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STRUCTURE OF COUPLE CATEGORIES

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1. Introduction

Let \mathbf{S} be the category of all sets and \mathbf{C} be a category. A covariant functor $F: \mathbf{C} \rightarrow \mathbf{S}$ will be called a *grounding* of \mathbf{C} . Dually, a contravariant functor $F: \mathbf{C}^* \rightarrow \mathbf{S}$ is called a *cogrounding* of \mathbf{C} , where \mathbf{C}^* is the dual category of \mathbf{C} . $\text{Cat}(\mathbf{C}, \mathbf{S})$ denotes the functor category, which has functors from a small category \mathbf{C} to \mathbf{S} for objects and for morphisms the natural transformations between two functors from \mathbf{C} to \mathbf{S} , and $\text{Co}(\mathbf{C}, \mathbf{S})$ denotes the couple category of \mathbf{C} [5].

It is known that any small category \mathbf{C} can be embedded into $\text{Cat}(\mathbf{C}, \mathbf{S})$ and $\text{Co}(\mathbf{C}, \mathbf{S})$

In this note, we shall be to establish the results on the conjugates of groundings and cogroundings of a category \mathbf{C} and some results on couple categories. Next we shall prove that the functor category $\text{Cat}(\mathbf{C}^*, \mathbf{S})$ and the full subcategory $\text{Con}(\mathbf{C}, \mathbf{S})$ of $\text{Co}(\mathbf{C}, \mathbf{S})$, has the natural couple (G, G^*, m) for objects and for morphisms the conjoint transformations $(G, G^*, m) \rightarrow (G', G'^*, m)$, are equivalent.

2. Conjugates of groundings and cogroundings

Any cogrounding F of a category \mathbf{C} is said to be *dominated* by a set S of objects of \mathbf{C} if every set $F(A)$, A is an objects in \mathbf{C} , is a union of sets $F(f) [F(B)]$, B ranging over the elements of S and f ranging over morphisms in $\text{Hom}_{\mathbf{C}}(A, B)$.

If F is dominated by some set of objects, it is called *proper*.

DEFINITION 1. Let $F: \mathbf{C}^* \rightarrow \mathbf{S}$ be a cogrounding of a category \mathbf{C} . The grounding $F^*: \mathbf{C} \rightarrow \mathbf{S}$, such that for each object A of \mathbf{C} , $F^*(A)$ is the set $\text{Hom}(F, h_A)$ of all natural transformations from F to $h_A = \text{Hom}_{\mathbf{C}}(_, A)$ and for each morphism $f: A \rightarrow A'$ in \mathbf{C} , $p \in F(W)$ and $\phi \in F^*(A)$,

$$[F^*(f)(\phi)]_W(p) = f \cdot \phi_W(p)$$

is called the *conjugate* of F .

Dually we can define the conjugate G_* of grounding G as follows;

$$G_*(A) = \text{Hom}(G, h^A),$$

$$[G_*(f)(\xi)]_W(p) = \xi_W(p) \cdot f$$

for $p \in G(W)$, $f: A \rightarrow B$, $\xi \in G_*(B)$.

PROPOSITION 1. *If a cogrounding F of \mathbf{C} and its conjugate F^* are proper then there is a natural transformation from F to F_{**} .*

Proof. Let $\eta: F \rightarrow F_{**}$ be as follows: For each $A \in \text{Ob}(\mathbf{C})$ (=the class of objects of \mathbf{C}),

$$\begin{array}{ccc} \eta_A: F(A) & \longrightarrow & F_{**}(A) = \text{Hom}(F^*, h_A) \\ \Downarrow & & \Downarrow \\ p & \longmapsto & \hat{p} \end{array}$$

where $\hat{p}_B(p^*) = p^*_{A'}(p)$ for all $p^* \in F^*(B)$. Then for $g: A \rightarrow A'$ in \mathbf{C} and $p' \in F(A')$,

$$\begin{aligned} [\eta_A \cdot F(g)](p') &= \eta_A[F(g)(p)] \\ [F_{**}(g) \cdot \eta_{A'}](p') &= F_{**}(g)[\eta_{A'}(p')] = h^{(g)} \cdot \eta_{A'}(p'), \end{aligned}$$

where $h^{(g)}: h^{A'} \rightarrow h^A$.

