THE PERIOD AND THE LINEAR COMPLEXITY OF CERTAIN LINEAR RECURRING SEQUENCES IN THE FINITE FIELD GF(q)

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1. Introduction

Let GF(q) be a finite field with q elements, where q is a prime power. A sequence

$$s_0, s_1, s_2, \cdots, s_t, s_{t+1}, \cdots$$

of elements of GF(q) is denoted by $\{s_t\}$. The set $GF(q)^{\omega}$ of all sequences in GF(q) forms an algebra over GF(q) under the addition, multiplication and the scalar multiplication defined as follows:

$$\{u_t\} + \{v_t\} = \{u_t + v_t\},\$$
$$\{u_t\} \{v_t\} = \{u_t v_t\},\$$
$$c\{u_t\} = \{cu_t\}$$

If there exists a positive integer r such that $s_{t+r} = s_t$ for all $t \ge 0$, then $\{s_t\}$ is said to be *periodic*, and the smallest positive integer r such that $s_{t+r} = s_t$ for all $t \ge 0$ is called the *period* of the sequence.

Let c_0, c_1, \dots, c_{n-1} be given elements of the finite field GF(q). A sequence $\{s_t\}$ in GF(q) satisfying the linear recurrence relation

$$s_{t+n} = \sum_{i=1}^{n-1} c_i s_{t+i}$$
 $(t = 0, 1, 2, \cdots)$

is called an n-th homogeneous linear recurring sequence in GF(q) and the polynomial

$$f(x) = -c_0 - c_{1x} - c_2 x^2 - \dots - c_{n-1} x^{n-1} + x^n$$

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is called the *characteristic polynomial* of $\{s_t\}$. And for any monic polynomial

$$q(x) = -a_0 - a_1 x - a_2 x^2 - \dots - a_{n-1} x^{n-1} + x^n$$

over GF(q) of degree $n \geq 1$, a sequence $\{s_t\}$ in GF(q) satisfying

$$s_{t+n} = \sum_{i=0}^{n-1} a_i s_{t+i} \qquad (t = 0, 1, 2, \cdots)$$

is said to be generated by g(x). The set of all homogeneous linear recurring sequence in GF(q), generated by g(x) forms a subspace of the vector space $GF(q)^{\omega}$, which is called the solution space of g(x) and is denoted by S(g(x)).

Let $\{s_t\}$ be a homogeneous linear recurring sequence in GF(q). Then there exists a unique monic polynomial $h(x) \in GF(q)[x]$, deg $h(x) \ge 1$, such that for any monic polynomial $g(x) \in GF(q)[x]$ we have

$$\{s_t\} \in S(g(x))$$
 if and only if $h(x)|g(x)$.

The polynomal h(x) is called the minimum polynomial of $\{s_t\}$ and the degree of h(x) is called the linear complexity or linear equivalence of $\{s_t\}$.

In this paper we will prove some theorems on the period and the linear complexity of certain sequences in GF(q) which are generated by combining two sequences in a reasonable way. In fact these theorems are generalizations of the main result in [1].

A sequence of elements of GF(2) is called a binary sequence. In recent years considerable interest has been shown in the generation of binary sequences which have good properties. Such binary sequences play an important role in a stream cipher system.

The terminology and the notation in this paper are standard, and they are taken from [3]. Throughout this paper, GF(q) denotes the finite field with q elements, where q is a prime power.

2. Preliminaries

In this section we will state some propositions which are useful in proving theorems in section 3. The proofs of these propositions will be omitted and they can be found in [3] and [6].

Let GF(q) be the finite field with q elements, where q is a prime power, and let $f(x) \in GF(q)[x]$ be a polynomial of degree $n \geq 1$ with $f(0) \neq 0$. Then there exits a positive integer $e \leq q^n - 1$ such that $f(x)|(x^e-1)$ in GF(q)[x]. The smallest positive integer e such that $f(x)|(x^e-1)$ in GF(q)[x] is called the order of f(x) and is denoted by ord f(x). If f(x) is an irreducible polynomial over GF(q) of degree n, then the order of f(x) divides $q^n - 1$.

