon \rightarrow 0} \|A - A^\epsilon\|_2 = 0$ and $\text{rank}(A) = \text{rank}(A^\epsilon)$ for sufficiently small ϵ . Therefore, applying Lemma 2.10, one obtains

$$\begin{aligned} \|f_M - f_M^\epsilon\| &\leq \|A^\dagger\|_2 \left[\frac{\|A - A^\epsilon\|_2 \|z\|_2 + \|b - b^\epsilon\|_2}{1 - \|A - A^\epsilon\|_2 \|A^\dagger\|_2} \right. \\ &\quad \left. + \frac{\|A - A^\epsilon\|_2 \|A^\dagger\|_2 \|b - Az\|_2}{1 - \|A - A^\epsilon\|_2 \|A^\dagger\|_2} + \|A - A^\epsilon\|_2 \|z\|_2 \right] \end{aligned}$$

for sufficiently small ϵ . Since $\lim_{\epsilon \rightarrow 0} \|A - A^\epsilon\|_2 = 0$ and $\lim_{\epsilon \rightarrow 0} \|b - b^\epsilon\|_2 = 0$, we see that $\lim_{\epsilon \rightarrow 0} \|f_M - f_M^\epsilon\| = 0$. Hence, the GGS(δ) method provides the M -solution to $Kf = g$ which is well-posed under perturbations in both K and g .

EXAMPLE 2.11. Let $H_1 = H_2$, ϕ_1 and ϕ_2 be orthonormal vectors in H_1 , $M = \langle \phi_1, \phi_2 \rangle$, and let P_i be the orthogonal projection of H_1 onto $\langle \phi_i \rangle$ for $i = 1, 2$. Define $K = P_1$ and $K^\epsilon = P_1 + \epsilon P_2$. Let $g = \alpha_1 \phi_1 + \alpha_2 \phi_2$ with $\alpha_2 \neq 0$ and $g^\epsilon = \alpha_1 \phi_1 + (\alpha_2 + \epsilon) \phi_2$. Then $\|K - K^\epsilon\| = \|g - g^\epsilon\| = \epsilon$, $K\phi_1 = \phi_1$, $K\phi_2 = 0$, $K^\epsilon\phi_1 = \phi_1$, and $K^\epsilon\phi_2 = \epsilon\phi_2$. Choose $\delta = \frac{1}{2}$. Then, it is easy to see that both the GGS(δ) and the GGS provide the exact M -solution $f_M = \alpha_1 \phi_1$ of $Kf = g$. We now consider the perturbed problem $K^\epsilon f = g^\epsilon$. Using the GGS method, $y_1^\epsilon = \psi_1^\epsilon = \phi_1$, $y_2^\epsilon = \epsilon\phi_2$, and $\psi_2^\epsilon = \phi_2$ are generated, and so $f_M^\epsilon = \alpha_1 \phi_1 + (1 + \frac{\alpha_2}{\epsilon}) \phi_2$. When ϵ converges to 0, f_M^ϵ never converges to f_M even if K^ϵ and g^ϵ converge to K and g respectively.

On the other hand, when using the GGS(δ) method, $y_1^\epsilon = \psi_1^\epsilon = \phi_1$, $y_2^\epsilon = \epsilon\phi_2$, and $\psi_2^\epsilon = \phi_2$ (if $\epsilon \geq \delta$) or 0 (if $\epsilon < \delta$) are generated. Hence, f_M^ϵ is equal to f_M for every ϵ which is less than δ . Therefore, the approximate M -solution f_M^ϵ for $K^\epsilon f = g^\epsilon$ obtained from the GGS(δ) method converges to the M -solution f_M of $Kf = g$ when ϵ converges to 0.