For $B \in \text{Ob}(\mathbf{C})$ and $q \in F^*(B)$

$$[\eta_A[F(g)(p')]]_B(q) = q_A[F(g)(p')]$$

while

$$\begin{aligned} [h^{(g)} \cdot \eta_{A'}(p')]_B(q) &= h_B^{(g)}[\eta_{A'}(p')]_B(q) \\ &= q_{A'}(p') \cdot g. \end{aligned}$$

Since $q: F \rightarrow h_B$ is a natural transformation, we have

$$q_A[F(g)(p')] = h^{(g)}[q_{A'}(p')] = q_{A'}(p') \cdot g$$

Therefore $\eta: F \rightarrow F_{**}$ is a natural transformation.

We say that a grounding (cogrounding) F of \mathbf{C} is *reflexive* if F is proper, and F^* is proper and the natural transformation $\eta: F \rightarrow F_{**}$ is the natural equivalence.

PROPOSITION 2. *Let $\Gamma_1: \text{Cat}(\mathbf{C}^*, \mathbf{S}) \rightarrow \text{Cat}(\mathbf{C}, \mathbf{S})$ and $\Gamma_2: \text{Cat}(\mathbf{C}, \mathbf{S})^* \rightarrow \text{Cat}(\mathbf{C}^*, \mathbf{S})$ be two contravariant functors such that $\Gamma_1(F) = F^*$ for all $F \in \text{Ob}(\text{Cat}(\mathbf{C}^*, \mathbf{S}))$ and for $\eta: F_1 \rightarrow F_2$ in $\text{Cat}(\mathbf{C}^*, \mathbf{S})$, $\xi \in F_2^*(X)$, $[\Gamma_1(\eta)]_X(\xi) = \xi \cdot \eta$, and for $G \in \text{Ob}(\text{Cat}(\mathbf{C}, \mathbf{S}))$, $\mu: G_1 \rightarrow G_2$ in $\text{Cat}(\mathbf{C}, \mathbf{S})$, $\phi \in G_2^*(X)$, $\Gamma_2(G) = G_*$, $[\Gamma_2(\mu)]_X(\phi) = \phi \cdot \mu$. Then the two functors $\text{Hom}_{\text{Cat}(\mathbf{C}^*, \mathbf{S})}(-, \Gamma_2)$ and $\text{Hom}_{\text{Cat}(\mathbf{C}, \mathbf{S})}(-, \Gamma_1)$ from the product category $\text{Cat}(\mathbf{C}^*, \mathbf{S})^* \times \text{Cat}(\mathbf{C}, \mathbf{S})^*$ to \mathbf{S} are naturally equivalent, where for $(F, G) \in \text{Cat}(\mathbf{C}^*, \mathbf{S})^* \times \text{Cat}(\mathbf{C}, \mathbf{S})^*$,*

$[\text{Hom}(-, \Gamma_1-)](F, G) = \text{Hom}(G, \Gamma_1(F) = F^*)$ and $[\text{Hom}(-, \Gamma_2-)](F, G) = \text{Hom}(F, \Gamma_2(G) = G_*)$ respectively.

Proof. For each $(F, G) \in \text{Cat}(\mathbf{C}^*, \mathbf{S})^* \times \text{Cat}(\mathbf{C}, \mathbf{S})^*$, we define a mapping $\psi_{(F,G)} : \text{Hom}(G, F^*) \rightarrow \text{Hom}(F, G_*)$ as follows; for $\phi \in \text{Hom}(G, F^*)$, $p \in F(X)$, $q \in G(W)$ and $X, W \in \text{Ob}(\mathbf{C})$

$$[[\psi_{(F,G)}(\phi)]_X(p)]_W(q) = [\phi_W(q)]_X(p).$$

Then for $\phi_1, \phi_2 \in \text{Hom}(G, F^*)$, if

$$\psi_{(F,G)}(\phi_1) = \psi_{(F,G)}(\phi_2),$$

$$[[\phi_1]_W(q)]_X(p) = [[\phi_2]_W(q)]_X(p)$$

for all $p \in F(X)$ and $q \in G(W)$. We have $[\phi_1]_W = [\phi_2]_W$ and $\phi_1 = \phi_2$. Next we shall show that the diagram