The multiplicative group $GF(q)^* = GF(q) - \{0\}$ is cyclic. A generator of the cyclic group $GF(q)^*$ is called a primitive element of GF(q). A polynomial $f(x) \in GF(q)[x]$ of degree $n \geq 1$ is called a primitive polynomial if it is the minimum polynomial over GF(q) of a primitive element of $GF(q^n)$. A monic polynomial $f(x) \in GF(q)[x]$ of degree n is a primitive polynomial over GF(q) if and only if it is irreducible over GF(q) and ord $f(x) = q^n - 1$.

Let $\{s_t\}$ be any sequence in GF(q) and let d be a positive integer. A sequence $\{u_t\}$ such that

$$u_t = s_{t+d} \qquad (t = 0, 1, 2, \cdots)$$

is called the translate of $\{s_t\}$ by d, and it is denoted by $d\{s_t\}$. And a sequence $\{v_t\}$ such that

$$v_t = s_{dt} \qquad (t = 0, 1, 2, \cdots)$$

is called the decimation of $\{s_t\}$ by d, and it is denoted by $\{s_t\}^{(d)}$.

PROPOSITION 2.1. Let $\{s_t\}$ be a homogeneous linear recurring sequence generated by a polynomial $f(x) \in GF(q)[x]$ of degree $n \geq 1$. Then the period of the sequence divides ord f(x) and so it is $\leq q^n - 1$.

If $h(x) \in GF(q)[x]$ is the minimum polynomial of $\{s_t\}$, then the period of the sequence is equal to ord h(x).

Seung Ahn Park

PROPOSITION 2.2. Let f(x) be a primitive polynomial over GF(q) of degree $n \geq 1$. Then any nonzero sequence in GF(q) generated by f(x) is of period $q^n - 1$.

Such a sequence is called a maximal period sequence in GF(q).

PROPOSITION 2.3. Let $f(x) \in GF(q)$ be a primitive polynomial of degree $n \geq 1$ and let $\{s_t\}$ be a nonzero sequence generated by f(x). Then the sequence $\{s_t\}$ in GF(q) is a maximal period sequence of period $q^n - 1$ and

$$S(f(x)) = \{i\{s_t\} \mid 0 \le i \le q^n - 2\} \cup \{\{0\}\}.$$

PROPOSITION 2.4. Let $f(x) \in GF(q)[x]$ be an irreducible polynomial of degree $n \geq 1$ such that $\alpha \in GF(q^n)$ is a root of f(x), and let d be a positive integer. If h(x) is an irreducible polynomial over GF(q) of degree n with $h(\alpha^d) = 0$, then

$$S(h(x)) = \{ \{s_t\}^{(d)} | \{s_t\} \in S(f(x)) \}.$$

Proposition 2.5. Let

$$f_1(x), f_2(x), \cdots, f_r(x)$$

be all the distinct monic irreducible polynomials in GF(q)[x] of degree n and order N, and let $M \geq 2$ be an integer which satisfies the following two conditions:

- (i) All prime factors of M divide N but not $(q^n 1)/N$.
- (ii) $q^n \equiv 1 \pmod{4}$ if $M \equiv 0 \pmod{4}$.

Then

$$f_1(x^M), f_2(x^M), \cdots, f_r(x^M)$$

are all the distinct monic irreducible polynomials in GF(q)[x] of degree nM and order NM.

3. The period and the linear complexity of certain sequences

Let GF(q) be the finite field with q elements, where q is a prime power. In this section we will study the period and the linear complexity of certain sequences which are generated by combining two maximal period sequences in GF(q).

Let f(x) and g(x) be primitive polynomials over GF(q) of degree m and n respectively, where m > 1 and n > 1. Let k be an integer such that

$$1 \le k < m$$

and let i_1, i_2, \dots, i_{k-1} be integers such that

$$0 < i_1 < i_2 < \cdots < i_{k-1} < m$$
.

