Analyses of A Lightweight Stream Cipher for RFID Encryption Model

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RFID 암호 모델을 위한 경량화 스트림 암호 방식의 해석

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요 약

WG-7 is a stream cipher based on WG Stream Cipher and is designed by Y. Luo, Q. Chai, G. Gong, and X. Lai in 2010. This cipher is designed to implement in low cost and lightweight application such as RFID tags. In this paper, we survey and compare cryptographic module such as stream and block cipher. We can estimate security performance suitable to system.

I. Introduction

For security applications in wireless sensor networks (WSNs), choosing best algorithms in terms of energy-efficiency and of small-storage requirements is a real challenge because the sensor networks must be design with limited resources. Sensor networks are made by the tremendous advances and convergence of micro-electro-mechanical systems (MEMS), wireless communication technologies and digital electronics. Sensor networks are composed of a large number of tiny devices or sensors which monitor their surrounding area to measure environmental information, detect to movements, vibrations, etc [1].

We take into consideration the problem of efficiently generating sequences in hardware for use in certain cryptographic algorithms. We show that sequences generated by linear feedback shift (LFSRs) registers in stream cipher can tailored to pursuit the appropriate be algorithms. For hardware implementation,

this reduces both time and chip area. WG-7 is proposed as a fast, lightweight and secure stream cipher inspired by family of WG stream ciphers design principles. WG is a synchronous stream to ECRYPT call cipher submitted for stream ciphers. WG-7 and WG are hardware-oriented stream ciphers that use a word-oriented Linear Feedback Shift Register (LFSR) and a filter function based on Welch-Gong (WG) transformation [1].

II. Characteristics of LFSR

The results (computed using the skyeye + eSimu tools) concerning the number of cycles required to perform all the tests are summarized in the table 1.

			nJ/byte				nJ/key	nJ/IV	nJ/byte
Algo.	Key	IV	Stream	40 bytes	576 bytes	1500 bytes	Key setup	IV setup	agility
Сору	80	80	38.32	60.85	16.84	142.07	70.54	67.29	145.35
RC4	128	0	465.17	9843.25	948.49	542.06	1243.66	379636.24	354.43
SNOW v2.0	128	128	438.34	1093.46	280.59	414.20	2656.66	41749.08	365.26
AES CTR	128	128	3587.00	2197.89	3437.36	3384.26	11378.81	2861.89	3499.45
DRAGON	128	128	514.26	2912.69	1144.53	1064.58	6846.80	74109.24	575.67
HC-256	128	128	471.39	102473.69	7577.28	3112.48	2540.02	2705307.85	864.11
HC-128	128	128	342.29	24264.78	1838.04	897.21	2540.20	950661.16	559.97
LEX	128	128	804.03	1186.80	670.42	714.16	8250.66	23850.60	868.13
Phelix	128	128	421.15	1470.51	461.14	454.71	20622.78	35111.32	461.26
Py	128	64	3894.22	5822.63	827.52	1101.62	145194.31	154181.03	1141.65
Руру	128	56	817.35	6008.43	1859.92	1361.36	145194.15	161834.16	1300.67
Salsa20	128	64	952.19	1394.11	907.17	1275.82	6884.19	2215.93	1268.12
SOSEMANUK	128	64	247.93	6727.04	648.50	528.97	286119.29	20860.01	365.30

Table 1. Number of CPU cycles for the stream ciphers using the testing framework

Table 2. Number of nJ for the stream ciphers using the testing framework

			cycles/byte				cycles/key	cycles/IV	cycles/byte
Algo.	Key	IV	Stream	40 bytes	576 bytes	1500 bytes	Key setup	IV setup	agility
Сору	80	80	2.19	3.72	1.00	7.58	4.40	4.19	7.78
RC4	128	0	26.97	610.95	58.53	33.29	76.41	23581.61	21.24
SNOW v2.0	128	128	25.08	66.38	16.82	23.71	163.41	2273.35	20.87
AES CTR	128	128	206.19	131.52	198.73	195.76	636.49	157.52	202.23
DRAGON	128	128	30.89	177.05	69.76	64.91	421.42	4497.61	33.60
HC-256	128	128	27.00	6044.76	446.11	183.17	141.75	198126.10	49.30
HC-128	128	128	19.35	1484.72	112.12	53.70	141.76	58194.93	31.67
LEX	128	128	47.07	71.41	40.32	41.92	501.41	1415.57	50.71
Phelix	128	128	25.61	90.15	28.36	26.77	1271.42	2154.61	26.99
Py	128	64	214.25	349.23	47.58	60.88	$7713.83 \ z$	9327.43	64.40
Руру	128	56	44.78	360.95	103.91	74.72	7713.82	9660.11	73.46
Salsa20	128	64	57.54	84.57	55.05	73.07	367.70	118.07	72.60
SOSEMANUK	128	64	14.81	385.63	37.95	30.48	16374.01	1264.09	20.78

IV. Conclusion

The comparison between the results obtained in concerning the performances of block ciphers using several modes of operation and the stream ciphers presented will be more pertinent for its demands. We also can estimate the general characteristics for estimating performance produced by the addition of a stream cipher in a real sensor communication environment.

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