

ON PROJECTIVE LIMITS IN PRIMITIVE CATEGORIES

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1. Introduction

Let \mathbf{C} be a category and let \otimes be a functor from the product category $\mathbf{C} \times \mathbf{C}$ into \mathbf{C} such that the functors

$$\otimes(\otimes \times \mathbf{C}): (\mathbf{C} \times \mathbf{C}) \times \mathbf{C} \rightarrow \mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$$

and

$$\otimes(\mathbf{C} \times \otimes): \mathbf{C} \times (\mathbf{C} \times \mathbf{C}) \rightarrow \mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$$

are identical and I be an object of \mathbf{C} such that the functors obtained by \otimes , fixing the variables in I , are identical to $1_{\mathbf{C}}$. Then a triple (\mathbf{C}, \otimes, I) is called a multiplicative category of \mathbf{C} .

For a multiplicative category (\mathbf{C}, \otimes, I) of \mathbf{C} , a triple (A, m, k) is called a tensor object or \otimes -object of \mathbf{C} , where A is an object of \mathbf{C} , m is a morphism from $A \otimes A$ to A such that $m(m \otimes A) = m(A \otimes m)$ and k is a morphism from I to A such that $m(k \otimes A) = m(A \otimes k) = 1_A$.

If we now define morphisms f from (A, m, k) to (A', m', k') as follows; for $f: A \rightarrow A'$,

$$f \cdot m = m' \cdot (f \otimes f), \quad k' = f \cdot k$$

and their composition is the composition in \mathbf{C} , then the class of objects, which is the class of all \otimes -objects of \mathbf{C} , together with the class of morphisms $f: (A, m, k) \rightarrow (A', m', k')$, and the composition in \mathbf{C} forms a category. This category is called a primitive category and we denote it by $\Phi(\mathbf{C}, \otimes, I)$. The functor $\psi: \Phi(\mathbf{C}, \otimes, I) \rightarrow \mathbf{C}$ is forgetful if for every object (A, m, k) and $f: (A, m, k) \rightarrow (A', m', k')$, $\psi[(A, m, k)] = A$, $\psi(f) = f$ (For other equivalent definitions see [1], [3]). The notion of multiplicative category was established by J. Bénabou [2] and the structure of the primitive categories was studied by B. Eckmann and P. J. Hilton [4].

The purpose of this note is to establish the simple results on the projective limits in primitive categories.

2. Primitive categories

Let (\mathbf{C}, \otimes, I) and $(\bar{\mathbf{C}}, \bar{\otimes}, \bar{I})$ be two multiplicative categories of \mathbf{C} and $\bar{\mathbf{C}}$ respectively. A triple (G, φ, δ) is called a morphism from (\mathbf{C}, \otimes, I) to $(\bar{\mathbf{C}}, \bar{\otimes}, \bar{I})$, where G is a functor from \mathbf{C} into $\bar{\mathbf{C}}$, φ is a natural transformation from $\otimes(G \times G)$ to $G \bar{\otimes}$ such that for all triple (A, B, C) of objects of \mathbf{C} ,

$$\varphi_{(A \otimes B, C)} \cdot (\varphi_{(A, B)} \otimes G(C)) = \varphi_{(A, B \otimes C)} \cdot (G(A) \bar{\otimes} \varphi_{(B, C)})$$

and $\delta: \bar{I} \rightarrow G(I)$ (in $\bar{\mathbf{C}}$) is a morphism such that for all objects $A \in \mathbf{C}$,

$$\varphi_{(A, \delta)} \cdot (\delta \otimes G(A)) = \mathbf{1}_{G(A)} = \varphi_{(A, I)} \cdot (G(A) \times \delta).$$

For two morphisms $(G, \varphi, \delta) : (\mathbf{C}, \otimes, I) \rightarrow (\bar{\mathbf{C}}, \otimes, \bar{I})$ and $(\bar{G}, \bar{\varphi}, \bar{\delta}) : (\bar{\mathbf{C}} \otimes, \bar{I}) \rightarrow (\bar{\bar{\mathbf{C}}}, \otimes, \bar{\bar{I}})$, we shall define the composition of (G, φ, δ) and $(\bar{G}, \bar{\varphi}, \bar{\delta})$ as follows;

$$(\bar{G} \cdot G, (\bar{G} * \varphi) (\bar{\varphi} * (G \times G)), \bar{G}(\delta) \cdot \bar{\delta})$$

where,

$$(G * \varphi) (\bar{\varphi} * (G \times G)) : [\otimes (\bar{G} \times \bar{G})] (G \times G) = \otimes (\bar{G} G \times \bar{G} G) \rightarrow [\bar{G} \otimes] (G \times G) \\ \rightarrow \bar{G} (G \otimes).$$