3. Numerical implementation of the GGS(δ)

In this section, it is assumed that $K : L_2[a, b] \rightarrow L_2[c, d]$ is an integral operator defined by $(Kf)(s) = \int_a^b k(s, t)f(t)dt$ with the kernel function $k(s, t) \in L_2([c, d] \times [a, b])$. The numerical method commonly used for integration of the form $\int_{-1}^1 f(x)dx$ is the Gauss-Legendre quadrature method. To this end, it is convenient to change variables so that we can work on $L_2[-1, 1]$ instead of $L_2[a, b]$ and $L_2[c, d]$. First, we show how the change of variables should be handled for the problem of finding the M -solution of $(Kf)(s) = g(s)$ for all $s \in [c, d]$, where $g \in L_2[c, d]$.

Let $s = h_1(x) = \frac{c+d}{2} + \frac{d-c}{2}x$ and $t = h_2(y) = \frac{a+b}{2} + \frac{b-a}{2}y$, where $-1 \leq x, y \leq 1$. Put $k_0(x, y) = k(h_1(x), h_2(y))$, $G(x) = g(h_1(x))$, and $F(y) = f(h_2(y))$. Define an integral operator $K_0 : L_2[-1, 1] \rightarrow L_2[-1, 1]$ by $(K_0F)(x) = \int_{-1}^1 k_0(x, y)F(y)dy$ for $F \in L_2[-1, 1]$ and $x \in [-1, 1]$. Then, $(Kf)(s) = g(s)$ is equivalent to $(K_0F)(x) = \frac{2}{b-a}G(x)$. Suppose that N is a closed subspace of $L_2[a, b]$ with orthonormal basis $\{\phi_j(t)\}_{j \in J}$, where J is an index set which is at most countable. Define $\psi_j(y) = \phi_j(h_2(y))$ for each $j \in J$. Let N_0 be a subspace of $L_2[-1, 1]$ which is the closed linear

span of $\{\psi_j(y)\}_{j \in J}$. It is clear that $(\phi_i, \phi_j)_{[a,b]} = \delta_{ij} = \frac{b-a}{2}(\psi_i, \psi_j)_{[-1,1]}$, where δ_{ij} represents the Kronecker delta, and $(\cdot, \cdot)_{[a,b]}$ and $(\cdot, \cdot)_{[-1,1]}$ refer to the inner products on $L_2[a, b]$ and $L_2[-1, 1]$ respectively. Using the above notations, we obtain the following.

LEMMA 3.1. $g \in K(N)$ if and only if $G \in K_0(N_0)$. Moreover, $g \in \overline{K(N)}$ if and only if $G \in \overline{K_0(N_0)}$, where the bars denote the closure of the given subspaces.

Let P be the orthogonal projection of $L_2[c, d]$ onto $\overline{K(N)}$. Define $P_0 : L_2[-1, 1] \rightarrow L_2[-1, 1]$ by $(P_0G)(x) = (Pg)(h_1(x))$, where $g = G \circ h_1^{-1} \in L_2[c, d]$. Then, it is easy to show that P_0 is the orthogonal projection of $L_2[-1, 1]$ onto $\overline{K_0(N_0)}$. Since $Kf = Pg$ is equivalent to $K_0F = P_0(\frac{2}{b-a}G)$ and $(f, f)_{[a,b]} = \frac{b-a}{2}(F, F)_{[-1,1]}$, the following theorem is immediately obtained.

THEOREM 3.2. $f_N(t)$ is the N -solution of $Kf = g$ if and only if $F_{N_0}(y) = (f_N \circ h_2)(y)$ is the N_0 -solution of $K_0F = \frac{2}{b-a}G$.

By taking $N = L_2[a, b]$ in Theorem 3.2, we obtain the following.

COROLLARY 3.3. $f_0(t)$ is the LSSMN of $Kf = g$ if and only if $F_0(y) = (f_0 \circ h_2)(y)$ is the LSSMN of $K_0F = \frac{2}{b-a}G$.