And let $\mathbf{Z}_q = \{0, 1, 2, \dots, q-1\}$ be the set of all nonnegative integers less than q, and let

$$\tau: GF(q) \longrightarrow \mathbf{Z}_q, \ \tau(a) = a'$$

be any injective map such that $\tau(0) = 0$.

Let $\{a_t\}$ and $\{b_t\}$ be nonzero sequences in GF(q) generated by f(x) and g(x), respectively. Thus $\{a_t\}$ and $\{b_t\}$ are maximal period sequences of period $q^m - 1$ and $q^n - 1$, respectively. For each vector

$$(a_t, a_{t+i_1}, \cdots, a_{t+i_{k-1}}) \in GF(q)^k$$
 $(t = 0, 1, 2, \cdots)$

let R_t be an integer given by

$$R_t = a_t' q^{k-1} + a_{t+i}' q^{k-2} + \dots + a_{t+i-1}'$$
 $(t = 0, 1, 2, \dots)$

and let $\{u_t\}$ be a sequence in GF(q) such that

$$u_t = b_d, \qquad (t = 0, 1, 2, \cdots)$$

where

$$d_t = t + \sum_{i=0}^t R_i$$
 $(t = 0, 1, 2, \cdots)$

From now on we will study the period and the linear complexity of the sequence $\{u_t\}$ in GF(q). The notation introduced in the above sentences will keep its meaning throughout this section.

THEOREM 3.1. Let

$$M = q^m - 1$$
, $N = q^n - 1$, $K = M + \sum_{t=0}^{M-1} R_t$.

Then the following hold.

(1)
$$K = M + \frac{q^m(q^k - 1)}{2} = \frac{q^m(q^k + 1)}{2} - 1$$

(2) $(i\{u_t\})^{(M)} = (d_i\{b_t\})^{(K)}$ $(i = 0, 1, \dots, M - 1)$

Proof. (1) Since $\{a_t\}$ is a maximal period sequence of period M, we have

$$\{(a_t, a_{t+1}, \dots, a_{t+m-1}) \mid 0 \le t \le M-1\} \cup \{(0, \dots, 0)\} = GF(q)^m$$

by Proposition 2.3. Hence

$$\{(a_t, a_{t+i_t}, \cdots, a_{t+i_{k-1}}) \mid 0 \le t \le M-1\} = GF(q)^k$$

and every nonzero vector in $GF(q)^k$ occurs exactly q^{m-k} times while t varies from 0 to M-1. Since

$$\tau: GF(q) \longrightarrow \mathbf{Z}_q. \ \tau(a) = a'$$

is injective map such that $\tau(0) = 0$, it follows that

$$\{(a'_t, a'_{t+i_1}, \cdots, a'_{t+i_{k-1}}) \mid 0 \le t \le M-1\} = \mathbf{Z}_{q^k}$$

and every nonzero vector in \mathbf{Z}_{q^k} occurs exactly q^{m-k} times while t varies from 0 to M-1. Therefore, we have

$$\sum_{t=0}^{M-1} R_t = \sum_{t=0}^{M-1} a_t' q^{k-1} + a_{t+i_1}' q^{k-2} + \dots + a_{t+i_{k-1}}'$$

$$= q^{m-k} \{1 + 2 + \dots + (q^k - 1)\} = \frac{q^m (q^k - 1)}{2}$$

and

$$K = M + \sum_{t=0}^{M-1} R_t = q^m - 1 + \frac{q^m(q^k - 1)}{2} = \frac{q^m(q^k + 1)}{2} - 1$$

(2) Since $\{a_t\}$ is a periodic sequence of period M, we have

$$R_{i+jM} = R_i$$
 $(0 \le i \le M-1, j = 0, 1, 2, \cdots)$

and so

$$d_{i+jM} = i + jM + \sum_{t=0}^{i+jM} R_t = jM + j \sum_{t=0}^{M-1} R_t + i + \sum_{t=0}^{i} R_t$$
$$= jK + d_i \qquad (0 \le i \le M-1, \ j = 0, 1, 2, \cdots)$$

Therefore, we have

$$u_{i+jM} = b_{jK+d_i}$$
 $(0 \le i \le M-1, j = 0, 1, 2, \cdots)$

and so
$$(i\{u_t\})^{(M)} = (d_i\{b_t\})^{(K)}$$
 $(i = 0, 1, \dots, M-1).$

THEOREM 3.2. Let

$$M = q^m - 1$$
, $N = q^n - 1$, $K = M + \frac{q^m(q^k - 1)}{2}$

Let β be a primitive element of $GF(q^n)$ such that $g(\beta) = 0$ and let h(x) be the minimum polynomial of β^K over GF(q).