$$\begin{array}{ccc} \text{Hom}(G_1, F_1^*) & \xrightarrow{\psi_{(F_1, G_1)}} & \text{Hom}(F_1, G_1^*) \\ \text{Hom}(\eta, \Gamma_1(\nu)) \downarrow & & \downarrow \text{Hom}(\nu, \Gamma_2(\eta)) \\ \text{Hom}(G_2, F_2^*) & \xrightarrow{\psi_{(F_2, G_2)}} & \text{Hom}(F_2, G_2^*) \end{array}$$

where $\eta : G_2 \rightarrow G_1 \in \text{Cat}(\mathbf{C}, \mathbf{S})$ and $\nu : F_2 \rightarrow F_1 \in \text{Cat}(\mathbf{C}^*, \mathbf{S})$, is commutative. For each $\phi \in \text{Hom}(G_1, F_1^*)$,

$$[\text{Hom}(\nu, \Gamma_2(\eta)) \cdot \psi_{(F_1, G_1)}(\phi)](\phi) = \Gamma_2(\eta) \cdot \psi_{(F_1, G_1)}(\phi) \cdot \nu$$

and for $p_2 \in F_2(X)$, $q_2 \in G_2(W)$,

$$\begin{aligned} & [[\Gamma_2(\eta) \cdot \psi_{(F_1, G_1)}(\phi) \cdot \nu]_X(p_2)]_W(q_2) \\ &= [[\Gamma_2(\eta)]_X \cdot [\psi_{(F_1, G_1)}(\phi)]_X \cdot \nu_X(p_2)]_W(q_2) \\ &= [[\psi_{(F_1, G_1)}(\phi)]_X \nu_X(p_2)]_W \cdot [\eta_W(q_2)] \\ &= [\phi_W[\eta_W(q_2)]]_X(\nu_X(p_2)). \end{aligned}$$

While

$$\begin{aligned} & [[[\psi_{(F_2, G_2)} \cdot \text{Hom}(\eta, \Gamma_1(\nu))](\phi)]_X(p_2)]_W(q_2) \\ &= [[\psi_{(F_2, G_2)} \cdot (\Gamma_1(\nu) \cdot \phi \cdot \eta)]_X(p_2)]_W(q_2) \\ &= [(\Gamma_1(\nu) \cdot \phi \cdot \eta)_W(q_2)]_X(p_2) \\ &= [[\Gamma_1(\nu)]_W \cdot \phi_W \cdot \eta_W(q_2)]_X(p_2) \\ &= [\phi_W \cdot \eta_W(q_2)]_X(\nu_X(p_2)). \end{aligned}$$

Therefore ψ is a natural equivalence.

A *diagram* in a category \mathbf{C} is a functor $D : \mathbf{I} \rightarrow \mathbf{C}$ whose domain category \mathbf{I} is a small category, [2], [3].

DEFINITION 2. The *limit functor* \bigwedge_D of a diagram D in a category \mathbf{C} is the functor

$$\begin{array}{ccccc} \wedge_D = h_D \cdot K : \mathbf{C} & \longrightarrow & \text{Cat}(\mathbf{I}, \mathbf{C}) & \longrightarrow & \mathbf{S} \\ \Downarrow & & \Downarrow & & \Downarrow \\ W & \longmapsto & K_W & \longmapsto & \text{Hom}(K_W, D) \end{array}$$

where for all $W \in \text{Ob}(\mathbf{C})$, K_W is the constant functor, and a *limit* of the diagram D is a representation (W_0, ϕ) of \wedge_D (cf. [6]), where ϕ is a natural equivalence from h_W to \wedge_D . Dually we can define the colimit functor and colimit of a diagram in \mathbf{C} .