Let's now consider the numerical implementation of the GGS(δ) for approximating the M -solution of the form

$$(Kf)(s) = \int_0^1 k(s, t)f(t)dt = g(s), \quad 0 \leq s \leq 1.$$

First, we will transform it to the equivalent problem

$$(K_0F)(x) = \int_{-1}^1 k_0(x, y)F(y)dy = 2G(x), \quad -1 \leq x \leq 1,$$

by using the change of variables $s = \frac{1+x}{2}$ and $t = \frac{1+y}{2}$, where $k_0(x, y) = k(\frac{1+x}{2}, \frac{1+y}{2})$, $F(y) = f(\frac{1+y}{2})$, and $G(x) = g(\frac{1+x}{2})$. Next, we will choose an n dimensional subspace $(M_0)_n$ of $L_2[-1, 1]$ so that $(M_0)_n = \langle \psi_1(y), \dots, \psi_n(y) \rangle$, where $\{\psi_i(y)\}_{i=1}^n$ is an appropriate orthonormal set in

$L_2[-1, 1]$. As a typical example for such an orthonormal set in $L_2[-1, 1]$, the normalized Legendre polynomials can be chosen. Then, the $(M_0)_n$ -solution $F_n(y)$ of the transformed problem $K_0F = 2G$ is approximated using the GGS(δ) method. Let $\hat{F}_n(y)$ denote the approximate $(M_0)_n$ -solution to $K_0F = 2G$ obtained from the GGS(δ). Note that from Theorem 3.2, $f_n(t) = F_n(2t - 1)$ is the M_n -solution of the given problem $Kf = g$, where $M_n = \langle \phi_1(t), \dots, \phi_n(t) \rangle$ and $\phi_i(t) = \psi_i(2t - 1)$. Thus, the approximate M_n -solution to $Kf = g$ computed from the GGS(δ) is given by $\hat{f}_n(t) = \hat{F}_n(2t - 1)$.

All computations were done in Fortran using double precision arithmetic which means about fifteen decimal digits of accuracy. The number δ for the GGS(δ) method was taken to be 10^{-13} . The second step of the GGS(δ) requires the computation of the minimum norm solution of $Ax = b$. If A has full rank, then A is an upper triangular matrix and so the minimum norm solution of $Ax = b$ is solved by the Linpack subroutine DTRSL [3, Chapter 6] which solves a linear system with triangular matrix. If A does not have full rank, then the minimum norm solution of $Ax = b$ is solved using the Linpack subroutine DSVDC [3, Chapter 11] which carries out the singular value decomposition of a matrix. The decision of whether or not A has full rank was made in the following way: If $\|y_i\| < \delta$ for an $i \in I$, then all elements of the i -th row of A consist of zeros and hence A does not have full rank. If $\|y_i\| \geq \delta$ for every $i \in I$, then A has full rank.

The DSVDC supplies $\{\sigma_1, \dots, \sigma_n\}$, the singular values of A . In order to decide which of σ_i 's are nonzero, the same tolerance δ mentioned above is used. In other words, if $\sigma_i < \delta$, then σ_i is set to zero. All inner products required during the GGS(δ) were done using the repeated 4-point Gauss-Legendre quadrature method. The examples tested in this paper are described below. For each example, f_0 denotes the LSSMN to $Kf = g$, and $(M_0)_n = \langle \psi_1(y), \dots, \psi_n(y) \rangle$ is chosen as an n -dimensional subspace of $L_2[-1, 1]$, where $\psi_i(y)$ is the normalized Legendre polynomial of degree $(i - 1)$. The exact M_n -solution f_n is also provided in each of the examples below if possible.

EXAMPLE 3.4. Let $k(s, t) = s + t$ and $g(s) = s$. Then $f_0(t) = 4 - 6t$, and $f_n(t) = f_0(t)$ for $n \geq 2$.

EXAMPLE 3.5. Let $k(s, t) = \cos(st)$ and $g(s) = \frac{\sin s}{s} + \frac{\cos s - 1}{s^2}$. Then $f_0(t) = t$, and $f_n(t) = f_0(t)$ for $n \geq 2$.