THEOREM 1. Let $(G, \varphi, \delta) : (\mathbf{C}, \otimes, I) \rightarrow (\bar{\mathbf{C}}, \otimes, \bar{I})$ and $\phi : \Phi(\mathbf{C}, \otimes, I) \rightarrow \mathbf{C}$ and $\phi' : \Phi(\bar{\mathbf{C}}, \otimes, \bar{I}) \rightarrow \bar{\mathbf{C}}$ be forgetful functors. If for a morphism (G, φ, δ) , every \otimes -object (A, m, k) and $f : (A, m, k) \rightarrow (A', m', k')$,

$$\Phi(G, \varphi, \delta) [(A, m, k)] = (G(A), G(m) \varphi_{(A, \delta)}, G(k) \delta) \in \Phi(\bar{\mathbf{C}}, \otimes, \bar{I})$$

and

$$\Phi(G, \varphi, \delta) [f] = G(f) : (G(A), G(m) \varphi_{(A, \delta)}, G(k) \delta) \\ \rightarrow (G(A'), G(m') \varphi_{(A', \delta)}, G(k') \delta) (\in \Phi(\bar{\mathbf{C}}, \otimes, \bar{I})),$$

then $\Phi(G, \varphi, \delta)$ is a functor from $\Phi(\mathbf{C}, \otimes, I)$ into $\Phi(\bar{\mathbf{C}}, \otimes, \bar{I})$ such that $G\phi = \phi' \Phi(G, \varphi, \delta)$. And $\Phi(G, \varphi, \delta)$ is a faithful functor if and only if so is G .

Proof. For two morphisms $f : (A, m, k) \rightarrow (A', m', k')$ and $f' : (A', m', k') \rightarrow (A'', m'', k'')$ in $\Phi(\mathbf{C}, \otimes, I)$,

$$\Phi(G, \varphi, \delta) [f' \cdot f] = G[f' \cdot f] = G[f'] \cdot G[f] \\ = \Phi(G, \varphi, \delta) [f'] \cdot \Phi(G, \varphi, \delta) [f].$$

$$\Phi(G, \varphi, \delta) [\mathbf{1}_A] = G[\mathbf{1}_A] = \mathbf{1}_{G(A)} = \mathbf{1}_{\Phi(G, \varphi, \delta) [(A, m, k)]},$$

and

$$G \cdot \phi [(A, m, k)] = G[A],$$

$$\phi' \cdot \Phi(G, \varphi, \delta) [(A, m, k)] = \phi' (G(A), G(m) \varphi_{(A, \delta)}, G(k) \delta) = G(A)$$

Hence

$$G\phi = \phi' \Phi(G, \varphi, \delta).$$

Since the function

$$\chi : \text{Hom}_{\Phi(\mathbf{C}, \otimes, I)} [(A, m, k), (A', m', k')] \longrightarrow \\ \text{Hom}_{\Phi(\bar{\mathbf{C}}, \otimes, \bar{I})} [(G(A), G(m) \varphi_{(A, \delta)}, G(k) \delta), (G(A'), G(m') \varphi_{(A', \delta)}, G(k') \delta)]$$

is univalent iff $\chi_0 : \text{Hom}_{\mathbf{C}} (A, A') \rightarrow \text{Hom}_{\bar{\mathbf{C}}} (G(A), G(A'))$ is univalent, $\Phi(G, \varphi, \delta)$ is faithful iff so is G .

DEFINITION. For a morphism $(G, \varphi, \delta) : (\mathbf{C}, \otimes, I) \rightarrow (\bar{\mathbf{C}}, \otimes, \bar{I})$, if $G\otimes = \otimes(G \times G)$ and $\bar{I} = G(I)$, i. e. $\varphi = \mathbf{1}_{\otimes}, \delta = \mathbf{1}_{\bar{I}}$ then we say that the functor G commutes to tensor product $[\mathbf{1}]$.

In this case we denote the morphism $(G, 1_{\mathcal{C}}, 1_I)$ by G and $\Phi(G, 1_{\mathcal{C}}, 1_I) = \Phi(G)$. Then we have

$$\Phi(G)[(A, m, k)] = (G(A), G(m), G(k)).$$

For a functor $G: \mathbf{C} \rightarrow \bar{\mathbf{C}}$, a morphism $T: G(A) \rightarrow \bar{A}$ in \mathbf{C} is called a universal morphism if for all morphism $f: G(A') \rightarrow \bar{A}$ there is a unique morphism $g: A' \rightarrow A$ such that $f = T \cdot G(g)$.