PROPOSITION 3. *There exists a natural equivalence from the limit functor of a diagram D in a category \mathbf{C} to the conjugate of a grounding of \mathbf{C} .*

Proof. Let $D : \mathbf{I} \rightarrow \mathbf{C}$ be a diagram in \mathbf{C} and $\wedge_D = h_D \cdot K$ be the limit functor of D . If we define $G_D : \mathbf{C} \rightarrow \mathbf{S}$ as follows; $G_D(X)$ is a disjoint union $\bigcup_{i \in \mathbf{I}} h^{D(i)}(X)$ of the sets $h^{D(i)}(X)$ for $X \in \text{Ob}(\mathbf{C})$, and $[G_D(p)](f) = p \cdot f$ for $p : X \rightarrow Y$ in \mathbf{C} , $f \in G_D(X)$, then G_D is a grounding of \mathbf{C} . Let define the mapping $\phi_W : A_0(W) \rightarrow G_{D*}(W)$ for $W \in \text{ob}(\mathbf{C})$ and $\varphi \in A_D(W)$ as follows;

$$[\phi_W(\varphi)]_X(x) = x \cdot \varphi \cdot (i)$$

where $x \in h^{D(i)}(X) \subset G_D(X)$ and $i \in \mathbf{I}$. On the other hand if we define the mapping $\phi_0(i) : K_W(i) \rightarrow D(i)$ as follows; $\phi_0(i) = [\phi_{D(i)}][1_{D(i)}]$ for an element $\phi \in G_{D*}(W)$ and $i \in \mathbf{I}$, then since ϕ is a natural transformation, $h^W(D(m)) \cdot \phi_{D(i)} = \phi_{D(i)} \cdot G_D(D(m))$ for $m : i \rightarrow j$ in \mathbf{I} . Hence for $1_{D(i)} : D(i) \rightarrow D(i)$ (in \mathbf{C}),

$$\begin{aligned} [h^W(D(m)) \cdot \phi_{D(i)}](1_{D(i)}) &= h^W(D(m))[\phi_{D(i)}(1_{D(i)})] \\ &= h^W((D(m))\phi_0(i) = D(m)\phi_0(i) \\ [\phi_{D(i)} \cdot G_D(D(m))](1_{D(i)}) &= \phi_{D(i)}(1_{D(i)}) = \phi_0(j), \end{aligned}$$

that is, the diagram

$$\begin{array}{ccc} K_W(i) = W & \longrightarrow & D(i) \\ \parallel & & \downarrow D(m) \\ K_W(j) = W & \longrightarrow & D(j) \end{array}$$

commutes. Therefore ϕ_0 is a natural transformation from K_W to D and if we take $[\phi_W(\phi_0)] = \phi$, then ϕ_W is an isomorphism. For a morphism $\alpha : V \rightarrow W \in \mathbf{C}$, consider the diagram

$$\begin{array}{ccc} L_D(W) & \xrightarrow{\phi_W} & G^*(W) \\ L_D(\alpha) \downarrow & & \downarrow G_{D*}(\alpha) \\ L_D(V) & \xrightarrow{\phi_V} & G_*(V). \end{array} \quad (1)$$

For $x \in (\text{Hom}(D(j), X)) (\subset G_D(x))$ and $\varphi \in L_D(W)$

$$\begin{aligned}
& [[G_{D*}(\alpha)\phi_w](\varphi)]_X(x) \\
& = [G_{D*}(\alpha)[\phi_w(\varphi)]_X(x) \\
& = [h^{(\alpha)} \cdot (\varphi)]_X(x) \\
& = h_X^{(\alpha)}[\phi_w(\varphi)]_X(x) \\
& = h_X^{(\alpha)}(x \cdot \varphi(j)) = x \cdot \varphi(j) \cdot \alpha,
\end{aligned} \tag{2}$$

while

$$\begin{aligned}
& [[\phi_V \cdot L_D^{(\omega)}](\varphi)]_X(x) \\
& = [\phi_V \cdot (\varphi \cdot K(\alpha))]_X(x) \\
& = x \cdot [(\varphi \cdot K(\alpha))(j)] \\
& = x \cdot [\varphi(j) \cdot K(\alpha)(j)] = x \cdot \varphi(j) \cdot \alpha
\end{aligned} \tag{3}$$

From (2) and (3) we prove that the diagram (1) is commutative. Hence we have the natural equivalence $\phi : L_D \cong G_{D*}$. This completes the proof.

3. Couple categories

This section of the paper is a sequel to my paper on couple category [5].