If (N, K) = 1, then the following hold.

- (1) The polynomial h(x) is a primitive polynomial over GF(q) such that $\deg h(x) = n$, $\operatorname{ord} h(x) = N$, and $h(x) = \prod_{i=0}^{n-1} (x \beta^{Kq^i})$.
- (2) For each $j = 0, 1, 2, \dots$, the sequence $(j\{b_t\})^{(K)}$ in GF(q) has h(x) as its miminum polynomial and it is of period N.
- (3) The sequence $\{u_t\}$ in GF(q) is generated by the polynomial $h(x^M)$ over GF(q), and

$$h(x^M) = \prod_{i=0}^{n-1} (x^M - \beta^{Kq^i}), \qquad \deg h(x^M) = nM.$$

Proof. By assumption g(x) is a primitive polynomial over GF(q) of degree n, and so there exists a primitive element $\beta \in GF(q^n)$ such that $g(\beta) = 0$.

- (1) We have $GF(q^n)^* = \langle \beta \rangle$ and $|GF(q^n)^*| = q^n 1 = N$. Since (N, K) = 1 it follows that $GF(q^n)^* = \langle \beta^K \rangle$. Therefore, β^K is a primitive element of $GF(q^n)$ and the order of h(x) is N. Hence the assertion (1) holds.
- (2) Since g(x) is a primitive polynomial over GF(q) with $g(\beta) = 0$ and $\{b_t\}$ is a nonzero sequence generated by g(x), we have

$$S(g(x)) = \{j\{b_t\} \mid 0 \le j \le N-1\} \cup \{\{0\}\}\$$

by Proposition 2.3. Hence we have

$$S(h(x)) = \{(j\{b_t\})^{(K)} \mid 0 \le j \le N-1\} \cup \{\{0\}\}\$$

by Proposition 2.4. Therefore, each sequence $(j\{b_t\})^{(K)}$ has h(x) as its miminum polynomial and it is of period N.

Thus the assertion (2) holds.

(3) Arrange elements of $\{u_t\}$ as follows:

In this way we obtain M sequences

$$\{u_t\}^{(M)}, ({}_1\{u_t\})^{(M)}, \cdots, ({}_{M-1}\{u_t\})^{(M)}$$

in GF(q). On the other hand, we have

$$(i\{u_t\})^{(M)} = (d_i\{b_t\})^{(K)} \qquad (i = 0, 1, \dots, M-1)$$

by Theorem 3.1. Hence each $(i\{u_t\})^{(M)}$ has h(x) as its minimum polynomial and $(i\{u_t\})^{(M)} \in S(h(x))$ by the assertion (2).

Suppose that

$$h(x) = -c_0 - c_1 x - \dots - c_{n-1} x^{n-1} + x^n \in GF(q)[x].$$

Then, for each $i = 0, 1, \dots, M - 1$, we have

$$u_{M(t+n)+i} = \sum_{j=0}^{n-1} c_j u_{M(t+j)+i}, \qquad (t=0,1,2,\cdots)$$

In particular we have

$$u_{t+nM} = \sum_{j=0}^{n-1} c_j u_{t+jM}, \qquad (t=0,1,2,\cdots)$$

and so $\{u_t\}$ is generated by the polynomial

$$h(x^M) = -c_0 - c_1 x^M - \dots - c_{n-1} x^{(n-1)M} + x^{nM} \in GF(q)[x].$$

Hence the assertion (3) holds.