LEMMA. Let (\mathbf{C}, \otimes, I) and $(\bar{\mathbf{C}}, \otimes, \bar{I})$ be two multiplicative categories and $G: \mathbf{C} \rightarrow \bar{\mathbf{C}}$ be a functor that is commutative to tensor product. Then if for a \otimes -object $(\bar{A}, \bar{m}, \bar{k})$ of $\Phi(\bar{\mathbf{C}}, \otimes, \bar{I})$ the morphism $T: G(A) \rightarrow \bar{A}$ in $\bar{\mathbf{C}}$ is universal, there exist unique morphisms $m: A \otimes A \rightarrow A$ and $k: I \rightarrow A$ such that $(A, m, k) \in \Phi(\mathbf{C}, \otimes, I)$ and $T: \Phi(G)[(A, m, k)] \rightarrow (\bar{A}, \bar{m}, \bar{k})$ is universal.

Proof. Since T is a universal morphism, for a morphism $f: G(A \otimes A) \rightarrow \bar{A}$ there is a unique morphism $m: A \otimes A \rightarrow A$ such that

$$\begin{array}{ccc} G(A) & \xrightarrow{T} & \bar{A} \\ \uparrow G(m) & \nearrow f & \\ G(A \otimes A) & & \end{array}$$

commutes. Similarly, there is a unique $k: I \rightarrow A$ with $T \cdot G(k) = \bar{k}$. By the uniqueness, $m \otimes (m \otimes A) = m \otimes (A \otimes m)$ and $m(A \otimes k) = m(k \otimes A) = 1_A$. If T is universal then by the definition of the morphism $T: (G(A), G(m), G(k)) \rightarrow (\bar{A}, \bar{m}, \bar{k})$ we can easily know that the latter is also universal.

3. Projective limits

For any two categories \mathbf{A} and \mathbf{B} , let $\mathbf{Cat}(\mathbf{A}, \mathbf{B})$ denote the category of all covariant functors from \mathbf{A} into \mathbf{B} . Let $K_A: \mathbf{J} \rightarrow \mathbf{C}$ be the constant functor which assigns each object of \mathbf{J} to the fixed object A of \mathbf{C} and each morphism of \mathbf{J} to 1_A and let $E_J: \mathbf{C} \rightarrow \mathbf{Cat}(\mathbf{J}, \mathbf{C})$ be a functor such that $E_J(A) = K_A$ for all $A \in \mathbf{C}$. For any functor $F \in \mathbf{Cat}(\mathbf{J}, \mathbf{C})$ a natural transformation $\rho: E_J(A) = K_A \rightarrow F$ is called the projective limit of F if for any natural transformation $\sigma: E_J(A') \rightarrow F$ there exist the unique morphism $\mu: A' \rightarrow A$ in \mathbf{C} such that

$$\sigma = \rho \cdot E_J(\mu).$$

Now we suppose that for any two objects (A, m, k) and (A', m', k') of $\Phi(\mathbf{C}, \otimes, I)$,

$$(A \otimes A') \otimes (A \otimes A') = (A \otimes A) \otimes (A' \otimes A').$$

Let us denote a morphism from $(A \otimes A') \otimes (A \otimes A')$ to $A \otimes A'$ by $m \otimes m'$ and a morphism for I to $A \otimes A'$ by $k \otimes k'$. Then the morphism $m \otimes m'$ is associative and

$$\begin{aligned} (m \otimes m')[(k \otimes k') \otimes (A \otimes A')] &= [m(k \otimes A)] \otimes [m'(k' \otimes A')] \\ &= 1_A \otimes 1_{A'} \\ &= 1_{A \otimes A'}. \end{aligned}$$

Similarly,

$$(m \otimes m')[(A \otimes A') \otimes (k \otimes k')] = 1_{A \otimes A'}.$$

Hence, in this case we have a functor

$$\odot: \Phi(\mathbf{C}, \otimes, I) \times \Phi(\mathbf{C}, \otimes, I) \rightarrow \Phi(\mathbf{C}, \otimes, I)$$

such that

$$\odot[(A, m, k), (A', m', k')] = (A, m, k) \odot (A', m', k') = (A \otimes A', m \otimes m', k \otimes k')$$

and for $f: (A, m, k) \rightarrow (B, n, h)$ and $f': (A', m', k') \rightarrow (B', n', h')$

$$\odot[(f, f')] = f \otimes f': (A \otimes A', m \otimes m', k \otimes k') \rightarrow (B \otimes B', n \otimes n', h \otimes h').$$

The functor \odot is associative and

$$(A, m, k) \odot (I, 1_I, 1_I) = (A, m, k).$$

Therefore given a multiplicative category (\mathbf{C}, \otimes, I) we can also define a multiplicative category $\{\Phi(\mathbf{C}, \otimes, I), \odot, \mathbf{I} = (I, 1_I, 1_I)\}$ of a primitive category $\Phi(\mathbf{C}, \otimes, I)$. And if

$$\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)) \times \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)) = \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I) \times \Phi(\mathbf{C}, \otimes, I)),$$

then we can define a functor

$$\mathbf{Cat}(\mathbf{J}, \odot): \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)) \times \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)) \rightarrow \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I))$$

as follows:

$$\mathbf{Cat}(\mathbf{J}, \odot)[(\phi, \psi)] = \phi \odot \psi,$$

$$[\phi \odot \psi](j) = \phi(j) \odot \psi(j),$$

for all $j \in \mathbf{J}$. Since the functor \odot is associative, so is the functor $\mathbf{Cat}(\mathbf{J}, \odot)$. Let us denote $\mathbf{Cat}(\mathbf{J}, \odot)$ by \odot .