DEFINITION 3. A *coupling* of a cogrounding F and a grounding G of a category \mathbf{C} is a function.

$$m : \bigcup_{(X,Y) \in \mathbf{C} \times \mathbf{C}} F(X) \times G(Y) \longrightarrow \bigcup_{(X,Y) \in \mathbf{C} \times \mathbf{C}} \text{Hom}_{\mathbf{C}}(X, Y)$$

such that

$$\begin{array}{ccc}
F(X) \times G(Y) & \longrightarrow & \text{Hom}_{\mathbf{C}}(X, Y) \\
\Downarrow & & \Downarrow \\
(p, q) & \longmapsto & m(p, q)
\end{array}$$

and for $f : W \longrightarrow X$ and $g : Y \longrightarrow Z$ in \mathbf{C} ,

$$m(F(f)(p), G(g)(q)) = g \cdot m(p \cdot q) \cdot f.$$

PROPOSITION 4. *There is a one-to-one correspondence between couplings m of F and G , and natural transformation $\mu : G \longrightarrow F^*$ defined by*

$$[\mu_Y(q)]_X(p) = m(p, q) \tag{4}$$

for $p \in F(X)$, $q \in G(Y)$.

Proof. Let K be a set of all couplings of F, G , and $\phi : K \longrightarrow \text{Hom}(G, F^*)$ such that $\phi(m) = \mu$ where μ satisfies the condition (4). Then $\phi(m) = \phi(n)$ implies $m = n$. For any $\text{Hom}(G, F^*)$, let m be a function from $\bigcup F(X) \times G(Y)$ to $\bigcup \text{Hom}(X, Y)$ satisfying the condition (4). Then for $f : W \longrightarrow X$ and $g : Y \longrightarrow Z$

$$m(F(f)(p), G(g)(q)) (= [\mu_Z(G(g)(q))]_W(F(f)(p)))$$

$$\begin{aligned}
&= [h_x \cdot \mu_Y(q)]_W(F(f)(p)) \\
&= g[[\mu_Y(q)]_W(F(f)(p))] \\
&= g[h_Y(f) \cdot [\mu_Y(q)]_X(p)] \\
&= g[[\mu_Y(q)]_X(p) \cdot f] \\
&= g \cdot m(p \cdot q) \cdot f
\end{aligned}$$

Hence m is a coupling of F, G . This completes the proof.

A *grounding couple* on a category \mathbf{C} is a triple $F = ({}'F, F', m_F)$ consisting of a cogrounding $'F$, grounding F' of \mathbf{C} and coupling m_F of $'F, F'$.

Morphisms $F = ({}'F, F', m_F) \longrightarrow G = ({}'G, G', m_G)$ of grounding couples are *conjoint transformations* $\eta = ({}'\eta, \eta')$ which are ordered pair of natural transformations $'\eta : {}'F \rightarrow {}'G$, $\eta' : G' \rightarrow F'$, satisfying $m_G(\eta'_X(p), q) = m_F(p, \eta'(q))$ for $p \in {}'F(X)$, $q \in G'(Y)$.

The category, which has these grounding couples for objects and for morphisms the conjoint transformations, is called a *couple category* of \mathbf{C} and we denote it by $\text{Co}(\mathbf{C}, \mathbf{S})$ [5].

Let G be a cogrounding of a category \mathbf{C} and a mapping

$$m : \bigcup_{(X, Y) \in \mathbf{C} \times \mathbf{C}} G(X) \times G^*(Y) \longrightarrow \bigcap_{(X, Y) \in \mathbf{C} \times \mathbf{C}} \text{Hom}_{\mathbf{C}}(X, Y)$$

satisfy $m(p, q) = q_X(p)$ for $p \in G(X)$ and $q \in G^*(Y)$. Then for $f : W \longrightarrow X$ and $g : Y \longrightarrow Z$ in \mathbf{C} , we have

$$\begin{aligned}
m(G(f)(p), G^*(g)(q)) &= [G^*(g)(q)]_W(G^*(f)(p)) \\
&= [h^{(g)}(q)]_W((G(f)(p)) \\
&= h_W^{(g)}[q_W(G(f)(p))] \\
&= q_W(G(f)(p)) \\
&= h_W^{(g)}q_X(p) \\
&= g \cdot q_X(p) \cdot f \\
&= g \cdot m(p \cdot q) \cdot f
\end{aligned}$$