THEOREM 3.3. Assume that the following three conditions hold.

- (i) $\left(q^{n}-1, q^{m}-1+\frac{q^{m}(q^{k}-1)}{2}\right)=1.$
- (ii) $q^n \equiv 1 \pmod{4}$ if $q^m \equiv 1 \pmod{4}$
- (iii) All prime factors of q^m-1 divide q^n-1 .

Then the sequence $\{u_t\}$ in GF(q) is of period $(q^n-1)(q^m-1)$ and linear complexity $n(q^m-1)$.

Proof. We will use the notation in Theorem 3.2. Note that

$$M = q^m - 1$$
, $N = q^n - 1$, $K = M + \frac{q^m(q^k - 1)}{2}$

The condition (i) implies that (N, K) = 1. Hence $\{u_t\}$ is generated by $h(x^M) \in GF(q)[x]$, by Theorem 3.2.

On the other hand, since the conditions (ii) and (iii) hold, $h(x^M)$ is a monic irreducible polynomial in GF(q)[x] of degree nM and order NM by Proposition 2.5.

Hence $h(x^M)$ is the minimum polynomial of $\{u_t\}$, and so $\{u_t\}$ is of period $(q^n-1)(q^m-1)$ and linear complexity $n(q^m-1)$.

In the case when q = 2 Theorem 3.3 can be rewritten as the following corollary, which is the main theorem in [1].

COROLLARY 3.4. Assume that q = 2 and the following two conditions hold.

- (i) $(2^{m}-1, 2^{m}-1+2^{m-1}(2^{k}-1))=1.$
- (ii) All prime factors of $2^m 1$ divide $2^n 1$.

Then the sequence $\{u_t\}$ in GF(2) is of period $(2^n-1)(2^m-1)$ and linear complexity $n(2^m-1)$.

Now consider the case when m = n in Theorem 3.3.

In Theorem 3.3, if m = n then the conditions (ii) and (iii) are satisfied automatically, and the condition (i) is equivalent to the condition

(*)
$$\left(q^n - 1, \frac{q^n(q^{k-1})}{2}\right) = 1.$$

If q is even, then the condition (*) is equivalent to the condition $(q^n - 1, q^k - 1) = 1$. Since $q^{(n,k)} - 1$ is a common divisor of $q^n - 1$ and $q^k - 1$, the condition $(q^n - 1, q^k - 1) = 1$ holds if and only if q = 2 and (n, k) = 1.

If q is odd, then (*) is equivalent to the condition

$$(**) \left(q^n - 1, \frac{q^k - 1}{2}\right) = 1.$$

Now assume that q is odd and that the condition (**) holds. Then $q^{(n,k)} - 1$ is divided by 2, and so q = 3 and (n,k) = 1. And k is odd if and only if $(3^k - 1)/2$ is odd. Therefore, if q is odd and the condition (**) holds then q = 3, (n, k) = 1 and k is odd.

Conversely, assume that q = 3, (n, k) = 1 and k is odd. And suppose that there exists an odd prime p which divides both $3^n - 1$ and $(3^k - 1)/2$. Then p is a common divisor of $3^n - 1$ and $3^k - 1$ and so

$$3^n \equiv 1 \pmod{p}$$
 and $3^k \equiv 1 \pmod{p}$.

Let r be the order of 3 modulo p. Then the above congruences yield that r|(n,k) and so r=1. This implies that $3 \equiv 1 \pmod{p}$ and p=2, which is a contradiction. Hence if q=3, (n,k)=1 and k is odd, then the condition (**) holds.

Therefore, we obtain the following corollary.

COROLLARY 3.5. Assume that m = n. Then the following hold.

- (1) If q = 2 and (n, k) = 1, then the sequence $\{u_t\}$ in GF(2) is of period $(2^n-1)^2$ and linear complexity $n(2^n-1)$.
- (2) If q = 3, (n, k)=1 and k is odd, then the sequence $\{u_t\}$ in GF(3) is of period $(3^n-1)^2$ and linear complexity $n(3^n-1)$.

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