If $E_{\mathbf{J}}: \Phi(\mathbf{C}, \otimes, I) \rightarrow \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I))$ be a functor such that for each object (A, m, k) of $\Phi(\mathbf{C}, \otimes, I)$ and for all objects j of \mathbf{J}

$$[E_{\mathbf{J}}(A, m, k)](j) = (A, m, k),$$

then for $E_{\mathbf{J}}(I, 1_I, 1_I) = \bar{\mathbf{I}}$ and $\phi \in \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I))$ and $\phi(j) = (B, n, h)$,

$$(\bar{\mathbf{I}} \odot \phi)(j) = (I \otimes B, 1_I \otimes n, 1_I \otimes h) = (B, n, h) = \phi(j).$$

Similarly, $\phi \odot \bar{\mathbf{I}} = \phi$. Hence we have a multiplicative category $[\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)), \odot, \bar{\mathbf{I}}]$ of a category $\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I))$.

For each ordered pairs $\{(A, m, k), (A', m', k')\}$ of objects of $\Phi(\mathbf{C}, \otimes, I)$ and for each $j \in \mathbf{J}$,

$$\begin{aligned} \odot \cdot E_{\mathbf{J}} \times E_{\mathbf{J}}[\{(A, m, k), (A', m', k')\}](j) &= [E_{\mathbf{J}}(A, m, k) \odot E_{\mathbf{J}}(A', m', k')](j) \\ &= E_{\mathbf{J}}(A, m, k)(j) \odot E_{\mathbf{J}}(A', m', k')(j) \\ &= (A, m, k) \odot (A', m', k') \\ &= (A \otimes A', m \otimes m', k \otimes k'). \end{aligned}$$

Similarly,

$$E_{\mathbf{J}} \odot [\{(A, m, k), (A', m', k')\}](j) = (A \otimes A', m \otimes m', k \otimes k').$$

The functor \odot commutes to the tensor product.

Let us denote

$$(E_{\mathbf{J}}, \mathbf{1}_{\odot}, \mathbf{1}_{\bar{\mathbf{I}}}) : [\Phi(\bar{\mathbf{C}}, \otimes, I), \odot, (I, \mathbf{1}_I, \mathbf{1}_I)] \rightarrow [\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)), \odot, \bar{\mathbf{I}}]$$

by $E_{\mathbf{J}}$. Then we have a functor

$$\Phi(E_{\mathbf{J}}) : \Phi[\Phi(\mathbf{C}, \otimes, I), \odot, (I, \mathbf{1}_I, \mathbf{1}_I)] \rightarrow \Phi[\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)), \odot, \bar{\mathbf{I}}]$$

such that for each $[(A, m, k), \bar{m}, \bar{k}] \in \Phi[\Phi(\mathbf{C}, \otimes, I), \odot, (I, \mathbf{1}_I, \mathbf{1}_I)]$, where

$$\bar{m} : (A, m, k) \odot (A, m, k) \rightarrow (A, m, k) \text{ and } \bar{k} : (I, \mathbf{1}_I, \mathbf{1}_I) \rightarrow (A, m, k).$$

Hence we have the

THEOREM 2. *Let (\mathbf{C}, \otimes, I) be a multiplicative category of \mathbf{C} and $(A \otimes A') \otimes (A \otimes A') = (A \otimes A) \otimes (A' \otimes A')$. Then for $\phi \in \mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I))$ and $(\phi, m_{\phi}, k_{\phi}) \in \Phi[\mathbf{Cat}(\mathbf{J}, \Phi(\mathbf{C}, \otimes, I)), \odot, \bar{\mathbf{I}}]$, the functor ϕ has a projective limit $\rho : E_{\mathbf{J}}[(A, m, k)] \rightarrow \phi$ if and only if $\Phi(\rho) : \Phi(E_{\mathbf{J}})[(A, m, k), \bar{m}, \bar{k}] \rightarrow (\phi, m_{\phi}, k_{\phi})$ is universal.*

Proof. Let $\Phi(\rho)$ be universal, then by the definition of $\Phi(\rho)$, ρ is universal. Conversely, if ρ has a projective limit ρ , then by the Lemma $\Phi(\rho)$ is universal only if ρ is universal.

References

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