Hence m_G is a coupling of G and G^* . A couple $(G, G^*, m_G) (\in \text{Co}(\mathbf{C}, \mathbf{S}))$ is called a *natural couple*. For two natural couples (G, G^*, m_G) and $(G', G'^*, m_{G'})$ let $\eta : G \longrightarrow G'$ and $\eta^* : G'^* \longrightarrow G^*$ be natural transformations such that $\eta^*_Y(q) = q \cdot \eta$ for $q \in G'^*(Y)$. Then for $p \in G(X)$ and $q \in G'^*(Y)$

$$\begin{aligned}
m(p, \eta^*_Y(q)) &= [\eta^*_Y(q)]_X(p) \\
&= [q \cdot \eta]_X(p) \\
&= q_X \eta_X(p) \\
&= m(\eta_X(p), q).
\end{aligned}$$

Hence (η, η^*) is conjoint transformation from (G, G^*, m_G) to $(G', G'^*, m_{G'})$. Thus we have the full subcategory of $\text{Co}(\mathbf{C}, \mathbf{S})$, which has natural couples for objects and for morphisms conjoint transformations (η, η^*) . The full subcategory of natural couples is called the *natural couple category* and denoted by $\text{Con}(\mathbf{C}, \mathbf{S})$.

LEMMA. A functor $T : \mathbf{A} \rightarrow \mathbf{B}$ is an equivalence if and only if there is a functor $S : \mathbf{B} \rightarrow \mathbf{A}$ together with natural equivalences

$$\varphi : \mathbf{1}_{\mathbf{B}} \cong T \cdot S \text{ and } \psi : ST \cong \mathbf{1}_{\mathbf{A}}, \quad [4]$$

PROPOSITION 5. The functor $\phi : \text{Cat}(\mathbf{C}^*, \mathbf{S}) \rightarrow \text{Con}(\mathbf{C}, \mathbf{S})$ such that $\phi(G) = (G, G^*, m_G)$ for all $G \in \text{Ob}(\text{Cat}(\mathbf{C}^*, \mathbf{S}))$ and $\phi(\eta) = (\eta, \eta^*)$ for all morphisms η in $\text{Cat}(\mathbf{C}^*, \mathbf{S})$ is an equivalence functor.

Proof. Let $\phi : \text{Con}(\mathbf{C}, \mathbf{S}) \rightarrow \text{Cat}(\mathbf{C}^*, \mathbf{S})$ be a functor such that $\phi(G, G^*, m_G) = G$ and $\phi(\eta, \eta^*) = \eta$ for $(\eta, \eta^*) : (G, G^*, m_G) \rightarrow (G', G'^*, m_{G'})$ and $\varphi : \mathbf{1}_{\text{Cat}(\mathbf{C}^*, \mathbf{S})} \rightarrow \phi \cdot \phi$. Then $\varphi_G : \mathbf{1}_{\text{Cat}(\mathbf{C}^*, \mathbf{S})}(G) \rightarrow \phi \cdot \phi(G) = G$ is an identity morphism and φ is a natural equivalence. Similarly we can obtain the natural equivalence $\varphi' : \phi \cdot \phi \sim \mathbf{1}_{\text{Con}(\mathbf{C}, \mathbf{S})}$. By the lemma the functor ϕ is an equivalence.

COROLLARY 1. The functor category $\text{Cat}(\mathbf{C}^*, \mathbf{S})$ is embedded into the couple category $\text{Co}(\mathbf{C}, \mathbf{S})$ and \mathbf{C} is embedded into $\text{Con}(\mathbf{C}, \mathbf{S})$.

A grounding couple $F = (F, F', m_F)$ is said to be *separated* if $\mu' : F' \rightarrow F^*$ and $\mu : F \rightarrow F_*'$ are monomorphic ([5], [3]). A cogrounding G is said to be *separated* if its natural couple (G, G^*, m_G) is separated.

DEFINITION 4. A subcategory \mathbf{A} of a category \mathbf{B} is called a *separating subcategory* if for every two distinct morphisms $f, g : X \rightarrow Y$ in \mathbf{B} , there exist an object Z in \mathbf{A} and $h : Y \rightarrow Z$ such that $h \cdot f \neq h \cdot g$.

We can know that a cogrounding F of \mathbf{C} is separated iff F and F^* are reflexive. By Proposition 5, we have the

COROLLARY 2. The natural couple category $\text{Con}(\mathbf{C}, \mathbf{S})$ is separated iff $\text{Cat}(\mathbf{C}, \mathbf{S})$ is separated.

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