EXAMPLE 3.6. Let $k(s, t) = (s - t)^2$ and $g(s) = s^2 - \frac{2}{3}s + \frac{1}{4}$. Then $f_0(t) = 15t^2 - 17t + \frac{9}{2}$, $f_2(t) = \frac{29}{12} - 2t$, and $f_n(t) = f_0(t)$ for $n \geq 3$.

EXAMPLE 3.7. Let $k(s, t) = \sin 2(s + t - 1)$ and $g(s) = 5 \sin(2s - 1) + 3 \cos(2s - 1)$. Then $f_0(t) = \left(\frac{10}{2 + \sin 2}\right) \cos(2t - 1) + \left(\frac{6}{2 - \sin 2}\right) \sin(2t - 1)$, $f_2(t) = \frac{5}{2 \sin 1} + \frac{3}{2(\sin 1 - \cos 1)}(2t - 1)$, and so on.

The numerical results for the above examples are summarized in Table 1. $\|f_n - \hat{f}_n\|$ is computed to see how close the approximate M_n -solution \hat{f}_n obtained from the GGS(δ) is to the exact M_n -solution f_n , and $\|f_0 - \hat{f}_n\|$ is computed to see how well the \hat{f}_n approximates f_0 . Since the exact M_n -solution f_n for Example 3.7 is not available by hand calculation if $n > 3$, the value of $\|f_n - \hat{f}_n\|$ is not listed in Table 1 for $n = 4$ and 5. The numbers in Table 1 under the column $\|f_n - \hat{f}_n\|$ range from 10^{-16} to 10^{-13} , which shows numerically that the GGS(δ) is a well-posed method for finding the M -solution to $Kf = g$. As can be seen from Table 1, for Examples 3.4 to 3.6 \hat{f}_n converges very fast to f_0 , whereas for Example 3.7 \hat{f}_n converges very slowly to f_0 . This is because we cannot approximate $\sin y$ and $\cos y$ very accurately using low-degree Legendre polynomials.

If we choose $(M_0)_2 = \langle \sin y, \cos y \rangle$ for Example 3.7, then it is easy to see that $f_2(t)$ equals $f_0(t)$. For this choice of $(M_0)_2$, the computed value of $\|f_0 - \hat{f}_2\| = \|f_2 - \hat{f}_2\|$ is about 10^{-14} , i.e., the computed M_2 -solution \hat{f}_2 is a good approximation to f_0 as well as f_2 . Hence, in order for the computed M_n -solution \hat{f}_n to approximate f_0 well, it is natural to choose M_n to reflect a priori knowledge of the f_0 . Since the goal of this paper is to develop a well-posed method for finding the M -solution to $Kf = g$, the problem of choosing a suitable subspace M_n so that f_n can approximate f_0 well is not considered herein and it will be studied in the future work.

Table 1. Numerical results for the GGS(δ)

Example	n	$\ f_n - \tilde{f}_n\ $	$\ f_0 - \tilde{f}_n\ $
3.4	2	5.10×10^{-15}	5.10×10^{-15}
	3	5.89×10^{-15}	5.89×10^{-15}
	4	5.89×10^{-15}	5.89×10^{-15}
3.5	2	1.09×10^{-13}	1.09×10^{-13}
	3	1.09×10^{-13}	1.09×10^{-13}
	4	1.09×10^{-15}	1.09×10^{-15}
3.6	2	3.74×10^{-16}	1.1932
	3	1.96×10^{-15}	1.96×10^{-15}
	4	1.96×10^{-15}	1.96×10^{-15}
3.7	2	6.53×10^{-14}	7.08×10^{-1}
	3	1.98×10^{-13}	5.75×10^{-1}
	4		5.54×10^{-1}
	5		5.53×10^{-1}

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